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| FR 001 | | |  | |  | | | |  | | | | ge | | | | There have been many papers in the last few months exploring the coroutines design space including P1745, P1713, P1662, P1663, P1344. These evolutions seems worth pursuing, notably because coroutines cannot be used in environments where heap allocations are impossible. However, it is not clear that there will be a reasonable non-breaking path towards these evolutions as they almost all assume a separation of coroutine\_handle and suspend points (as described in P1745), which would be a breaking change if adopted after the publication of the C++20 standard. | | | | Ensure that the final specification of coroutines do not preclude further extensions including those outlined in the paper P1745. | | | | |  | | | | | |
| DE 002 | | |  | |  | | | |  | | | |  | | | | P1020R1 ((https://wg21.link/p1020r1) proposed some helper functions for smart pointers that do not necessarily initialize values. Providing better performance but bring more risk.  However, the technical term "default initialization" used in their name does not always initialize (as value initialization does). Therefore, it is highly confusing for programmers to call an ordinary function  make\_unique\_default\_init()  make\_shared\_default\_init()  allocate\_shared\_default\_init()  If you ask what is safer:  make\_shared() or  make\_shared\_default\_init() almost every programmer says the latter.  For this reason, the more risky version should have a name that demonstrates the risk rather than giving a false sense of safety. | | | | We strongly request to rename them to  make\_unique\_nonvalue\_init()  make\_shared\_nonvalue\_init()  allocate\_shared\_nonvalue\_init() (with or without a \_ between non and value). | | | | |  | | | | |
| DE 003 | | |  | |  | | | |  | | | |  | | | | C++ has a key problem being far too complicated for beginners. Now we make it even worse. Instead of size(), which is self-explanatory but creates warnings, we introduce ssize(), which avoids the warnings but is way harder to teach and NOT self-explanatory.  There was a nice alternative proposal:  count() as proposed in https://wg21.link/p1764.  This is both self-explanatory and avoids the confusing warnings. | | | | We request to use the name count() or at least remove ssize() otherwise. | | | | |  | | | | |
| CZ 004 | | |  | |  | | | |  | | | | ge | | | | With introducing of P1152 (Deprecating volatile) we deprecated most usages of volatile keyword. We should reflect the changes in library by adopting the changes there. | | | | Adopt P1831 (deprecating volatile: library) | | | | |  | | | | |
| CZ 005 | | |  | |  | | | |  | | | | ge | | | | Please consider adopting P0593 (Implicit creation of objects for low-level object manipulation) to fix long-standing problems with low level object creation and UB. | | | | Adopt P0593 (Implicit creation of objects for low-level object manipulation) | | | | |  | | | | |
| RU 006 | |  | | | | | [atomics.types.operations] | | | | Paragraph 2 | | | | te | | | | Fix default initialization of C++ atomic variables to fit user expectations. Default constructor of std::atomic has to initialize the atomic object with value of T{}. Default constructor of std::atomic\_flag has to initialize to clear state. | | | | | Apply wording from P0883. | | | | |  | | | | | |
| RU 007 | |  | | | | | [basic.life] | | | | Paragraph 8, bullet 3 | | | | te | | | | In many cases it is impossible to use a pointer returned from placement new or std::launder (for example in std::vector, std::variant, std::optional, std::uninitialized\_\*+std::destroy\*).  Because of that issue all the standard libraries have undefined behaviors in widely used types. The only way to fix that issue is to adjust the lifetime rules to auto-launder the placement new.  Dropping the “const“ and “reference” requirement from paragraph 8 [basic.life] removes UB from std::vector and std::optional. Additional removing of the “potentially-overlapping” requirement removes UB from std::variant. | | | | | Apply the following changes to the [basic.life] paragraph 8 bullet 3:  – the type of the original object is not const-qualified~~, and, if a class type, does not contain any non-static data member whose type is const-qualified or a reference type~~, and  Optionally remove the to the bullet 4 from [basic.life] paragraph 8:  – ~~neither the original object nor the new object is a potentially-overlapping subobject ([intro.object])~~. | | | | |  | | | | | |
| RU 008 | |  | | | | | [class.cdtor] | | | | Paragraph 2 | | | | te | | | | Clarify that \*this in constructors could not alias with reference input parameters of the same type. This is useful for optimizing all the constructors of all the classes. | | | | | Apply the following changes to the [class.cdtor] paragraph 2:  During the construction of an object, if the value of the object or any of its subobjects is accessed through a glvalue that is not obtained, directly or indirectly, from the constructor's this pointer, the value of the object or subobject thus obtained is ~~unspecified~~undefined. | | | | |  | | | | | |
| RU 009 | |  | | | | | [compliance] | | | |  | | | | te | | | | Make <compare> freestanding. | | | | | Apply wording from P1855. | | | | |  | | | | | |
| FI 010 | | |  | | [compliance] | | | |  | | | |  | | | | | <coroutine> uses <compare>, <compare> is not freestanding | | | | Adopt P1855R0 | | | | |  | | | | | |
| RU 011 | |  | | | | | [dcl.fct.def.general] | | | | Paragraph 8 | | | | te | | | | Make \_\_func\_\_ usable in constant expressions. Resolve CWG 2362 (<http://www.open-std.org/jtc1/sc22/wg21/docs/cwg_active.html#2362>) | | | | | Change the definition of \_\_func\_\_ to a static constexpr variable:  static constexpr char \_\_func\_\_[] = "function-name"; | | | | |  | | | | | |
| RU 012 | |  | | | | | [expr.prim.req] | | | | Paragraph 6 | | | | te | | | | Accessing an incomplete type in requires-expression results in substitution failure. This is a dangerous approach that provokes odr-violations. Making the program ill-formed is a better approach that protects from hard detectable odr-violations. | | | | | Add to the [expr.prim.req] paragraph 6:  In such cases, the requires-expression evaluates to false; it does not cause the program to be ill-formed if the requires-expression does not refer to incomplete type. | | | | |  | | | | | |
| RU 013 | |  | | | | | [string.cons] | | | | Paragraph 30 | | | | te | | | | Because of the implicit conversion of arithmetic types it is error prone to use the basic\_string::operator=(charT c):  double d = 3.14;  std::string s;  s = d; // Compiles  Make sure that the program is ill-formed if an implicit conversion from arithmetic type happens while assigning to std::basic\_string. Or at least make sure that the program is ill-formed if an implicit conversion from floating point type happens while assigning to std::basic\_string. | | | | | Apply the following changes to the [string.cons] paragraph 30:  template <class T>  constexpr basic\_string& operator=(~~char~~T c);  Constraints:  is\_same\_v<T, charT> is true  Effects: Equivalent to:  return \*this = basic\_string\_view<charT, traits>(addressof(c), 1);  Change the [basic.string] synopsys:  template <class T>  constexpr basic\_string& operator=(~~char~~T c); | | | | |  | | | | | |
| RU 014 | |  | | | | | [string.view.template] | | | |  | | | | te | | | | Forbid assigning an rvalue std::basic\_string to std::basic\_string\_view. It will allow to detect errors at compile time, rather than at runtime. Resolve LWG 3068 (https://cplusplus.github.io/LWG/issue3068) | | | | | Add to the [string.view.template] basic\_string\_view synopsis:  template <class A>  basic\_string\_view& operator=(const basic\_string<charT, traits, A>&&) = delete; | | | | |  | | | | | |
| FI 015 | | |  | | [support.limits.general] | | | |  | | | |  | | | | Various feature-testing macros are missing; examples are ones for span and concepts. | | | | Discuss, modify and adopt Barry Revzin’s revision of SD-6, and adopt the resulting changes into the IS working draft. | | | | |  | | | | | |
| RU 016 | |  | | | | | [temp.alias] | | | |  | | | | te | | | | Fix pack expansion into fixed alias template parameter list by resolving CWG 1430 (<http://www.open-std.org/jtc1/sc22/wg21/docs/cwg_active.html#1430>).  Desired outcome is to allow the following code to compile:  #include <type\_traits>  template <typename... Xs>  auto f (Xs...)  -> std::invoke\_result\_t<Xs...>; | | | | |  | | | | |  | | | | | |
| GB 017 | | |  | | 01 | | | |  | | | | Te | | | | This is an explicit request to address all open issues on the core and library list. | | | |  | | | | |  | | | | |
| DE 018 | | |  | | 01.02 31.8 31.8.1 31.8.2 31.8.3 31.8.4 31.9  31.10 | | | |  | | | |  | | | | P0883: Fixing Atomic Initialization (https://wg21.link/p0883) went through SG1 and LEWG to be adopted for C++17. This is an important fix for the broken atomic initialization. | | | | Apply this fix through LWG.  According to N4830 this paper now applies to:  31.2 Header <atomic> synopsis [atomics.syn]  31.8 Class template atomic [atomics.types.generic]  31.8.1 Operations on atomic types [atomics.types.operations]:  31.8.2 Specializations for integers [atomics.types.int]  31.8.3 Specializations for floating-point types [atomics.types.float]  31.8.4 Partial specialization for pointers [atomics.types.pointer]  31.9 Non-member functions [atomics.nonmembers]  31.10 Flag type and operations [atomics.flag] | | | | |  | | | | |
| US 019 | | |  | | 02 | | | | p1 | | | | TE | | | | We normatively refer to ISO/IEC 9899:2011 which is not the latest version of the C Standard. | | | | We should be tracking against ISO/IEC 9899:2018 (aka, C17) instead. C17 is near-identical to C11 in that it was a bugfix release that only resolved DRs, with one exception (ATOMIC\_VAR\_INIT was deprecated in C17) | | | | |  | | | | |
| US 020 | | |  | | 02 | | | | p1 | | | | TE | | | | We normatively refer to ISO/IEC 9899:2011 which is not the latest version of the C Standard. | | | | We should be tracking against ISO/IEC 9899:2018 (aka, C17) instead. C17 is near-identical to C11 in that it was a bugfix release that only resolved DRs, with one exception (ATOMIC\_VAR\_INIT was deprecated in C17) | | | | |  | | | | |
| US 021 | | |  | | 03.06 [denfs.block] | | | |  | | | | Ed | | | | [intro.defs] defines 'block' only as a verb (3.6 [defns.block]), for the memory model. We also use 'block' to define a brace-enclosed region of code, per 8.3 [stmt.block]. The index entry for 'block' refers to the [defns.block] definition, but has sub-entries that apply only to the [stmt.block] definition. | | | | Add a second definition for block as a brace-enclosed region of code, and update misdirected index references to the new definition. Index entries for 'block' will want to clarify between the verb and the noun. | | | | |  | | | | |
| GB 022 | | |  | | 04.01 | | | |  | | | | Ed | | | | Allow access to standard library names via import  [intro.compliance] 4.1 p5  "The names defined in the library have namespace scope (9.7). A C++ translation unit (5.2) obtains access to these names by including the appropriate standard library header (15.2)." | | | | Add that standard library names can also be accessed by importing the appropriate header unit.:  … library header (15.2)<ins> or importing the appropriate standard library named header unit (10.3), [headers] 16.5.1.2 p4</ins>. | | | | |  | | | | |
| JP1 023 | | |  | | 04.02 | | | | p5 | | | | ed | | | | There are also many "[Note 1 to entry:" style notes in the document as well as "[Note:". They should also be referred to. | | | | Add "[Note 1 to entry:" and terminated by "--end note]" at the end of the 2nd sentence appropriately. | | | | |  | | | | |
| US 024 | | |  | | 05.02 [lex.phases] | | | | 2 | | | | Te | | | | Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. This concerns the following text: "Except for splices reverted in a raw string literal, if a splice results in a character sequence that matches the syntax of a universal-character-name, the behavior is undefined." The behavior should either be well defined, and produce the universal character; ill-formed (diagnostic required); or conditionally supported as forming the universal character (must be documented and diagnosed if not supported) | | | | Make this either a diagnosable error or produce the universal character (potentially as conditionally supported behaviour). | | | | |  | | | | |
| US 025 | | |  | | 05.02 [lex.phases] | | | | 4 | | | | Te | | | | Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. This concerns the following text: "If a character sequence that matches the syntax of a universal-character-name is produced by token concatenation (15.5.3), the behavior is undefined." The behavior should either be well defined, and produce the universal character; ill-formed (diagnostic required); or conditionally supported as forming the universal character (must be documented and diagnosed if not supported) | | | | Make this either a diagnosable error or produce the universal character (potentially as conditionally supported behaviour). | | | | |  | | | | |
| US 026 | | |  | | 05.04  [lex.pptoken],  6.5  [basic.link],  10  [module.unit],  10.4  [module.global],  15 [cpp],  15.4 [cpp.glob.frag] | | | | See [P1857r0](https://wg21.link/p1850r0) | | | | te | | | | import is a directive while module is not. This makes it difficult to do dependency scanning via partial preprocessing, which was the design intent of making import a directive. | | | | Make module a directive. See [P1857r0](https://wg21.link/p1850r0) for details. | | | | |  | | | | |
| US 027 | | |  | | 05.04 [lex.pptoken] | | | | 4 | | | | Te | | | | It is undefined behavior to have an unmatched ' or " character as a single-character token (i.e., surround by whitespace). Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. | | | | Make this a diagnosable error. | | | | |  | | | | |
| US 028 | | |  | | 05.06 [lex.token] | | | |  | | | | Te | | | | A token is defined as either an identifier, a keyword, a literal, an operator, or a punctuator. While the first 4 all have definitions for their grammar term, there is no definition of a punctuator. | | | | Add the missing definition for the grammar term punctuator. It probably should be "one of:" the set of tokens in preprocessing-op-or-punc that are not also present in the grammar production for operator | | | | |  | | | | |
| NL 029 | | | 1 | | 05.10 [lex.name] | | | | [tab:lex.name.allowed] | | | | te | | | | Allowed characters include those from U+200b until U+206x; these are zero-width and control characters that lead to impossible to type names, indistinguishable names and unusable code & compile errors (such as those accidentally including RTL modifiers). | | | | Disallow invisible characters in this range | | | | |  | | | | |
| JP2 030 | | |  | | 05.13.04 | | | |  | | | | ed | | | | At binary-exponent-part, a small p and a large P are mixed and hard to distinguish with current font. | | | | Use another font which has distinct typefaces for p and P. | | | | |  | | | | |
| US 031 | | |  | | 06.02  [basic.def.  odr] | | | | 09.2.2 | | | | Ed | | | | The last line of the example is (mildly) confusing:  { [n]{ return n; }; }; *// OK*  Yes ‘n’ is reachable but the code will not compile. | | | | I would suggest changing the line to:  { [n]{ return n != 0; }; }; *// OK* | | | | |  | | | | |
| GB 032 | | |  | | 06.04.2 | | | |  | | | | Ed | | | | [basic.lookup.argdep] example paragraph 5 lacks necessary namespace qualification  In section 6.4.2 paragraph 5:  Translation unit 1:  namespace R {  export struct X {};  export void f(X);  }  namespace S {  export void f(X, X);  }  the f() declaration in S in translation unit 1 lacks namespace qualification for X. | | | | Insert R:: as shown:  namespace R {  export struct X {};  export void f(X);  }  namespace S {  export void f(<ins>R::</ins>X, <ins>R::</ins>X);  } | | | | |  | | | | |
| US 033 | | |  | | 06.05 | | | |  | | | | te | | | | Import declarations may only appear at the top level, outside of any bracing. This means they cannot appear inside a linkage block (9.10) ‘extern “C” { … }’. Header translation (15.2/7) is specified as a rewrite of a #include directive into an import declaration. Unfortunately header files exist that contain (different sets of) #includes (of C headers) inside linkage blocks – even when the included header is itself C++-aware and encapsulates its contents in a linkage block. Such included headers may be included elsewhere from the top level and be suitable for include translation. It would be unfortunate if the decision as to whether include translation should be performed had to be made context-sensitive, examining previously lexed tokens. It could also be surprising to users. Thus, in practice, implementors must permit import declarations inside linkage blocks (at least for these translated cases). | | | | Permit import-declarations within linkage blocks at the global namespace level. There is a commonality with the next issue, which is noted as an editorial issue below.FYI both Clang’s module extension and GCC’s C++ Modules found it necessary to implement this extension, due to glibc header files. | | | | |  | | | | |
| US 034 | | |  | | 06.05 | | | |  | | | | te | | | | Import declarations may only appear at at the top level, outside of any bracing. This means they cannot appear inside an export block (10.2) ‘export { … }’. Such non-generality of export blocks can lead to user confusion. The original intent was to simplify module preamble parsing, along with the backtracking the original ATOM design envisioned. In the [p1103](https://wg21.link/p1103) specification it simplifies the recognition of header imports during phase 4, where their exported macros must be made available. However, as specified in the previous issue, implementors must account for some cases occurring within linkage blocks. Permitting imports within export blocks would not further inconvenience implementations. Note, relaxing this restriction does not break the status quo with lexing import declarations as preprocessing directives. The export semantics are applied at the c++-parser level. | | | | Permit import-declarations within export blocks at the global namespace level. There is a commonality with the next issue, which is noted as an editorial issue below. In a named module, an export block of imports would be restricted to the beginning of module purview, and perhaps only permitted to contain import declarations, to retain the current requirement of placing import declarations immediately after the module declaration. | | | | |  | | | | |
| US 035 | | |  | | 06.05 | | | |  | | | | te | | | | The CD requires implementation linkage promotion for internal-linkage entities that are referenced from locations made visible to other TUs. For example:  export module foo;  static void internal () {}  export template<typename T> void widget ()  { internal (); }  instantiations of ‘widget’ will need to reference ‘internal’, thus its symbol must be globally visible. This presents implementation difficulties and inhibits certain optimizations. Papers [p1347](https://wg21.link/p1347) and [p1395](https://wg21.link/p1395) describe the issue in some detail. | | | | [P1498](https://wg21.link/p1498) proposes making such references ill-formed. Wording should be completed to specify that. A decision should be made as to whether the error is indicated when compiling the module interface that exposes the internal entity, or whether referencing that internal entity from a different TU (instantiating a template, calling an inline function, etc). | | | | |  | | | | |
| US 036 | | |  | | 06.05 | | | |  | | | | te | | | | The semantics of the private module fragment (PMF) are underspecified. It appears well formed to declare an entity in the interface purview and define it in the PMF. How does this interact with inlining, instantiation and internal linkage definitions? The intent of the PMF is as-if it is a separate module implementation unit, but we do not define the boundary between the interface purview and the PMF as a translation unit boundary. Implementations may defer instantiation (of function definitions) to the end of translation. Similarly internal linkage and inline functions (that are ODR-used) must be defined at the end of translation. Perhaps emission of Compiled Module Interface (CMI) should be deferred to the end of translation – and not at the beginning of the PMF. It will therefore observe entities declared or defined in the PMF. This may create implementation difficulty to preserve the semantics of the PMF being as-if a separate translation unit. Alternatively, the boundary could be specified as a new kind of ‘end of translation’ point. This seems a dramatic change. | | | | Remove the private module fragment. It’s semantics are too ill-defined, and I do not believe there is sufficient experience to define them at this point. | | | | |  | | | | |
| US 037 | | |  | | 06.05 & 10.2 | | | |  | | | | ed | | | | Linkage and export blocks are the only two namespace-scope braced constructs that do not introduce a scope. As described above, implementations must at least permit some imports within linkage blocks. This leads to the confusing situation that imports are neither a top-level construct, nor one that is permitted at any global scope. If both the above restrictions are relaxed, import-declarations can be described simply as entities appearing at global scope. | | | | Accept the previous two changes. | | | | |  | | | | |
| CA 038 | | |  | | 06.05 [basic.link], 9.8 [namespace.udecl], 9.10 [dcl.link], 12.1 [over.load] | | | |  | | | | te | | | | Trailing *requires-clause*s are ignored in various places covering the identity of functions. | | | | Update appearances of “parameter-type-list” in [basic.link], [namespace.udecl], [dcl.link], and [over.load] to also take *require-clause*s into account. | | | | |  | | | | |
| FR 039 | | |  | | 06.05.2/4 | | | |  | | | | **te** | | | | The current wording allows a non-exported function to be visible to an importing translation unit via ADL. That is too broad and goes against the abstraction boundaries that modules are supposed to bring. | | | | Specify that these names are available only during the second phase of ADL during template instantiation and only if the module interface unit is on the path of instantiation. | | | | |  | | | | | |
| US 040 | | |  | | 06.06.2  [intro.object] | | | | n/a | | | | te | | | | The core language changes from [P0593R5](https://wg21.link/p0593r5) to enable implicit object creation are important to support P1004R2. These changes should be viewed as resolving a defect in C++. | | | | Apply the changes from [P0593R5](https://wg21.link/p0593r5), other than the addition of new standard library functions, and treat them as the resolution of a defect report. | | | | |  | | | | |
| US 041 | | |  | | 06.06.3 [basic.life] | | | | 8 | | | | te | | | | The subject paragraph is missing wording which disallows using a glvalue which referred to a member subobject from being used to refer to an object that is not similarly such a member subobject. This means that a subobject of one class might become a subobject of an altogether different class.  For example, with the wording as it is, the following should return typeid(Z) if the assertions pass:  #include <new> #include <memory> #include <cassert> #include <typeinfo>  struct X { int x; }; struct Y : X {  virtual ~Y() = default; }; struct Z : Y { };  const std::type\_info &f() {  Y \*yp = new Y;  int &yx = yp->x;  assert(sizeof(Y) == sizeof(Z) &&  alignof(Y) == alignof(Z));  Z \*zp = new ((void \*)yp) Z;  assert(std::addressof(yx) ==  std::addressof(zp->x));  return typeid(\*static\_cast<Y \*>(  reinterpret\_cast<X \*>(&yx))); }  It is surprising that the set of types associated with the most derived objects of the objects that are pointer interconvertible with the pointee of a pointer value may change between points at which the pointer value may be plainly dereferenced.  In this specific case, yx was obtained as a reference to an int subobject of an object whose most derived type is Y. Yet yx came to refer to a subobject that was not of any object whose most derived type is Y. yx should not be allowed to change identities in this manner.  This appears to be a defect in all prior versions of C++. | | | | Add a new bullets to 6.6.3 [basic.life] paragraph 8:   * the respective complete objects (6.6.2 [intro.object]) of the original object and the new object are either the same or meet these requirements in turn, and * the original object and the new object are each the same subobject in relation to their respective complete objects as each other. | | | | |  | | | | |
| US 042 | | |  | | 06.06.3 [basic.life] | | | | 8 | | | | te | | | | The restriction on automatically referring to the new object when a class type contains a non-static data member of const-qualified or reference type renders use of the pointer returned by data() of a std::vector of such class types unsafe without applying std::launder. In particular, if a pointer returned by data() on a one-element vector of a class type that is subject to the restriction is stored and then followed by a pop\_back() and a call to push\_back, then the stored pointer would not be usable to access the new element without applying launder. The need to apply launder is not currently indicated by the library wording and is presumably unwanted. | | | | In the bullet, strike:   * the type of the original object is not const-qualified~~, and, if a class type, does not contain any non-static data member whose type is const-qualified or a reference type~~, and   Remove, from subclause 6.6.2 [intro.object] paragraph 2, the note beginning with “[i]f the subobject contains a reference member or [ … ]” and the associated example. | | | | |  | | | | |
| US 043 | | |  | | 06.08  [basic.def]  9  [dcl.dcl]  9.1.8.3  [dcl.type.  elab]  9.6.2  [enum.udecl]  9.8  [namespace.udecl]  11.3  [class.mem]  15.10  [cpp.predefined] | | | | 02.17  1                  Table 17 | | | | te | | | | The syntax of the using enum feature is very subtle, which may be confusing to users. Consider, for example:  struct B {  protected:  enum class E { e1, e2 };  };  class D1 : public B {  public:  using B::E;  // Creates a public member using  // declaration in 'D1' for 'B::E'.  };  class D2 : public B {  public:  using enum B::E;  // Creates a public member using  // declarations in 'D2' for  // 'E::e1' and 'E::e2'.  };  void f() {  D1::E v1; // Ok  D2::E v2; // COMPILE FAILURE  D1::e1; // COMPILE FAILURE  D2::e2; // Ok  }  Additionally, the syntax looks similar to elaborated type specifiers, which may also lead to confusion. Please reconsider this feature; it was added very late in the C++20 design process, and it may be better to move it to C++23 to give us more time to work on the design. | | | | Possible resolutions include:   * Disallow class member using enum declarations. * Move using enum from C++20 to C++23 to give us additional time to work on it. | | | | |  | | | | |
| CZ 044 | | |  | | 07 | | | | 07.7 | | | | te | | | | Small Buffer Optimization is a common and widespread container implementation strategy for containers that do not have as strong a guarantee for iterator invalidation as std::vector, such as std::string. With both vector and string being made constexpr, Small Buffer Optimization cannot be done legally at constexpr time due to the inability to invoke std::construct\_at or std::destroy\_at on memory that has automatic storage duration. That is, the wording does not seem to bless the following program as legal:  #include <memory>  constexpr int foo () {  int x = 0;  std::destroy\_at(&x);  std::construct\_at(&x, 2);  return x;  }  constexpr int bar () {  int x[2]{};  std::destroy\_at(&x[0]);  std::construct\_at(&x[0], 3);  return x[0];  }  int main () {  static\_assert(foo() == 2);  static\_assert(bar() == 3);  }  This presents problems for a future wherein someone wishes to work with SBO containers, and may present a serious problem for the future work with getting constexpr-capable containers to escape compile-time and serialize into something that can be used at runtime. | | | | Change [expr.const/5.1-2] to read as follows:  —  for a call to std::construct\_­at or std::ranges::construct\_­at, the first argument, of type T\*, does not point to either storage allocated with std::allocator<T> or an automatic storage duration variable of type T, or the evaluation of the underlying constructor call is not a core constant expression, or  —  for a call to std::destroy\_­at or std::ranges::destroy\_­at, the first argument, of type T\*, does not point to either storage allocated with std::allocator<T> or an automatic storage duration variable of type T, or the evaluation of the underlying destructor call is not a core constant expression. | | | | |  | | | | |
| JP3 045 | | |  | | 07.02.2 | | | | p3.6 | | | | ed | | | | "C1 is reference-related to C1" looks like a typo. | | | | C1 is reference-related to C2 | | | | |  | | | | |
| GB 046 | | |  | | 07.05.4 | | | |  | | | | Te | | | | Implementations should be allowed to cache the results of concept specialisation evaluations  The current wording of [expr.prim.id]/4 precludes the possibility of caching concept specialisations (which has led to demonstrable performance issues in implementations that are standard-conforming contrasted against implementations that cache results). | | | | Add wording to the effect of:  A program which causes a given set of template arguments to change their satisfaction of a constraint, and relies upon this change, is ill-formed, no diagnostic is required. | | | | |  | | | | |
| US 047 | | |  | | 07.05.4.1 | | | | 2 | | | | ed | | | | If the entity is a template parameter object for a template parameter of type T (13.1), the type of the expression is const T." follows from  [temp.param]/1.6 (which establishes that the template parameter object has that type). | | | | Strike the sentence. | | | | |  | | | | |
| BG7 048 | | | P 117-119 | | 07.06.2.3 | | | |  | | | | ge | | | | The same comments as 6, this time for the Awaitable concept. await\_ready, await\_suspend, await\_resume are too uniformly named, with no differentiation b/w the observer await\_ready and the other functions. | | | | Rename await\_suspend to await\_on\_suspend and await\_resume to await\_on\_resume. | | | | |  | | | | |
| BG2 049 | | | P 118 | | 07.06.2.3 | | | | (3.7) | | | | ge | | | | The variants of await\_suspend() that return void and bool can introduce unexpected stack frames. Because of this coroutine authors cannot rely that co\_await will not introduce stack frames. | | | | Leave only the await\_suspend() variant that returns a coroutine\_handle. | | | | |  | | | | |
| BG1 050 | | | P 117-119, P 134-135, P 150, P 204-206, P 520-524, etc. | | 07.06.2.3, 7.6.17, 8.6.4, 9.4.4, 17.12, etc. | | | |  | | | | ge | | | | The coroutines design that was adopted has serious flaws:   * The language-provided constructs cannot be used without helper types (“coroutine types” and “awaiters”). It is expected that these helper types will be provided by library writers. We believe that major and subtle differences will often make it hard or impossible to compose helper classes written by different authors. * The semantics of the co\_await operator enforces a “trampoline” style transfer of control and hinders the adoption of a tail-call-style transfer of control in the future. * The coroutine\_handle::resume() method is untyped which makes transfer of control type-unsafe. * Furthermore, the untyped resume operation separates the transfer of control from the transfer of data. This affects performance, jeopardizes safety and increases the risk of bugs. * There can be multiple copies of the same coroutine\_handle raising the risk of inadvertent incorrect resumption of a coroutine. * The asymmetric variants of await\_suspend() cause calls between coroutines to accumulate stack frames. This is contrary to the very definition of coroutines and precludes certain synchronous usage patterns. * The symmetric variant of await\_suspend() is forced to use a trampoline which presents a challenge to the optimizer. * Having both the symmetric and asymmetric variants of await\_suspend() means that coroutine authors cannot rely that using co\_await will prevent unbounded accumulation of stack frames. * The type-erased coroutine frame introduces allocations that are out of the control of the user. The proposed solution (to declare an unimplemented override of new for the coroutine frame that causes compilation errors when the compiler fails to elide the allocations) ensures only the detection of these allocations but doesn’t give a recipe for eliminating them.   Additionally, after the TS was merged into the WD, early adopters identified problems (in P1662R and P1745R0) in the design related to cancellation and cleanup which the Bulgarian NB finds valid. The authors have indicated that eliminating these issues will require breaking changes to the current design. We strongly support resolving these problems without breaking changes.  More importantly, these findings suggest that although the TS has been used in production, the usage experience has been limited to just a few specific use cases.  Last but not least, at least three promising alternatives have been presented in the last two years - P1063, P1430 and P1745 - and for the most part these alternative designs cannot be reached from the current design without breaking changes. The Bulgarian NB is seriously worried that the work in the design space is not mature and that important design directions have not been considered. | | | | Delay coroutines from C++ 20 to C++ 23 and work actively to reach a qualitatively better design and substantial usage experience beyond the current limited set of users and use cases. | | | | |  | | | | |
| GB 051 | | |  | | 08.05.4 | | | |  | | | | Te | | | | Range-based for loops should look for ranges::begin and ranges::end  ranges::begin and ranges::end perform similar work to the range-based for loop. It would be good to have these as a bullet in the lookup for what a range-for looks for to help with the deprecation process of std::begin and std::end in a future standard. | | | | Add a sub-item that allows begin-expr to be determined from a call to ranges::begin(range) and end-expr to be determined from a call to ranges::end(range).  Consider adding this sub-item between 1.3.1 and 1.3.2. | | | | |  | | | | |
| US 052 | | |  | | 08.06.4 | | | | 1 | | | | ed | | | | ‘A coroutine shall not return to its caller or resumer by a return statement (8.6.3)’ suggests this is a dynamic restriction – i.e. a coroutine may contain return-statements, provided they are not reached during execution. This is believed a wording error. | | | | Replace with ‘A coroutine shall not contain a return statement’. I believe it clear that ‘contain’ already includes unreachable statements of constexpr-ifs. Thus one cannot use constexpr-ifs to define a function that is conditionally a coroutine. | | | | |  | | | | |
| US 053 | | |  | | 08.06.4 [stmt.return.coroutine] | | | | 3 | | | | TE | | | | Apply remaining coroutine TS issues to the working paper. | | | | Apply resolution of the issue 34 “Mandate the return type for return\_void and return\_value to be void.” From [p0664r8](https://wg21.link/p0664r8) | | | | |  | | | | |
| GB 054 | | |  | | 09.01.5 | | | |  | | | | Te | | | | P1331R2 intended to allow uninitialized state in constant evaluation, but left behind two bullets in the requirements for a constexpr function that still require explicit initialization for variant members. | | | | Strike [dcl.constexpr]/4.1 and /4.2:  — if the class is a union having variant members (11.4), exactly one of them shall be initialized;  — if the class is a union-like class, but is not a union, for each of its anonymous union members having variant members, exactly one of them shall be initialized;  (This is CWG2424) | | | | |  | | | | |
| US 055 | | |  | | 09.02.3.5  [dcl.fct] | | | | 18 | | | | te | | | | Defaulting an argument with a placeholder-type-specifier in an abbreviated function template is valid but useless. Calling the function without providing values for the default parameters will result in a deduction error.  void foo(std::integral auto i = 0) {}  foo(); // COMPILE FAILURE foo(1); // Ok  This is also a problem for generic lambdas:  auto foo = [] (auto i = 0) {};  foo(); // COMPILE FAILURE  foo(1); // Ok | | | | This problem could potentially be resolved if the invented template parameter was assigned an appropriate default value, e.g. decltype() of the provided default argument, although this could not work if the default argument references other parameters in the function. | | | | |  | | | | |
| US 056 | | |  | | 09.03  [dcl.init] | | | | n/a | | | | te | | | | C++20 intends to support parenthesized initialization of aggregates (eg, StructWithThreeElements(1, 2, 3)). The support for this also permits dropping trailing elements (eg, StructWithThreeElements(1)). Function-style casts from a single element are by definition equivalent to C-style casts, so this also unintentionally and problematically permits (StructWithThreeElements)1 and static\_cast<StructWithThreeElements>(1).  In addition, while the problem this feature is trying to address is real (container emplace is not sufficiently general), this change only tackles a corner of the problem (only aggregates), and in so doing, further increases the complexity of C++ initialization. This problem should be addressed by a big-picture consideration of C++ initialization; something like lazy parameters are probably the right way to provide clean simple syntax for a generalized emplace. | | | | Revert the application of [P0960R3](https://wg21.link/p0960r3).  Even without the more general concerns about this feature, we do not have any confidence that a different change to initialization, designed at the last minute, will do any better than [P0960](https://wg21.link/p1960) did, so we should not attempt to repair it, and should instead reconsider ways to tackle the problem for C++23. | | | | |  | | | | |
| JP4 057 | | |  | | 09.03.4 | | | | p3.10 | | | | ed | | | | "const B &b2{a};" in the example is in inconsistent style on the place of &. | | | | const B& b2{a}; | | | | |  | | | | |
| US 058 | | |  | | 09.04.1 | | | |  | | | | ge | | | | It is surprising (if consistent with if constexpr) that the bodies of non-templated functions whose constraints are not satisfied must be valid. | | | | Make such function bodies discarded statements; change [stmt.if]/2 to treat discarded statements as non-instantiated templated entities that are not subject to [temp.res]/8. | | | | |  | | | | |
| BG5 059 | | | P 204-206 | | 09.04.4 | | | |  | | | | ge | | | | There is a missed opportunity for consistency in naming get\_return\_object, initial\_suspend and final\_suspend.  These 3 functions are very similar in the way, their return value is actually the customization point (and not so much the implementation of the functions themselves). | | | | Rename initial\_suspend to get\_initial\_suspend and final\_suspend to get\_final\_suspend.  Note: Is not just about consistency. initial\_suspend and final\_suspend can be very confusing to newcomers as they are novel customization concepts, that are worded similarly to regular observer functions.  By having a get\_ prefix - which observers in the Standard Library do not use - people will be reminded of get\_return\_object and the need to return a user-provided object as a customization. | | | | |  | | | | |
| BG6 060 | | | P 204-206 | | 09.04.4 | | | |  | | | | ge | | | | Most of the Promise function names are worded as observer functions, yet none of them are!  Example: unhandled\_exception reads exactly the same as uncaught\_exception(s), already in the standard library!  This will create confusion in people learning the API, as the same naming style is reused with a different meaning. | | | | Mark functions, which are user implementations of required behavior, with an "on\_" prefix. unhanded\_exception becomes on\_unhanded\_exception, return\_value becomes, on\_return\_value, return\_void becomes on\_return\_void, yield\_value becomes on\_yield\_value.  This naming might seem novel, but the API itself is quite different from the rest of the Standard Library already. The fact that, for example, yield\_value does not actually return the yielded value, the way std::begin does, is already a stark enough contrast that calls for different naming.  Note: await\_transform is a special, optional case and should be let as it is, unless the Committee decides otherwise. | | | | |  | | | | |
| US 061 | | |  | | 09.04.4 | | | | 10 | | | | te | | | | Coroutine allocation does not consider std::align\_val\_t overloads introduced in C++17 | | | | Add them to the sequence of operator new calls that are attempted using wording similar to 7.6.2.7/18 | | | | |  | | | | |
| US 062 | | |  | | 09.04.4  [dcl.fct.def.coroutine] | | | | 5 | | | | te | | | | Coroutine promise types are currently required to define an unhandled\_exception member function. For many simple promise types, the definition of this function is simple and trivial; it just rethrows. It is unfortunate that authors of promise types are required to write this boilerplate. Many other coroutine extension points are only used if they exist, but this one is required. It would be nice if it was not required for noexcept coroutines. | | | | If the coroutine is noexcept, do not wrap the invocation of the function body in a try/catch block and do not require that promise types define an unhandled\_exception method. | | | | |  | | | | |
| US 063 | | |  | | 09.04.4 | | | | 9 | | | | te | | | | The construction of the argument list for the call to the allocation function to allocate the 'coroutine state' does not call the overload of  operator new() that accepts a std::align\_val\_t in the case that the allocation required for the coroutine has 'new-extended alignment'. This means that allocations of coroutine frames may not be correctly aligned in cases where the coroutine state contains overaligned types. | | | | Apply similar wording from [expr.new]p18: Insert "If the coroutine state has new-extended alignment then the next argument is std::align\_val\_t." after "has type size\_t."  Insert at end of paragraph: If no matching function is found and the allocated coroutine state has new-extended alignment, the alignment argument is removed from the argument list, and overload resolution is performed again. | | | | |  | | | | |
| US 064 | | |  | | 09.04.4 [dcl.fct.def.  coroutine]. | | | | 11 | | | | TE | | | | Apply remaining coroutine TS issues to the working paper. | | | | Apply resolution of the issue 33. “Parameter copy wording does not capture the intent.” From [p0664r8](https://wg21.link/p0664r8) | | | | |  | | | | |
| US 065 | | |  | | 09.04.4 [dcl.fct.def.  coroutine] | | | | 3 | | | | TE | | | | Apply remaining coroutine TS issues to the working paper. | | | | Apply resolution of the issue 24 “The specification of initial suspend point does not correctly captures the intent.” From [p0664r8](https://wg21.link/p0664r8) | | | | |  | | | | |
| FR 066 | | |  | | 09.05.4 | | | |  | | | | **te** | | | | Defining a coroutine’s promise type forces people who do not care about exceptions to write more boilerplate than necessary. | | | | std::coroutine\_traits<....>::promise\_type::unhandled\_exception should be optional and equivalent to exception re-throw when not declared. | | | | |  | | | | | |
| FR 067 | | |  | | 09.05.4 | | | |  | | | | **te** | | | | The number of customization points used by coroutines can be reduced while reducing boilerplate, increasing safety and consistency. | | | | Adopt P1477 and remove coroutine\_handle::from\_promise (which is not needed after adopting P1477) | | | | |  | | | | | |
| FR 068 | | |  | | 09.05.4 | | | |  | | | | **te** | | | | promise\_type::final\_suspend can be renamed to promise\_type::done to be more consistent with the proposed executor design. | | | | Rename the customization point promise\_type::final\_suspend to done | | | | |  | | | | | |
| US 069 | | |  | | 09.06 | | | |  | | | | te | | | | Unnamed unscoped enumerations may be defined in multiple header units, and have no linkage. The merging rules of multiple definitions from header units or appearing textually outside of module purview require implementations to determine if two particular definitions are for the same entity. There is no mechanism specified to determine whether two such enumerations are for the same entity. Unnamed, *untypedefed*, enums are common in header files, as the enumeration values also appear in the containing (namespace) scope. A mechanism should be specified. | | | | Use the first enumerator as the key. If two unnamed unscoped enumeration definitions in the same scope have the same identifier for their first enumerator, they are defining the same enumerated type. (It therefore is an ODR violation if the enumerators are not the same.)FYI this is the heuristic independently implemented in the Clang Modules extension and GCC C++ Modules. It is expected in Clang C++ Modules. | | | | |  | | | | |
| US 070 | | |  | | 09.06.2  [enum.  udecl] | | | |  | | | | te | | | | Keep the “using enum” language feature in C++20. This is a very useful feature. | | | | If it helps consensus, consider the alternative spelling “using enumerators E;” instead. | | | | |  | | | | |
| US 071 | | |  | | 09.09.4 [dcl.fct.def.coroutine] | | | | 1 | | | | te | | | | Coroutine keywords are prefixed with `co\_` which looks ill-designed. Papers to rename them weren't seriously considered. | | | | Adopt [P1485R1](https://wg21.link/p1485r1). | | | | |  | | | | |
| US 072 | | |  | | 1 - 15  Language | | | |  | | | | ge | | | | Please address [open CWG issues](http://www.open-std.org/jtc1/sc22/wg21/docs/cwg_active.html). | | | | Appropriate action would include making changes to the CD, identifying an issue as not requiring a change to the CD, or deferring the issue to a later point in time. | | | | |  | | | | |
| US 073 | | |  | | 10  and others | | | |  | | | | te | | | | Because of the number of serious bugs that have been discovered in modules, including some that are requiring design changes at this late stage, modules do not seem sufficiently ready to be included in C++20. | | | | Remove modules from the working draft with the intention to add them back soon after C++20 ships. | | | | |  | | | | |
| US 074 | | |  | | 10 [module.unit] | | | |  | | | | te | | | | Module names can contain ., but they have no semantic meaning. This seems to be a holdover from when they did have semantic meaning in previous designs. Keeping them will cause user confusion and may prevent future addition of sematnics. | | | | Disallow . in module-names. See [P1873r0](https://wg21.link/p1873r0) for details. | | | | |  | | | | |
| FR 075 | | |  | | 10.01 | | | |  | | | | **te** | | | | As specified, modules names can contain spaces and comments  module foo /\* \*/.bar; | | | | Specify that modules names are formed of a single entity with different rules than other identifiers. Consider not forbidding modules names containing C++ keywords between dots - but keep the ability to separate by dots so that modules can be organized logically. | | | | |  | | | | | |
| US 076 | | |  | | 10.01  [module.  unit] | | | |  | | | | te | | | | Remove *module-name-qualifier*. It removes our ability to have proper submodules, if we choose to do so at a later date. | | | | Remove *module-name-qualifier*. | | | | |  | | | | |
| US 077 | | |  | | 10.01 | | | |  | | | | te | | | | Allowing semanticless . in module names closes off the possibility of providing semantics in the future. Additionally . is a poor choice of separator in C++, where we currently use :: to represent the semanticless hierarchy used for code organization. | | | | Disallow . in module names by removing the module-name-qualifier from the module-name grammar. Alternatively, replace the . in module-name-qualifier with :: | | | | |  | | | | |
| GB 078 | | |  | | 10.01 | | | |  | | | | Ed | | | | Ensure consistent use of "digits"  'std + digits' being reserved is referenced in 3 places (wrt. modules and namespaces): [namespace.future] 16.5.4.2.3, [diff.cpp14.library] C.4.7, [module.unit] 10.1.  In [module] "digits" appears in italics (indicating a definition by [intro.defs], though italic "digits" is also used in [cpp.include] 15.2).  In [namespace] and [diff] "digits" does not appear in italics.  Other references in the standard qualify "digits" with a base (e.g. "decimal digits"), or use it in a radix-neutral context (e.g. "number of digits").  Presumably the intention is that "digit" refers to the grammar in [lex.name] 5.10, e.g. that "stda" is not intended to be a reserved module name or namespace name and "std8" is.  These references should be consistent. | | | | Either of:  1) Remove italic formatting from "digits" in [module.unit] 10.1; or  2) Apply italic formatting to "digits" in [namespace.future] 16.5.4.2.3 and [diff.cpp14.library] C.4.7  and optionally qualify "digits" with "decimal" (or other base as appropriate). | | | | |  | | | | |
| GB 079 | | |  | | 10.01 | | | |  | | | | Ed | | | | The private-module-fragment is mentioned many times in the standard but its usage its obscure; as is in the papers. Nobody in the BSI knows its intended usage.  Please provide an example. | | | | Example added | | | | |  | | | | |
| PL 080 | | |  | | 10.01 [module.unit] | | | |  | | | | te | | | | Dots in module names suggest a hierarchy, and that introduces a possibility of confusion among users. Several members of the committee have already suggested they intend to implement such a perceived hierarchy in the modules they ship, reexporting all "submodules" from the "main" module, and this will make it hard to invent other semantics for "submodules" in the future, and will also cause users to think that those semantics are a part of the language.  We should allow ourselves more time to consider the possible semantics for actual hierarchical modules and submodules. Allowing dots in the module names makes it impossible for us to introduce such semantics. If we decide that we do not want to introduce hierarchical module semantics in the future, we can always relax the rules in the standard and allow dots in module names in the future. | | | | Adopt P1873. | | | | |  | | | | | |
| US 081 | | |  | | 10.02 | | | | 3 | | | | ge | | | | Lifting the restriction on enclosing sequences of existing declarations that might include (say) a static\_assert or might be rendered empty by the preprocessor, with "export {...}" in C++23, as suggested by SG2, would not help after existing code has been updated to use modules. | | | | Apply the change recommended by §3 of [P1766R1](https://wg21.link/p1766r1). | | | | |  | | | | |
| US 082 | | |  | | 10.03  [module.  import] | | | |  | | | | te | | | | The order of initialization of globals in interface and header units is not defined for imported modules. This has the potential to break code, including iostream, when moving to modules. | | | | Define an ordering. See [P1874r0](https://wg21.link/p1874) for details. | | | | |  | | | | |
| US 083 | | |  | | 10.03 | | | |  | | | | te | | | | Header units make macros available to their importers. A header unit, A, may export an imported header unit, B – indeed, they cannot not export. An importer of A will observe B’s exported macro set.   However, if a named module, C, reexport imported header unit, B, C’s importers will not observe B’s macros via that path.   This is confusing, and may make B’s API unusable, without indirect importers directly importing B.  This partial indirect importation has similar issues to today’s unplanned reliance of indirect include files. | | | | Disallow exportation of header unit imports from named modules. Importers of a named module that relies on a header unit in its user interface will need to explicitly import that header unit. | | | | |  | | | | |
| US 084 | | |  | | 10.03 | | | |  | | | | te | | | | export import foo; should not be allowed outside of module interface purview. | | | | Ban using similar wording as 10.2/1 | | | | |  | | | | |
| US 085 | | |  | | 10.03 | | | | 6 | | | | ed | | | | The paragraph discusses import declarations importing Translation Units. Translation Units are not imported, modules are (para 3). Para 5, describing header units, gets close to an import of a translation unit. The paragraph is describing the transitive nature of exported import declarations. The wording is confusing. | | | | Reword paragraph 6 to avoid equating ‘module’ and ‘translation unit’. | | | | |  | | | | |
| US 086 | | |  | | 10.03  10.6 | | | | 6  1 | | | | te | | | | When a translation unit U imports translation unit T, the wording in 10.3/6 treats non-exported imports in T differently depending on whether T and U are in the same module. If they are in the same module those imports are also imported by U. This inconsistency is presumably to allow U to implicitly import implementation partitions imported by T, which cannot be exported. However this also applies to unrelated imports in T, including those within the global module fragment caused by implicit translation of #includes. This has an effect on the reachability rules in 10.6/1 | | | | A few options, some of which can be sensibly combined:   1. Apply this only to imports within the module purview of T 2. Apply this only for imports of partitions of the same module as T. Possibly only implementation partitions, because interface partitions could be directly exported 3. Eliminate this implicit import rule, and only implicitly import explicitly exported imports. 4. Eliminate the distinction between interface and implementation partitions. Allow any partition unit to be exported. Only require partition units containing export declarations to be exported by the primary interface unit. | | | | |  | | | | |
| US 087 | | |  | | 10.03 | | | | 9 | | | | te | | | | The wording restricts interface dependencies to module units. This allows header units to have cyclic dependencies and even allows them to be used to break dependencies among module units. | | | | Replace all usages of “module unit” in 10.3/9 with “translation unit” | | | | |  | | | | |
| US 088 | | |  | | 10.04 [module.  global],  15.4 [cpp.glob.  frag] | | | |  | | | | ed | | | | [module.global] and [cpp.glob.frag] refer to the same thing, but use different names. | | | | Use the same name. Either glob.frag or global. | | | | |  | | | | |
| GB 089 | | |  | | 10.06 | | | |  | | | | Ed | | | | Mark translation units in [module.reach] example  [module.reach] 10.6 p5  The example in this paragraph comprises two translation units. They are currently separated by a blank line. These should be marked up as being separate translation units as per other examples, such as that in p4. | | | | Add plain text line "Translation unit #1:" before "export module A;"  Add plain text line "Translation unit #2:" before "module B;" | | | | |  | | | | |
| US 090 | | |  | | 11.03.1 | | | | 1 | | | | te | | | | If [P1815](https://wg21.link/p1875) is applied to C++20 (as is expected), but [P1779](https://wg21.link/p1779) is applied to C++23 (as suggested by SG2), the latter will be a silent performance regression from changing language versions rather than being a concern to be addressed while converting header files to modules. | | | | Apply [P1779](https://wg21.link/p1779) (after appropriate CWG review). | | | | |  | | | | |
| US 091 | | |  | | 11.10.01 | | | | 04.1 | | | | Te | | | | A user-defined operator<=> can affect whether an enumeration type has strong structural equality. | | | | Specify that all enumeration types have that property. | | | | |  | | | | |
| US 092 | | |  | | 11.10.01 | | | | 04.2.1  [class.  compare.  default] | | | | te | | | | Class types with a non-static data member of array type do not have strong structural equality. | | | | Require only that the (ultimate) element type of such an array type have strong structural equality. | | | | |  | | | | |
| US 093 | | |  | | 11.10.01 | | | | 4 | | | | te | | | | The definition of strong structural equality is used only in [temp.param]. | | | | Move it to [temp.param] | | | | |  | | | | |
| US 094 | | |  | | 11.10.01 | | | | 4 | | | | te | | | | Class types can have strong structural equality even if their operator== is deleted (since overload resolution "succeeds" even if it finds a deleted function). | | | | Say "that is defaulted in the definition of C and is not defined as deleted". | | | | |  | | | | |
| US 095 | | |  | | 12.02 | | | | 1 | | | | ge | | | | A declaration redeclares a constrained function if its requires-clause is equivalent. No atomic constraint expression can be equivalent to any other unless it is accessed via a concept (even within a single translation unit); for functions with no template parameters, all requires-clauses are functionally equivalent to either "requires true" or "requires false". In either case, the program is ill-formed NDR because of constructs that are functionally equivalent but not equivalent. | | | | Document these severe restrictions, change the definition of (functionally) equivalent for atomic constraints to rely on the ODR, and/or eagerly evaluate non-dependent (portions of) constraints. | | | | |  | | | | |
| CA 096 | | |  | | 12.02 [over.dcl] | | | | Paragraph 1 | | | | te | | | | Declaration matching ([over.dcl]) is based upon whether trailing *requires-clause*s are equivalent; however, *equivalent*, with respect to expressions ([temp.over.link]), is defined only for expressions involving template parameters. | | | | Extend the definitions of equivalent and functionally equivalent to cover expressions subject to normalization in general (not just those involving template parameters).  Further, make the determination of expression equivalence treat concept definitions as opaque by adding a condition that an expression that may be subject to constraint normalization is functionally equivalent only if each *qualified-concept-name* that may be expanded by normalization would be considered to name the same type if, instead of a concept, a class template was named. | | | | |  | | | | |
| US 097 | | |  | | 13 | | | | 1 | | | | Te | | | | It does not seem useful to allow a type-constraint of Concept<> (directly, rather than via pack expansion). | | | | Make the template-argument-list non-optional. | | | | |  | | | | |
| US 098 | | |  | | 13 | | | | 6 | | | | Ge | | | | It is surprising that the single syntax Concept<X> can be a type-constraint (which becomes Concept<T,X>, not the Concept<X><T> that would result from adding <T> as for the other kind of type-constraint) or a very different id-expression. | | | | Add a syntactic disambiguator, perhaps in the trivial form of Concept<,X> for the type-constraint case. | | | | |  | | | | |
| US 099 | | |  | | 13.01 | | | | 01.6 | | | | Te | | | | "When a non-type template-parameter of non-reference and non-class type is used as an initializer for a reference, a temporary is always used."  follows from its prvalue status. | | | | Strike the sentence. | | | | |  | | | | |
| US 100 | | |  | | 13.01 | | | | 04.1 | | | | Te | | | | Reference types (of which there are no glvalues) seem to vacuously have strong structural equality, which would allow an rvalue reference (which is a literal type) as a template parameter. | | | | Explicitly exclude all reference types from strong structural equality. | | | | |  | | | | |
| US 101 | | |  | | 13.01 | | | | 04.1 | | | | Ge | | | | It is surprising that "template<int&,char&> int t;" is allowed, but that "struct S {int &i; char &f; bool operator==(const S&)=default/\*delete\*/;}; template<S> int t;" is not. | | | | Don't use == to define the equivalence of class-type non-type template arguments  (see comment on [temp.type]/1.5). | | | | |  | | | | |
| US 102 | | |  | | 13.01 | | | | 04.1 | | | | Ge | | | | If it is decided not to use == to define equivalence of class-type non-type template arguments (just as it is not used for references and pointers to members), some of the uncertainty surrounding non-type template parameters of floating-point type will no longer pertain, whereas the (very real) availability of undesirable workarounds involving std::bit\_cast<int>(-0.f) will persist. | | | | Apply [P1714R1](https://wg21.link/p1714r1) (as already approved by EWG and CWG). | | | | |  | | | | |
| PL 103 | | |  | | 13.01 [temp.param] | | | |  | | | | te | | | | The current syntax for constrained type template parameters, especially after the recent change of the naming convention for the standard library, causes confusion about the difference between the following two templates (one takes a value parameter, the other takes a constrained type parameter):  template<bool B> struct foo {};  template<std::boolean B> struct foo {};  This is also inconsistent with the requirement to use the keyword `auto` for variable and parameter declarations with deduced constrained type:  template<std::boolean auto B> struct foo {};  The author of this comment believes that, regardless of what the naming convention for standard library concepts ends up being, this is going to be confusing in real life code, and introduce a place where the knowledge of whether a name designates a type or a concept is necessary to be able to tell the \*kind\* of a template parameter. Therefore, this comment proposes that the second definition above would have to be written as following  template<std::boolean class B> struct foo {};  In the future, if we decide that the perceived possible confusion between the first two definitions in this comment is not actually a problem, this can be further relaxed to allow the current syntax, simiarly to how `Concept auto foo` is expected to be possible to relax in the future if the committee finds that to be desirable. | | | | In [temp.param]/1, replace the definition of the production rule *type-parameter* with:  *type-parameter:  type-constraintopt type-parameter-key ...opt identifieropt  type-constraintopt type-parameter-key identifieropt = type-id  template-head type-parameter-key ...opt identifieropt  template-head type-parameter-key identifieropt = id-expression*  Throughout the rest of the draft, replace all uses of `ConceptName TypeParameterName` with `ConceptName class TypeParameterName` (or `ConceptName typename TypeParameterName`). | | | | |  | | | | | |
| CA 104 | | |  | | 13.04 [temp.constr] | | | |  | | | | te | | | | The interaction between constraints and substitution has been the subject of some confusion. Declaration matching and partial ordering may require substitution that is not otherwise required to determine satisfaction; however, the wording does not make this clear in an accessible manner. | | | | Add a least a note, likely with examples, indicating that declaration matching and partial ordering may require substitution into constraints. Since these substitutions are not being performed as part of determining viability of candidates for overload resolution, the SFINAE process does not apply. | | | | |  | | | | |
| US 105 | | |  | | 13.04.1 | | | | 2 | | | | te | | | | Nothing prohibits forming a pointer to a non-overloaded non-template function whose constraints are not satisfied. | | | | Extend [over.over] to perform trivial overload resolution even when a function is named without a target type, obviating the need for [dcl.fct.def.delete]/2. | | | | |  | | | | |
| US 106 | | |  | | 13.04.1.2  [temp.constr.atomic] | | | |  | | | | ed | | | | Concepts use the term "atomic", which is already a term of art within the C++ standard, as evidenced by clause [atomics]. | | | | Use a term other than "atomic" for concepts. | | | | |  | | | | |
| CA 107 | | |  | | 13.04.1.4 [temp.constr.atomic] | | | | Paragraph 2 | | | | te | | | | The rules in 13.6.6.1 that the subject paragraph defers to does not handle parameter mapping for type template parameters, template template parameters, and non-type template parameters where substitution has made the expression non-dependent. | | | | For non-dependent (after substitution) members of the parameter mapping, consider types by type identity, and expressions by type and value. P1624 describes a treatment for dependent cases that defer to the declaration matching rules through alias template and variable template proxies. | | | | |  | | | | |
| US 108 | | |  | | 13.04.2 | | | | 2 | | | | Ed | | | | The possibility of type-constraints appearing in a parameter-type-list  is omitted (until /3.3.3). | | | | Mention it alongside "template-parameter-list". | | | | |  | | | | |
| US 109 | | |  | | 13.04.2 | | | | 3 | | | | Te | | | | Only templates are described as having associated constraints, but  [over.match.viable]/3 and  (via [temp.constr.order]/3) [over.match.best]/2.6  need them for non-template functions. | | | | Replace "template" with "declaration"; other declarations will simply always match one of the first two bullets. | | | | |  | | | | |
| CA 110 | | |  | | 13.04.2 [temp.constr.decl] | | | | Paragraph 3 | | | | te | | | | Overload resolution ([over.match.best]) asks us to prefer a more constrained non-template function using rules that order declarations based on their associated constraints ([temp.constr.order]), but “associated constraints” are defined for templates ([temp.constr.decl]) and not for functions. | | | | Add the following as a new paragraph before the subject paragraph:  The *associated constraints* of a non-template function is the normal form of the *constraint-expression* introduced by the trailing *requires-clause*, if any; otherwise, the function has no associated constraints. | | | | |  | | | | |
| US 111 | | |  | | 13.04.3 | | | | 1 | | | | Ge | | | | It is surprising that the very special "pseudo)" evaluation semantics of && and || are not extended to !, and in particular that !A || !B is not at all the same as !(A && B) in case of substitution failure or for subsumption. | | | | Assuming there is a rationale for the omission, add it as a note along with an example illustrating the failure of ! to invert a substitution failure. | | | | |  | | | | |
| CA 112 | | |  | | 13.04.4 [temp.constr.order] | | | |  | | | | te | | | | How template parameters from one template is to be matched against template parameters in another template when they appear in substituted parameter mappings is not clearly defined. | | | | In 13.6.6.2 [temp.func.order], candidates that are specializations of function templates should be ordered based on their constraints only when the templates have the same name (including for *conversion-function-id*s), parameter-type-list, and template parameter lists. | | | | |  | | | | |
| US 113 | | |  | | 13.05 | | | | 01.5 | | | | Te | | | | The == operator is inappropriate for comparing non-type template arguments of enumeration type, since it may be overloaded for them. | | | | Compare the values of the underlying type. | | | | |  | | | | |
| US 114 | | |  | | 13.05 | | | | 01.5 | | | | Te | | | | The == operator is inappropriate for template arguments of a class type with a member of enumeration or pointer-to-member type (and for object pointers would be incompatible with plausible extensions to [temp.arg.nontype]/2; see CWG 2043). | | | | Approach #1: Forbid class types with members of such types (to make operator== equivalent to template-argument equivalence).  Approach #2:  Apply a suitable revision of [P1837R0](https://wg21.link/p1837) that reverts not only [P0732R2](https://wg21.link/p0732r2) but also part of [P1185R2.](https://wg21.link/p1185r2)  Approach #3: Define equivalence of class-type non-type template arguments directly in terms of the (template-argument) equivalence of their base class subobjects and non-static data members (which allow to be references, but not mutable or volatile). Remove the definition of strong structural equality; restore from C++17's [temp.param]/4 bullets 1, 2, 4, and 5, or else use "a literal non-class type C for which, given an glvalue...". Directly forbid non-type template parameters of union-like class types. | | | | |  | | | | |
| US 115 | | |  | | 13.06.4  [temp.friend] | | | |  | | | | te | | | | Hidden friends that are non-templates currently cannot have a requires-clause, but this functionality is important and used throughout Ranges. | | | | Change [temp.friend]/9 to refer only to those friend declarations that are not any kind of templated entity. | | | | |  | | | | |
| US 116 | | |  | | 13.06.6.1 | | | | 6 | | | | Te | | | | There is no specification for equivalence among constraint-expressions. | | | | Presumably, define it in terms of /5's expression equivalence. | | | | |  | | | | |
| US 117 | | |  | | 13.06.6.1 | | | | 6 | | | | Te | | | | Types and type-constraints are supposed to be compared by /5, but it handles only dependent expressions. | | | | Generalize /5 to support type-ids (by recursive decomposition). Compare non-dependent types by identity; compare non-dependent type-constraints according to the rules in [temp.type]. | | | | |  | | | | |
| US 118 | | |  | | 13.07 | | | | 8 | | | | Te | | | | No reasonable implementation needs the freedom extended by making uninstantiable templates ill-formed with no diagnostic required. | | | | With the exception of the last (long) bullet, specify instead that it is unspecified whether the program is ill-formed (with a required diagnostic) when the conditions pertain. | | | | |  | | | | |
| US 119 | | |  | | 13.09.2.4 | | | | 8 | | | | Te | | | | The check that "deduction succeeds for a given type" suggests that each P/A pair is considered separately, contradicting [temp.deduct.type]/2. | | | | Directly define "at least as specialized" in terms of the overall deduction succeeding, as seems to be the implementation consensus. | | | | |  | | | | |
| US 120 | | |  | | 14.04.1.2 | | | | 2 | | | | Te | | | | There is no specification for which template parameters in (the template arguments in) one parameter mapping are "the same template parameter" (from [temp.over.link]/5) as those in another parameter mapping. | | | | Use the mappings obtained by the partial-ordering deduction (which is required to have succeeded), augmented by matching by position for template parameters of the same kind that were not deduced per [temp.deduct.partial]/12 (to support constrained function templates like std::make\_unique). | | | | |  | | | | |
| US 121 | | |  | | 15 | | | |  | | | | te | | | | [P1703](https://wg21.link/p1703) requested in a change such that import declarations of header units were lexed as preprocessing directives – those for named modules were not.   Import declarations (of either kind) are now preprocessing directives that result in tokens passed through to the C++ parser proper. The rationale for covering all import declaration was to permit source scanners operating in a non-standard preprocessing mode to extract module dependencies.   This goal is not achieved.  To achieve this goal, module declarations too must be treated as preprocessing directives. Without that, such scanners will not be able, in general, to detect module unit creation, only consumption. | | | | Either: a) revert the changes inspired by [p1703](https://wg21.link/p1703), or b) extend the changes inspired by [p1703](https://wg21.link/p1703) to module-declarations, module-private-partitions and the global-module-fragment introducer. | | | | |  | | | | |
| US 122 | | |  | | 15 | | | |  | | | | te | | | | [P1703](https://wg21.link/p1703) resulted in the creation of a preprocessing directive that does not begin with ‘#’. This is likely to confuse users, as the restricted lexing requirements come without the mnemonic ‘#’ marker. It will also complicate code formatting tools, such as editors. | | | | Revert the changes inspired by [p1703](https://wg21.link/p1703) | | | | |  | | | | |
| US 123 | | |  | | 15 | | | |  | | | | te | | | | Source scanners that do not use the complete preprocessing algorithm employ heuristics to approximate that. They will fail in some cases, whatever the standard specifies. Users of such source scanners already have to constrain the format of pieces of code to permit the scanner to function. The p1703 approach enshrines a particular scanning heuristic, with its own particular set of failing cases. For instance, when the scanner reaches an explicit header unit import, it will need to read the exported macros of that header unit. Correctly doing this requires reading the header’s Compiled Module Interface. This may be done in several ways, amongst which are: 1) reading it directly, or 2) determining the CMI’s exported macro set by processing the header unit’s source file, or 3) approximating the exported set by reading the header unit’s source file and discarding non-directive text, within the current scan. #1 requires interleaving of source scanning and compilation. The motivation of source scanners was to not do that.#2 is essentially implementing a macro-only module system inside the source scanner, which is liable to be both complex and/or inaccurate. #3 is an approximation, as it will observe macro definitions and undefinitions that are not exported. It will also incorrectly determine the prevailing macro definition algorithm of 15.3 in certain circumstances. A proposal to specify prevailing macro more in keeping with traditional #include ordering (p1174) was rejected. Thus approach #3 has been deemed undesirable. Alternative approaches of scanners overestimating the set of imports, but permitting failures, have been described more than once at meeting. | | | | Revert the changes inspired by [p1703](https://wg21.link/p1703) | | | | |  | | | | |
| US 124 | | |  | | 15 | | | |  | | | | te | | | | [P1703](https://wg21.link/p1703) relaxes the context sensitivity of the import keyword. Surrounding braces are no longer relevant – only the formatting of the line beginning with the import token (and possible preceding ‘export’ token). Previous drafts of the standard recognized the import keyword only outside of any braces (other than extern “C” linkage blocks). This requirement was motivated by [p0795](https://wg21.link/p0795), which pointed out that ‘import’ and ‘module’ were used in the user interfaces of significant software. However, the import declaration’s C++ grammar is unchanged, and it must appear at the outermost scope. The scope-agnostic lexing of the preprocessing directive will result in a) confusing errors at the parser level, and b) frustration that such uses of ‘import’ must be locally protected from the preprocessor. | | | | Revert the changes inspired by [p1703](https://wg21.link/p1703) | | | | |  | | | | |
| US 125 | | |  | | 15 | | | |  | | | | te | | | | The control-line changes to make import declarations be a directive were insufficient to achieve the stated goals of that change | | | | Apply changes from [P1857](https://wg21.link/p1857) | | | | |  | | | | |
| GB 126 | | |  | | 15 | | | |  | | | | Te | | | | P1703R1 should be part of the Modules Tooling Technical Report, not part of the standard  P1703R1 removed important features from import declarations:  the context-sensitivity no longer takes braces into account and instead matches all lines starting 'import' (breaking compatibility with existing code — a codesearch.isocpp.org search for "import" finds many cases that will be broken by the new rule), and- line continuations are now required when import declarations span multiple lines (making use of attributes on import declarations ugly and awkward).  There is also evidence that the proposal does not fully solve the problem that it aims to solve, as it does not cover module declarations. | | | | Extend the new rules to also cover module declarations, allow import declarations to span multiple lines without backslash line continuations, and consider whether the context-sensitivity can be improved so that it doesn't reject the cases found by code search.  Alternatively, revert P1703R1 from the C++ standard draft and instead establish a direction to include the rules from P1703R1 in the modules tooling technical report (as guidance on how to write code that supports dependency extraction from the widest possible set of tools). | | | | |  | | | | |
| US 127 | | |  | | 15  [cpp] | | | | ❡1 | | | | te | | | | The import contextual keyword's context is too broad and breaks real code such as import->doImport(); . | | | | Add one addtional token of context. See [P1857r0](https://wg21link/p1857) for details. | | | | |  | | | | |
| US 128 | | |  | | 15.01 | | | |  | | | | te | | | | A control line of the form ‘[export] import ...’, is intended to be passed through to the C++ parser, after lexing and header-unit macro importation. However, this is never specified. The closest we get is in 15.3 where, for header unit imports, we specify that the ‘import’ keyword is replaced by a special token. | | | | Specify (in 15.1, a new subsection, or make 15.3 more general) that these tokens are passed through. | | | | |  | | | | |
| US 129 | | |  | | 15.01 | | | | p1 | | | | TE | | | | While the \_\_has\_cpp\_attribute feature was under development in WG21, WG14 added C++-compatible attribute support to C2x and that was not taken into consideration for this feature. WG14 is considering adopting the same functionality for C but are having difficulties with the identifier chosen by WG21.  While WG14 could always pick a name like \_\_has\_c\_attribute, such a solution is unsatisfying because users would have to write twice as much code. C and C++ do not need separate preprocessor conditional inclusion features for this functionality – a single feature will suffice. e.g.,  #if \_\_has\_cpp\_attribute(something) #define SOMETHING [[something]] #elif \_\_has\_c\_attribute(something) #define SOMETHING [[something]] #endif  is exactly equivalent to the shorter:  #if \_\_has\_attribute(something) #define SOMETHING [[something]] #endif  in cases where the code in question is shared between C and C++ compilers. | | | | Rename \_\_has\_cpp\_attribute to \_\_has\_attribute or some other language-agnostic name, or alternatively, keep the name \_\_has\_cpp\_attribute and introduce a second, language-agnostic name as a synonym. | | | | |  | | | | |
| US 130 | | |  | | 15.01 [cpp.cond] | | | | 19 | | | | Te | | | | Producing a token that might be reparsed as a 'defined' operator during macro replacement has undefined behavior. Undefined behavior lexing the program has no place in a modern standard, and this should either be a diagnosable error, or (perhaps conditionally) supported behavior to become that 'defined' operator. | | | | Make this either a diagnosable error, or (perhaps conditionally) supported behavior to become that 'defined' operator. | | | | |  | | | | |
| US 131 | | |  | | 15.02 [cpp.include] | | | | 4 | | | | Te | | | | It is undefined behavior for token replacement to produce an include directive that does not match either of the two well-defined forms in the grammar. Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. | | | | Make this a diagnosable error | | | | |  | | | | |
| US 132 | | |  | | 15.03 | | | |  | | | | te | | | | Header units may provide macro definitions and undefinitions to their importers. These are #define and #undefs (of imported macros) that are encountered ‘when preprocessing each translation unit’. There is ambiguity as to whether this includes: a) macros defined on the command line, b) macros defined by the implementation (including indirectly via command line option), c) macros defined in forced headers. Mailing list discussion concluded that such macros should NOT be exported. | | | | Add wording to explicitly exclude these macro definitions (as there can be no imported macros visible at the point they are defined, explicit undefinitions are irrelevant). | | | | |  | | | | |
| US 133 | | |  | | 15.03 | | | |  | | | | te | | | | It is not clear whether header-unit source code can contain internal-linkage entities.  For example the iostream header can contain:    static ios\_base::Ioinit \_\_ioinit;  Is that permitted, or does it make the header file incompatible with being a header unit? | | | | Preference for internal linkage entities to be an error.  Thus library implementors will need an (implementation-defined) mechanism to know whether the header file is being textually included, or whether it is being processed to create a header-unit. | | | | |  | | | | |
| US 134 | | |  | | 15.03 | | | |  | | | | te | | | | It is not clear whether header-unit source code can contain  definitions of external linkage entities.  For example:    int version () { return 5;}  1) Does that emit a definition of ‘version’ to an object file associated with the header-unit? 2) Is it ill-formed? 3) Does it emit ‘version’ in the object file of each importer. #3 will lead to multiple-definition linker errors. At least 2 implementors of module compilers had differing understandings of this. Users will need to know whether an object file is a possibility. | | | | Have a slight preference for permitting emission of an object file when creating a header-unit.  i.e. option #1.  However, option #2 would also be acceptable.  Option #3 does not seem a good choice. | | | | |  | | | | |
| DE 135 | | |  | | 15.03 | | | | paragraph 2 | | | | te | | | | "import" is not a language keyword to allow for backward compatibility with existing pre-modules source code, where "import" might be used as an identifier. The current status does not achieve the desired goal; for example "import->module = ENV;" on a line is considered as an (ill-formed) module import and cannot be parsed as an expression-statement. | | | | Revert P1703R1. Document syntax restrictions to aid tools in the form of recommendations, outside of the C++ language standard. | | | | |  | | | | |
| US 136 | | |  | | 15.04 | | | |  | | | | te | | | | The optional ‘export’ keyword of a module declaration must come from source file inclusion, as it is part of the pp-balanced-token-sequence. It may come from macro expansion. The first restriction is clearly an error. The second leniency is probably a difficulty for scanners. | | | | Change the pp-global-module-fragment reduction to: module ; pp-balanced-token-seq export-opt module | | | | |  | | | | |
| US 137 | | |  | | 15.04 | | | |  | | | | te | | | | The post-phase-4 token sequences for module declarations specify that the module keyword must not come from macro expansion. (And if other comments are accepted, this will be true in more cases, and for the export keyword too.)   There is no restriction of interspersing NULL macro expansions, which again conflicts with the needs of source scanners. For instance:   module;  #define empty  empty export empty module empty foo;  Is this acceptable? | | | | Either: 1) add an example showing such empty expansion is valid or 2) add text restricting the use of such use | | | | |  | | | | |
| US 138 | | |  | | 15.04 | | | |  | | | | te | | | | The global module fragment grammar defines a pp-balanced-token-sequence. It is unclear whether the tokens of an import control line passed through to the c++ parser are part of the balanced sequence. (And therefore naked import control lines cannot appear in the GMF, and the effect of any unbalanced token sequence it might contain extends beyond the control line.) | | | | Clarify that: a) the tokens of an import control line are, and b) the tokens of other control lines are not (because they do not emit pp-tokens). Note: Implementations might emit tokens to pass a pragma directive through, but the effect is as-if that is a single internal token. | | | | |  | | | | |
| US 139 | | |  | | 15.04 & 6.5 | | | |  | | | | te | | | | When there is no Global Module Fragment, a module declaration’s tokens may be the result of macro expansion. When a GMF is present [10.4], the module token of the module declaration must not be the product of macro expansion. There is no restriction on the module keyword introducing a private module fragment in either case. This is at best inconsistent, and believed to be an error in conveying design intent. It presents difficulty with source scanners, that must therefore perform complete preprocessing to detect the module declaration in the non-GMF case. Therefore there is nothing to gain by the restrictions placed on the GMF.However, scanners could gain advantage if the restriction was applied to all module declarations. There is no implementation difficulty with either approach when compiling as specified in the std, it is purely for processing source code in an extra-standard manner. Compilers may have implementation difficulty detecting erroneous macro expansion generation as currently specified, when being given already-preprocessed tokens, as they usually cannot tell whether tokens are the result of macro expansion in that case (i.e. -fpreprocessed). This is compiling source in a manner outside the standard, so arguably not a defect. | | | | Either: a) The export & module tokens of a module-declaration, private module-fragment & global module fragment introducers must all never be the product of macro expansion, or b) No restriction on producing any export & module tokens from macro expansion. | | | | |  | | | | |
| US 140 | | |  | | 15.04/1 | | | |  | | | | te | | | | The requirements here apply only to files that lexically start with module; which means they have no effect for files where on entry to phase 7, the first token sequence forms a module declaration. Among other things, this allows a module declaration to come from an #include or macro expansion. | | | | Require that TUs that don't start with module; either start with a module declaration at the start of phase 4, or they shall not contain any module declaration in phase 7. This would also be addressed by the changes in P1857 | | | | |  | | | | |
| US 141 | | |  | | 15.05 [cpp.replace] | | | | 11 | | | | Te | | | | It is undefined behavior to have a preprocessing directive inside the parens of a macro invocation. Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error, or (perhaps conditionally) supported behavior to immediately apply that directive. For example, existing practice on many compilers is to allow an if-section, although at least one compiler is known to diagnose an error in this case. | | | | Make this either a diagnosable error, or (perhaps conditionally) supported behavior to immediately apply that directive. | | | | |  | | | | |
| US 142 | | |  | | 15.05.2  [cpp.  stringize] | | | | 2 | | | | Te | | | | It is undefined behavior for token pasting with # to produce anything that is not a valid string literal. Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. | | | | Make this a diagnosable error | | | | |  | | | | |
| US 143 | | |  | | 15.05.3 [cpp.concat] | | | | 3 | | | | Te | | | | It is undefined behavior for token pasting with ## to produce anything that is not a valid preprocessing token. Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. | | | | Make this a diagnosable error | | | | |  | | | | |
| US 144 | | |  | | 15.06 [cpp.line] | | | | 3 | | | | Te | | | | If a #line directive is given a digit sequence outside the range 1..2,147,483,647 the behavior is undefined. Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. | | | | Make this a diagnosable error | | | | |  | | | | |
| US 145 | | |  | | 15.06 [cpp.line] | | | | 5 | | | | Te | | | | If, after macro replacement, a #line directive does not exactly match one of the two supported forms, the behavior is undefined. Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. | | | | Make this a diagnosable error | | | | |  | | | | |
| GB 146 | | |  | | 15.10 | | | |  | | | | Te | | | | Concepts are missing a feature test macro  There is no feature test macro for the concepts language facility. | | | | Add a suitable definition of \_\_cpp\_concepts to [tab:cpp.predefined.ft] | | | | |  | | | | |
| GB 147 | | |  | | 15.10 | | | |  | | | | Te | | | | Add a feature-test macro for consteval | | | | Add \_\_cpp\_consteval to Table 17 [tab:cpp.predefined.ft]. | | | | |  | | | | |
| US 148 | | |  | | 15.10  [cpp.  predefined] | | | | 4 | | | | Te | | | | If a user attempts to #undef or #define a macro named 'defined', the behavior is undefined. Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. | | | | Make this a diagnosable error | | | | |  | | | | |
| US 149 | | |  | | 15.10  [cpp.  predefined] | | | | 4 | | | | Te | | | | If a user attempts to #undef or #define a predefined macro named in this clause, the behavior is undefined. Undefined behavior lexing the program has no place in a modern standard, and this should be a diagnosable error. | | | | Make this a diagnosable error | | | | |  | | | | |
| US 150 | | |  | | 15.10  [cpp.  predefined] | | | | Table 17 | | | | GE | | | | Familiar template syntax for generic lambdas should have a feature test macro: it is a significant enough feature | | | | Add \_\_cpp\_lambda\_template\_parameters | | | | |  | | | | |
| US 151 | | |  | | 16 - 32  Library | | | |  | | | | ge | | | | Please address [open LWG issues](http://www.open-std.org/jtc1/sc22/wg21/docs/lwg-active.html). | | | | Appropriate action would include making changes to the CD, identifying an issue as not requiring a change to the CD, or deferring the issue to a later point in time. | | | | |  | | | | |
| US 152 | | |  | | 16.03.4  [defns.  comparison] | | | |  | | | | Te | | | | This definition should be updated to accommodate the new 3-way comparison operator (7.6.8 [expr.spaceship]) as well. | | | | Update the definition | | | | |  | | | | |
| US 153 | | |  | | 16.04.1.3  [structure.requirements]  (library) | | | | 5 | | | | ed | | | | [P0898](https://wg21.link/P0898) applied a Cpp98 prefix (which was editorially changed to Cpp17) to all named requirements, in order to avoid ambiguity with library-defined concepts that had the same names. The named requirements are frequently part of user-facing library documentation, so changing their spelling carries a substantial educational cost. But after the application of [P1754](https://wg21.link/P1754), the standard consistently uses capital letters in the spelling of all named requirements, and consistently avoids capital letters in the names of library-defined concepts, so the prefix is no longer necessary for disambiguation. In short, the benefits of the Cpp17 prefix have evaporated, but the costs remain. | | | | Delete the Cpp17 prefix from all named requirements, and update [structure.requirements] (16.4.1.3) paragraph 5 to reflect that change. | | | | |  | | | | |
| US 154 | | |  | | 16.04.1.3 [structure.requirements] and many others | | | | n/a | | | | ed | | | | [P0898R3](https://wg21.link/p0898r3) applied a "Cpp98" prefix (which was editorially changed to "Cpp17") to all named requirements, in order to avoid ambiguity with library-defined concepts that had the same names. The named requirements are frequently part of user-facing library documentation, so changing their spelling carries a substantial educational cost. But after the application of P1754R1, the standard consistently uses capital letters in the spelling of all named requirements, and consistently avoids capital letters in the names of library-defined concepts, so the prefix is no longer necessary for disambiguation. In short, the benefits of the "Cpp17" prefix have evaporated, but the costs remain. | | | | Delete the "Cpp17" prefix from all named requirements, and update 16.4.1.3/p5 to reflect that change. | | | | |  | | | | |
| GB 155 | | |  | | 16.04.1.4 | | | |  | | | | Ed | | | | Consider renaming the "Expects:" and "Ensures:" elements in Library wording  The choice of Expects: and Ensures: for library preconditions and postconditions was done for consistency with the C++ Contracts feature. Since Contracts are not in C++20, and if they return there's no guarantee that "expects" and "ensures" will be used, we should consider reverting to more conventional terms such as "preconditions" and "postconditions". | | | | Change "Expects:" to "Preconditions:" and "Ensures:" to "Postconditions:" everywhere. | | | | |  | | | | |
| US 156 | | |  | | 16.04.2.2.6 | | | | 2 | | | | te | | | | The text reads "The type of a customization point object shall model semiregular." However, the type of a customization point objecct is very likely to be const, and const types do not model semiregular. We should instead be testing the cv-unqualified type. | | | | Should read, "The type of a customization point object ignoring cv-qualifiers shall satisfy semiregular."  See <https://cplusplus.github.io/LWG/issue3285> | | | | |  | | | | |
| US 157 | | |  | | 16.05.1.2 [headers] | | | | 4 | | | | Te | | | | The header <cstddef> should be added to the set of importable C++ library headers. It contains important C++ features like std::byte and std::nullptr\_t that are more than just C compatibility. Likewise, the C compatibility layer is essentially important vocabulary typedefs that the C++ standard library relies on and are not otherwise exported from importable C++ library header units. | | | | Add <cstddef> to the list of importable headers. | | | | |  | | | | |
| US 158 | | |  | | 16.05.1.3  [compliance] | | | |  | | | | te | | | | <coroutine> is listed as a freestanding header, however, it includes <compare>, which is not freestanding. Please ensure that <coroutine> is a freestanding header. | | | | One possible resolution would be to make <compare> a freestanding header by adopting [P1855](https://wg21.link/P1855). If that is not possible, <coroutine> could be modified to remove the dependency on <compare>. We do not consider making <coroutine> non-freestanding an acceptable solution. | | | | |  | | | | |
| US 159 | | |  | | 16.05.1.3  [compliance] | | | |  | | | | te | | | | Please ensure that <compare> is a freestanding header. | | | | Adopt [P1855](https://wg21.link/P1855). | | | | |  | | | | |
| GB 160 | | |  | | 16.05.1.3 | | | |  | | | | Te | | | | <compare> should be in freestanding implementations.  The <compare> header is closely tied to a language feature, and should be defined even for freestanding implementations. | | | | Add <compare> to tab:headers.cpp.fs in [compliance]. | | | | |  | | | | |
| PL 161 | | |  | | 16.05.1.3 [compliance] | | | |  | | | | te | | | | <compare> is currently not a freestanding header. This causes two problems:  1. It is impossible to use the spaceship operator functionality in a minimal freestanding implementation.  2. The <coroutine> header, which is freestanding, uses <compare>, which is not. | | | | Adopt P1855. | | | | |  | | | | | |
| US 162 | | |  | | 16.05.3.5  [allocator.  requirements] | | | | [tab:allocator.req.var] | | | | Te | | | | The default behavior for a.destroy is now to call destroy\_at | | | | Replace "default" entry with: destroy\_at(c) | | | | |  | | | | |
| US 163 | | |  | | 16.05.3.5 [allocator.  requirements] | | | | [tab:allocator.req.var] | | | | Te | | | | The default behavior for a.construct is now to call construct\_at | | | | Replace "default" entry with: construct\_at(c, std::forward<Args>(args)...) | | | | |  | | | | |
| FR 164 | | |  | | 16.05.4.9 | | | |  | | | | **te** | | | | While char8\_t, char16\_t and char32\_t are assumed to encode utf-8, utf-16 and utf-32 code units respectively, the encoding of u8string, u16string and u32string objects is not specified. | | | | Adopt P1880 | | | | |  | | | | | |
| DE 165 | | |  | | 16.05.5.4 | | | |  | | | | te | | | | It is unclear whether friend functions declared in a class are intended to be found via argument-dependent lookup only (and not via regular unqualified lookup), or whether the implementation is permitted to add declarations of that function that would allow unqualified lookup to succeed. For an example, see 17.11.2.2. | | | | Clarify in the vicinity of 16.5.5.4 that friend functions are found via argument-dependent lookup only, unless a synopsis (but not a detailed specification, 16.4.1.4) expressly shows a namespace-scope declaration of that function. For existing friend functions, move non-trivial definitions from the synopses to regular descriptive elements. | | | | |  | | | | |
| GB 166 | | |  | | 17.03.1 | | | |  | | | | Te | | | | The new library span does not have a feature test macro | | | | Add a definition of \_\_cpp\_lib\_span to [tab:support.ft] | | | | |  | | | | |
| US 167 | | |  | | 17.03.1  Applies to Table 36 Standard Library Feature Test Macros [tab:support.ft] | | | | Table 36 | | | | te | | | | We forgot another feature test macro. | | | | Add a new entry to the table:  \_\_cpp\_lib\_nonmember\_signed\_size | 201907L | <iterator> | | | | |  | | | | |
| DE 168 | | |  | | 17.03.1 | | | | Table 36 | | | | te | | | | No consistent policy is applied to feature-test macros involving "constexpr" annotations and related features in the standard library.  Currently, we have  \_\_cpp\_lib\_array\_constexpr (note naming deviation) \_\_cpp\_lib\_constexpr \_\_cpp\_lib\_constexpr\_dynamic\_alloc \_\_cpp\_lib\_constexpr\_invoke \_\_cpp\_lib\_constexpr\_string \_\_cpp\_lib\_constexpr\_swap\_algorithms \_\_cpp\_lib\_constexpr\_vector  Some approved papers adding "constexpr" instructed to increase the value of the generic \_\_cpp\_lib\_constexpr macro, others introduced separate macros.  Paper P0202R3 instructed to add \_\_cpp\_lib\_constexpr\_algorithms, but that was apparently never reflected in the C++ Working Draft. Paper P1424R1 resolved to use \_\_cpp\_lib\_constexpr for all constexpr-related library features, but that was apparently incompletely implemented. | | | | Apply a consistent policy to constexpr-related library features: Either provide a single feature-test macro and remove all others, or create separate ones for each feature. | | | | |  | | | | |
| DE 169 | | |  | | 17.08.2 | | | | paragraph 3 | | | | te | | | | The expectation of the note that a default argument expression involving current() causes a source\_location to be constructed that refers to the site of a function call where that default argument is needed has no basis in normative text. In particular, 9.2.3.6 paragraph 5 seems to imply that the name "current" and its semantics are bound where it appears lexically in the function declaration. | | | | Add normative text to express the desired semantics. | | | | |  | | | | |
| US 170 | | |  | | 17.11  [cmp] | | | |  | | | | te | | | | The strong\_equality and weak\_equality comparison categories don’t make sense now that we split equality from ordering. It doesn’t make sense to declare an operator<=> that returns one of these – they just add needless complexity. | | | | Remove strong\_equality and weak\_equality. Simplify three\_way\_comparable{,\_with} to only deal with the ordering categories. | | | | |  | | | | |
| US 171 | | |  | | 17.11.02.1 | | | | 4 | | | | te | | | | Substitutability is ill-defined because it circularly depends on “comparison-salient state” and it is itself used to determine the correct return time of comparisons. Comparisons define what is “comparison-salient”, and if f can distinguish between a and b, it must be examining state that the comparison did not consider salient. | | | | Either clarify the definition of substitutability, or eliminate the distinctions between strong and weak comparisons, which are the only place that definition is used. | | | | |  | | | | |
| US 172 | | |  | | 17.11.02.1 | | | | 4 | | | | Ge | | | | The concept of "substitutability" is meaningless: beyond issues like failing to require that f is pure and referring to "public const members" of something that might not be a class type, the only plausible definition of "comparison-salient state" is "any member of any notional tuple whose comparison is equivalent to that of the type", in which case everything has the property tautologically. | | | | Replace the paragraph:  "For the purposes of this subclause, a type T is said to exhibit substitutability if, given two values of type T such that a == b is true, a and b represent the same abstract value (as defined by T).".  Alternatively, remove the definition as well as std::weak\_equality and std::weak\_ordering (which tellingly are never used except to propagate their use in user classes) and consider renaming std::strong\_equality to std::equality and std::strong\_ordering to std::total\_ordering. | | | | |  | | | | |
| CA 173 | | |  | | 17.11.02.2 [cmp.weakeq] | | | |  | | | | te | | | | With the separation of <=> and ==, weak\_equality has lost its primary use (of being a potential return type of <=>). Currently weak\_equality serves no useful purpose in the standard (i.e., nothing in std acts on it), and just causes confusion (what’s the difference between weak and strong, when should I use which?)  The difference between the two is ill-defined (involving substitutability and “salient” properties, which are also vaguely defined). The best definition of equality for a type is the type’s own == operator. We should not try to sub-divide the concept of equality. | | | | Remove weak\_equality and all references to it. Rename strong\_equality to just equality.  (New wording probably requires a paper, forthcoming). | | | | |  | | | | |
| GB 174 | | |  | | 17.11.04 | | | |  | | | | Te | | | | It's confusing for equality\_comparable[\_with] and totally\_ordered[\_with] to be in a completely different clause to three\_way\_comparable[\_with]. We recommend moving [cmp.concept] to the same location as the others. | | | | Move [cmp.concept] to Clause 18 [concepts.compare] and rename the sub-clause as [concepts.threewaycomparable] and have it included in <concepts>. | | | | |  | | | | |
| GB 175 | | |  | | 17.11.06 | | | |  | | | | Te | | | | Move [cmp.object] to [comparisons]  While it's nice to have <compare>, fragmenting function objects makes library organisation difficult. It would be good to migrate this type to [comparisons], where it will be with types that are similar. | | | | Move [cmp.object] to be a sub-clause of [comparisons] and have it (additionally) included in <functional>. | | | | |  | | | | |
| US 176 | | |  | | 17.11.06 [cmp.object]  20.14.7  [comparisons]  20.14.8 [range.cmp] | | | |  | | | | Te | | | | The library defines a consistent total order for pointers in three places, but demands that only two of them be consistent. The total order in [comparisons] should be required to be the same total order as the other two subclauses. Ideally, this wording on the total order could be consolidated into one place, possibly in the clause 16 library-wide wording, and cross-referenced from these three places, simplifying the wording. | | | | Make a consistent definition in a single place, and have all three uses refer to it. | | | | |  | | | | |
| JP5 177 | | |  | | 17.11.07 | | | | p1.3 | | | | ed | | | | "ISO/IEC/IEEE 60599" is a typo. | | | | ISO/IEC/IEEE 60559 | | | | |  | | | | |
| CA 178 | | |  | | 17.11.07 [cmp.alg] | | | |  | | | | te | | | | std::strong\_order, weak\_order, and partial\_order have special cases for floating point, but are missing special casing for pointers (whereas compare\_three\_way and std::less have the special casing for pointers) | | | | 1. Change [cmp.alg] bullet 1.4 from  "Otherwise, strong\_ordering(E <=> F) if it is a well-formed expression."  to  "Otherwise, strong\_ordering(compare\_three\_way()(E, F)) if it is a well-formed expression."  2. Change [cmp.alg] bullet 2.4 from  "Otherwise, weak\_ordering(E <=> F) if it  is a well-formed expression."  to  "Otherwise, weak\_ordering(compare\_three\_way()(E, F)) if it is a well-formed expression."  3. Change [cmp.alg] bullet 3.3 from  "Otherwise, partial\_ordering(E <=> F) if  it is a well-formed expression."  to  "Otherwise, partial\_ordering(compare\_three\_way()(E, F)) if it is a well-formed expression." | | | | |  | | | | |
| FR 179 | | |  | | 17.12.03.2 | | | |  | | | | **te** | | | | coroutine\_handle::from\_address  and coroutine\_handle::address limit future evolutions while providing limited benefits | | | | Remove the functions coroutine\_handle::from\_address  and coroutine\_handle::address | | | | |  | | | | | |
| BG4 180 | | | P 524 | | 17.12.05 | | | | 1 | | | | ge | | | | (Related to BG2) The code example uses the void-returning variant of await\_suspend(). | | | | Change suspend\_never::await\_suspend() to return its argument and change suspend\_always::await\_suspend() to return nullptr. | | | | |  | | | | |
| US 181 | | |  | | 17-32 | | | |  | | | | Te | | | | The spaceship operator<=> is typically not usable unless the library header <compare> is directly included by the user. Many standard library headers provide overloads for this operator. Worse, several standard classes have replaced their existing definition for comparison operators with a reliance on the spaceship operator, and existing code will break if the necessary header is not (transitively) included. In a manner similar to the mandated library headers transitively #include-ing <initializer\_list> in C++11, these headers should mandate a transitive #include <compare>. | | | | Add: #include <compare> to the header synopsis for each of the following headers:  <array> <chrono> <coroutine> <deque> <forward\_list> <filesystem> <iterator> <list> <map> <memory> <optional> <queue> <ranges> <regex> <set> <stack> <string> <string\_view> <system\_error> <thread> <tuple> <type\_index> <unordered\_map> <unordered\_set> <utility> <variant> <vector> | | | | |  | | | | |
| US 182 | | |  | | 18  [concepts] | | | |  | | | | te | | | | [P1754](https://wg21.link/P1754) changed the naming convention for concepts in standard library from PascalCase to snake\_case. Using snake\_case for standard library concepts creates confusion for users about which standard library facilities are concepts and which are concrete types. For example:   * function (type) & invocable (concept). * iterator (type) & range (concept). * iterator (type) & input\_iterator (concept). * bool (type) & boolean (concept).   Please consider better ways of disambiguating standard library concepts from types, functions, and other kinds of things. | | | | Possible resolutions include:   * Place all standard library concepts into a nested namespace, such as std::concepts. * Add a Hungarian-notation-style prefix or suffix to standard library concepts, e.g. c\_\* or \*\_c. * Be stricter about requiring that concept names be adjectives not nouns (for example, range is a noun). * Use a different casing style for standard library concepts. | | | | |  | | | | |
| GB 183 | | |  | | 18 | | | |  | | | | Te | | | | Adopt P1716  We're currently in a partial state between the old std::relation and what's currently in the CD.  We should adopt P1716 to move to complete the change. | | | | See P1716 | | | | |  | | | | |
| GB 184 | | |  | | 18 | | | |  | | | | Te | | | | `object` should be a concept  is\_object\_v is used in multiple places around the content that ranges introduces; it feels like a fundamental core concept, and we should probably introduce this as a concept so as to not shoot ourselves in the foot. | | | | Add to [concepts]:  template<class T>  concept object = is\_object\_v<T>;  Respecify movable so that it subsumes object.  Respecify incrementable\_traits, *cond-value-type*, iterator\_traits, empty\_view, single\_view, ref\_view, filter\_view, transform\_view, take\_while\_view, drop\_while\_view, and *semiregular-box*, to require object<T> instead of is\_object\_v<T>. | | | | |  | | | | |
| US 185 | | |  | | 18.02  [concepts.  equality] | | | |  | | | | te | | | | This section talks about “implicit expression variations” but it isn’t actually clear what any of this wording means or how it is intended to be used.  See also: <https://github.com/ericniebler/stl2/issues/536> and <https://github.com/ericniebler/stl2/issues/537> | | | | Clarify the meaning of the wording and its intended use. | | | | |  | | | | |
| GB 186 | | |  | | 18.02 | | | |  | | | | Te | | | | Rename "equality preserving"  The term "equality preserving" is often called "deterministic" in other contexts. Is there a better name here, particularly considering the negative of "not required to be equality preserving". | | | | Consider changing occurrences of "equality preserving" and "equality-preserving" and "equality-preservation" to something better, possibly based on 'deterministic'.  Affects the following (sub)subclauses:  [concept.convertible]  [concept.commonref]  [concept.swappable]  [concept.equalitycomparable]  [concept.invocable]  [concept.regularinvocable]  [range.cmp]  [iterator.synopsis]  [iterator.concept.readable]  [iterator.concept.writable]  [iterator.concept.winc]  [iterator.concept.output]  [range.range] | | | | |  | | | | |
| GB 187 | | |  | | 18.02 | | | |  | | | | Te | | | | What does equality-preservation imply for user-defined concepts?  Expressions declared in a requires-expression in this document are required to be equality-preserving, except for those annotated with the comment "not required to be equality-preserving."  While the wording concerns itself with standard concepts, it does not say anything about user-defined concepts. Should it be considered standard practice for user-defined requirements to be equality-preserving unless otherwise specified too?  Example:  template<typename T>  concept spaceship\_example\_dont\_use\_me =  std::regular<T> and  std::totally\_ordered<T> and  requires(std::remove\_reference\_t<T> const& x,  std::remove\_reference\_t<T> const& y) {  x <=> y; // is this required to be equality-preserving too?  }; | | | | Provide clarification for whether or not user-defined concepts are required to be equality-preserving unless otherwise specified. | | | | |  | | | | |
| GB 188 | | |  | | 18.02 | | | |  | | | | Te | | | | Is there a possible issue with stating that replacing a constant lvalue with a non-constant lvalue should work?  What if this involves binding to a const reference, where a non-const lvalue would fail (eg. a deleted overload)? | | | |  | | | | |  | | | | |
| US 189 | | |  | | 18.02 | | | | 1 | | | | Ge | | | | The "equal" relation used to define equality preservation (and modification) is completely unspecified (except implicitly when std::equality\_comparable must be modeled), even for scalar types. It is useful for application code to rely on different definitions (e.g., comparing pointers or through them) for different algorithm calls; see also comment on [concept.moveconstructible]/1. | | | | Specify that the relation (or the abstract value) is implicitly chosen by the program for each use of the library and that the library produces results consistent with the definition (so long as the associated  semantic requirements are satisfied). Specify the (strongest) notion of equality supported by each language and library type that models the appropriate concepts. | | | | |  | | | | |
| US 190 | | |  | | 18.02 | | | | 2 | | | | Ge | | | | Defining "domain" in terms of a requirement denies using it as a property of a type (e.g., in [concept.equalitycomparable]/1.1). | | | | Treat separately the set of values supported and the set of values used (which are just the input values to an algorithm and any values it computes). | | | | |  | | | | |
| GB 191 | | |  | | 18.04.13 | | | | 01.2 | | | | Te | | | | Clean up definition of "equal"  Should English phrases such as (1.1) "u is equal to u2" be replaced by a definition using assertions on == or strong equality. What does "equal" mean for types with no == or <=> defined?  How does this relate to equality\_comparable? | | | | Consider stronger wording for the definition of "equal" (e.g. "representationally equal"). | | | | |  | | | | |
| US 192 | | |  | | 18.04.13 | | | | 1 | | | | Ge | | | | With, for example, non-empty std::unique\_ptr objects for which "equal" is defined by operator==, the move\_constructible semantic requirements are vacuous since there is no equal u2 to consult. | | | | Replace "equal" in [concepts.equality]/1 with an equally abstract notion of "value" (that is defined by the program and propagated by the library). | | | | |  | | | | |
| US 193 | | |  | | 18.04.7  [concepts.  arithmetic] | | | |  | | | | te | | | | C++20 lacks a concept for arithmetic types. This omission is surprising, as this is a fairly common use case. For example, suppose I wish to write a function that squares a number. Pre C++20, I might write:  template <typename T>  auto square(T x) {  return x \* x;  }  In C++20, it would seem natural to be able to write:  auto square(std::arithmetic auto x) {  return x \* x;  }  However, such a standard library concept is missing! Instead, we must write the more verbose:  template <typename T>  requires std::is\_arithmetic\_v<T>  auto square(T x) {  return x \* x;  } | | | | Change [concepts.arithmetic] (18.4.7) as follows:  template<class T>  concept integral = is\_integral\_v<T>;  template<class T>  concept signed\_integral = integral<T>  && is\_signed\_v<T>;  template<class T>  concept unsigned\_integral = integral<T>  && !signed\_integral<T>;  template<class T>  concept floating\_point =  is\_floating\_point\_v<T>;  template<class T>  concept arithmetic = is\_arithmetic\_v<T>; | | | | |  | | | | |
| GB 194 | | |  | | 18.04.7 | | | |  | | | | Te | | | | Respecify integral and floating\_point  Types that model integral or floating\_point also model regular, and should refine regular at a minimum. To avoid overload resolution ambiguity, we should reconsider the definition of both concepts, and introduce two additional concepts: scalar and arithmetic, which form the basis of integral and floating\_point.  It was noted that this approach now causes there to be three times as many template instantiations for all integral and floating-point types. Given that there are a relatively small and finite number of integral and floating-point types, the author is not particularly concerned with this cost, especially as we move into a world of modules.  It is unclear to the author whether or not the proposed change can be made after C++20 ships. | | | | Minimal change:  template<class T>  concept scalar = is\_scalar\_v<T> && regular<T>;  template<class T>  concept arithmetic =  is\_arithmetic\_v<T> &&  scalar<T> &&  totally\_ordered<T>;  template<class T>  concept integral = is\_integral\_v<T> && arithmetic<T>;  template<class T>  concept floating\_point = is\_floating\_point\_v<T> && arithmetic<T>;  The preferred change is the same as the above, but refines arithmetic slightly further, to account for spaceship.  template<class T>  concept arithmetic =  is\_arithmetic\_v<T> &&  scalar<T> &&  totally\_ordered<T> &&  three\_way\_comparable<T>; | | | | |  | | | | |
| US 195 | | |  | | 18.05.2 | | | |  | | | | te | | | | The boolean concept is over-complicated and fails to capture what it intends because doing so would require it to be recursive (i.e., b satisfies boolean iff the expression b && b also satisfies boolean, etc.). LEWG decided in Cologne that the boolean concept should be removed and all uses of it in the library be replaced with convertible\_to<bool>. | | | | Remove the boolean concept and replace all uses of it with convertible\_to<bool>. | | | | |  | | | | |
| US 196 | | |  | | 18.05.2  [concept.  boolean] | | | |  | | | | te | | | | The boolean concept is super complicated. Even still, it is possible to have two types that separately model boolean that still can’t be used together. A simpler formulation would be easier to understand. | | | | Consider the formulation:  template <typename T>  concept boolean = integral<remove\_cvref\_t<T>>;    This won’t accept true\_type/false\_type but at least means you can write conditions without bool casts throughout. | | | | |  | | | | |
| GB 197 | | |  | | 18.05.2 | | | |  | | | | Te | | | | Remove concept `boolean`, replace that requirement with `convertible\_to<bool>`  The concept boolean simultaneously:  - has an overly-complex specification  - is costly to check  - doesn't achieve what it was designed to do  ​ericniebler/stl2 #389 expands on more of this problem.  The best solution forward is to remove std::boolean and replace it with std::convertible\_to<bool>. | | | | Strike [concept.boolean].  Replace all occurrences of boolean in the CD with convertible\_to<bool>. Known occurrences:  [cmp.concept]  [concept.equalitycomparable]  [concept.totallyordered]  [concept.predicate] | | | | |  | | | | |
| US 198 | | |  | | 18.05.2 | | | | 1 | | | | Ge | | | | It is meaningless to have two different values b1 and b2 in the definition of the concept. | | | | Rename b1 to b; remove b2 and use b (again) instead. | | | | |  | | | | |
| GB 199 | | |  | | 18.05.3 | | | |  | | | | Te | | | | Should equality-preservation concern itself with volatile or data races?  Note text: should this mention volatile as well? Or data races? | | | | Add a mention of volatile and/or data races to the note | | | | |  | | | | |
| GB 200 | | |  | | 18.05.3 | | | | 7 | | | | Ed | | | | Example (7.2) is incomplete and doesn’t show all of the combinations of pairs or arguments: there are 6 possible pairs but only 4 are shown. b == d and a == c are missing. Similarly, the a = c examples are incomplete. | | | | Account for the missing examples. | | | | |  | | | | |
| US 201 | | |  | | 18.05.4 | | | |  | | | | te | | | | The totally\_ordered\_with<T, U> redundantly requires both common\_reference\_with<const remove\_reference\_t<T>&, const remove\_reference\_t<U>&> and equality\_comparable\_with<T, U> (which also has the common\_reference\_with requirement). The redundant requirement should be removed. | | | | Change the definition of totally\_ordered\_with to:  template<class T, class U>   concept totally\_ordered\_with =     totally\_ordered<T> && totally\_ordered<U> &&     equality\_comparable\_with<T, U> &&     totally\_ordered<       common\_reference\_t<         const remove\_reference\_t<T>&,         const remove\_reference\_t<U>&>> &&     requires(const remove\_reference\_t<T>& t,                     const remove\_reference\_t<U>& u) {       [...as before...] | | | | |  | | | | |
| GB 202 | | |  | | 18.05.4 | | | |  | | | | Te | | | | Define `totally\_ordered[\_with]` in terms of //`partially-ordered-with`//  This will simplify the definition of both concepts (particularly totally\_ordered\_with), and make them in-line with equality\_comparable[\_with].  Now that we've defined partially-ordered-with for [cmp.concept], we should consider utilising it in as many locations as possible. | | | | template<class T> concept totally\_ordered = equality\_comparable<T> && partially-ordered-with<T, T>;  template<class T, class U>  concept totally\_ordered\_with =  totally\_ordered<T> &&  totally\_ordered<U> &&  common\_reference\_with<const  remove\_reference\_t<T>&, const  remove\_reference\_t<U>&> &&  totally\_ordered<  common\_reference\_t<  const remove\_reference\_t<T>&,  const remove\_reference\_t<U>&>> &&  equality\_comparable\_with<T, U> &&  partially-ordered-with<T, U>; | | | | |  | | | | |
| GB 203 | | |  | | 18.05.4 | | | |  | | | | Te | | | | `common\_reference\_with` requirement in `totally\_ordered\_with` is redundant.  This is already required by equality\_comparable\_with, so by reshuffling the requirements, we can simplify the definition of totally\_ordered\_with. | | | | Change totally\_ordered\_with to:  template<class T, class U>  concept totally\_ordered\_with =  totally\_ordered<T> && totally\_ordered<U> &&  equality\_comparable\_with<T, U> && // moved up  totally\_ordered<  common\_reference\_t<  const remove\_reference\_t<T>&,  const remove\_reference\_t<U>&>> &&  requires(const remove\_reference\_t<T>& t,  const remove\_reference\_t<U>& u) {  { t < u } -> boolean;  { t > u } -> boolean;  { t <= u } -> boolean;  { t >= u } -> boolean;  { u < t } -> boolean;  { u > t } -> boolean;  { u <= t } -> boolean;  { u >= t } -> boolean;  }; | | | | |  | | | | |
| FR 204 | | |  | | 18.06 | | | | 1 | | | | **te** | | | | The concepts semiregular and regular require default constructibility. While default constructibility can be convenient in some cases, it can also be very harmful when there is no obvious default value for a type.  Providing a default constructor for those types is a well known source of hard to find bugs, where the initial and meaningless value can be used as if it were a real one. Type with meaningless default constructor are even worse than use of uninitialized data, because this use can be detected by tools, while the use of meaningless data cannot.  Regular is a nice name for something that should be fairly common, and adding a default constructibility requirement for regular will lead to many user types being default-constructible with no good reasons.  See also <https://quuxplusone.github.io/blog/2018/05/10/regular-should-not-imply-default-constructible/> for a more detailed discussion on this subject | | | | Remove the *semiregular* concept  Change the definition of *regular* to:  template<class T>  concept regular = copyable<T> && equality\_comparable<T>;  Adjust the text in some places, such as:  16.4.2.2.6: The type of a customization point object shall model **default\_constructible and default constructible**  An alternative would be to totally remove both *semiregular* and *regular* from the standard, since anyways these concepts are not used much. | | | | |  | | | | | |
| GB 205 | | |  | | 18.07.3 | | | |  | | | | Te | | | | The regular\_invocable name is potentially misleading as being related to the regular concept.  Suggest pure\_invocable or similar to indicate that neither the function or the arguments are changed. Since this is only a semantic difference from invocable then a clear name would help  If the “equality preserving” term were changed to deterministic (or similar) then deterministically\_invocable might be suitable (and contrasts nicely in meaning with non\_deterministically\_invocable). | | | | Possible alternatives to regular\_invocable:  pure\_invocable  consistent\_invocable  Or  deterministically\_invocable  deterministic\_invocable  If regular\_invocable is renamed to either of these latter two, consider also renaming invocable to non\_deterministic(ally)\_invocable. | | | | |  | | | | |
| GB 206 | | |  | | 18.07.3 | | | |  | | | | Te | | | | What is the intention for regular\_invocable?  The definition of regular\_invocable states that calling invoke is equality-preserving and that neither the function object, nor its arguments are modified. A chat with Casey Carter in Cologne about why the function object can't also be const-qualified revealed that the intention of regular\_invocable is to refine invocable so that it's equality-preserving. The author is not convinced that the current wording is in sync with this hallway discussion (by one of its designers).  Examples:  auto eq1 = [](auto const x) { return x \* x; };  auto eq2 = [](auto& x) { x \*= x; };  auto eq3 = [&x]{ x \*= x; };  The author's understanding is that all three of these lambdas are equality-preserving, but only decltype(eq1) models regular\_invocable. | | | | Please confirm that the definition of regular\_invocable is correct.  If it is correct, please consider requiring that the function object also be const-qualified (this will help to prevent changes to the function object). | | | | |  | | | | |
| US 207 | | |  | | 19.05.2.5 [syserr.errcat.objects] | | | |  | | | | te | | | | The lifetime of the objects returned by functions like std::system\_category() is unclear. Because these objects are meant to be referred to by std::error\_code values, issues over the lifetime of the error category objects exposes use of std::error\_codes during program termination ([basic.start.term]) to undefined behavior. | | | | Provide a convenient mechanism to establish the lifetime of all similar error category objects associated with the implementation in one shot (perhaps in the style of [iostream.objects.overview]). Encourage implementations to allow references to the associated objects as-if their lifetime began during constant initialization before that of any object with a non-trivial destructor. | | | | |  | | | | |
| DE 208 | | |  | | 20.02.1 [utility.syn] | | | |  | | | | te | | | | Comparing and converting numbers of different numeric types is, should  be a trivial task, but it's not because of implicit conversion  P0586 was voted in Cologne, it adds free functions for comparing  different numeric types as if they where signed types. | | | | Adopt P0586 as discussed in Cologne and commented on GitHub  (<https://github.com/cplusplus/papers/issues/259>) | | | | |  | | | | |
| US 209 | | |  | | 20.04.2  [string.view] | | | |  | | | | te | | | | string\_view should be made to be constructible from any contiguous character range in the new Ranges world. | | | | Adopt [P1391](https://wg21.link/p1391) | | | | |  | | | | |
| CA 210 | | |  | | 20.05 [tuple], 20.7 [variant], 31 [atomics], Annex D [depr.\*] | | | |  | | | | te | | | | Deprecate some uses of volatile in the standard library. | | | | Adopt P1831. | | | | |  | | | | |
| US 211 | | |  | | 20.05 [tuple], 20.7 [variant], 31 [atomics], Annex D[depr.\*] | | | |  | | | | te | | | | Deprecate the library uses of volatile which were voted for deprecation by LEWG. | | | | Adopt [P1831R0](https://wg21.link/p1831r0). | | | | |  | | | | |
| US 212 | | |  | | 20.07.3.1  [variant.ctor] | | | |  | | | | te | | | | Resolve LWG 3228: surprising variant construction | | | | Resolve LWG 3228 | | | | |  | | | | |
| US 213 | | |  | | 20.10.08.2 | | | | 17 | | | | te | | | | uninitialized\_construct\_using\_allocator should use construct\_at instead of operator new | | | | *Effects:* Equivalent to: return ~~::new(static\_cast<void\*>(p))~~ construct\_at(p,  ~~T(~~make\_obj\_using\_allocator<T>(alloc, std::forward<Args>(args)...))); | | | | |  | | | | |
| US 214 | | |  | | 20.10.10 | | | | 1 | | | | te | | | | propagation traits for std::allocator are inconsistent: POCMA and POCS should never differ | | | | Add to allocator class definition:  using propagate\_on\_container\_move\_assignment = true\_type;  using propagate\_on\_container\_swap = true\_type;  using is\_always\_equal = true\_type; | | | | |  | | | | |
| US 215 | | |  | | 20.10.11 [specialized.  algorithms] | | | | 6 | | | | TE | | | | The ‘voidify’ change introduced in ‘The One Ranges Proposal’ damages const correctness. | | | | Delete *voidify*, and change places using *voidify* back to static\_cast<void\*>. | | | | |  | | | | |
| US 216 | | |  | | 20.11.08  [util.smartpr.atomic]  31  [atomics] | | | |  | | | | ed | | | | Please move the section specifying the atomic<shared\_ptr<T>> and atomic<weak\_ptr<T>> specializations from [utilities] (Clause 20) to [atomics] (Clause 31) so that it is located within the same section as the rest of atomic<T>. If this text is not relocated, it is more likely that the atomic<shared\_ptr<T>> specializations will be overlooked in future changes to atomic<T>. We have encountered this same issue in the past with the <numeric> algorithms, which previously lived in [numerics], and were frequently overlooked when updates were made to [algorithms]. Moving the section that the text is in is purely an editorial change and does NOT imply changing which header the specializations are in. | | | | Move [util.smartptr.atomic] (20.11.8) from [utilities] (Clause 20) to right after [atomics.types.memop] (31.8.5) in [atomics] (Clause 31). E.g. Make [util.smartptr.atomic] 31.8.6. The stable tag should not be changed. | | | | |  | | | | |
| US 217 | | |  | | 20.12.03 & 20.12.3.2 | | | | 2 | | | | te | | | | polymorphic\_allocator::allocate\_object and new\_object should be [[nodiscard]] | | | | Add [[nodiscard]] in front of the return type for allocate\_object and new\_object in class declaration and in member-function description for polymorphic\_allocator template. | | | | |  | | | | |
| JP6 218 | | |  | | 20.12.03.2 | | | | p1 | | | | ed | | | | It's better to use a C++ property than C standard library macro, SIZE\_MAX. | | | | Replace "SIZE\_MAX" with "numeric\_limits<size\_t>::max()" | | | | |  | | | | |
| JP7 219 | | |  | | 20.12.03.2 | | | | p8.1 | | | | ed | | | | It's better to use a C++ property than C standard library macro, SIZE\_MAX. | | | | Replace "SIZE\_MAX" with "numeric\_limits<size\_t>::max()" | | | | |  | | | | |
| US 220 | | |  | | 20.14.08 | | | | 2 | | | | Te | | | | The implementation-defined total order should be the same as that used by [comparisons]. | | | | State so. | | | | |  | | | | |
| US 221 | | |  | | 20.14.08 | | | | 3, 7 | | | | Ge | | | | Requiring the conversions to be equality-preserving is meaningless absent a definition of equality for the pointer type (which serves to constrain the definition for the parameter types). | | | | Define it: in this case, in terms of the implementation-defined total order. | | | | |  | | | | |
| FR 222 | | |  | | 20.15.10 | | | |  | | | | **te** | | | | std::is\_constant\_evaluated is easily misused, since it will always be true in if constexpr conditions, | | | | Make std::is\_constant\_evaluated a language feature by adopting P1938 (if consteval {}) | | | | |  | | | | | |
| GB 223 | | |  | | 20.20.02 | | | |  | | | | Te | | | | What does "not a format string" mean?  std::format throws when the relevant argument "is not a format string", but [format.string] doesn't clearly say when a given input is "not a format string". Is "{a}" a format string consisting of those verbatim characters (because it doesn't match the grammar for a replacement field) or is it not a format string? | | | | Define how "errors" in a format string are treated. | | | | |  | | | | |
| GB 224 | | |  | | 20.20.02 | | | |  | | | | Te | | | | Format string grammar is in terms of narrow characters only  The BNF grammar for format strings is specified in terms of char literals like '{' but it's not clear what that means for wide character strings such as L"{}". | | | | Clarify the (obvious) mapping from wide characters to terminals in the grammar, i.e. L'{' is equivalent to '{' etc.  Consider using the same grammar style as the core language, instead of a modified BNF. | | | | |  | | | | |
| GB 225 | | |  | | 20.20.02.2 | | | |  | | | | Ed | | | | std::format() alignment specifiers should be independent of text direction  The align specifiers for formatting standard integer and string types are expressed in terms of "left" and "right". However, "left alignment" as currently defined in the format() specification might end up being right-aligned when the resulting string is displayed in a RTL or bidirectional locale.  This ambiguity can be resolved by removing "left" and "right" and replacing with "start" and "end", without changing any existing implementation and without changing the intent of the feature. | | | | In [tab:format.align]:  Forces the field to be <del>left-aligned within</del> <ins>aligned to the start of</ins> the available space  and  Forces the field to be <del>right-aligned within</del> <ins>aligned to the end of</ins> the available space | | | | |  | | | | |
| GB 226 | | |  | | 20.20.02.2 | | | |  | | | | Te | | | | Make locale-dependent formats for std::format() congruent with default formatting  The design of format() prefers "locale-independent" formatting options for performance reasons. It provides very limited support for locale-dependent formatting via the 'n' specifier.  It's particularly problematic that the 'n' specifier for floating point numbers is specifically limited to the chars\_format::general presentation. It would be very useful to have access to chars\_format::scientific and chars\_format::fixed formatting with locale-dependent presentation.  Adding these features to std::format() at this stage would require significant wording changes that are too large to contain in a comment. However, one approach that could be taken in the future would be to make 'n' be an additional suffix that could be added to format specifiers, rather than being a lone format specifier. This would enable locale-dependent formatting of any of the conversions of any of the arithmetic types.  In order to keep the design space open for making this change in a future version of the standard, it would be ideal for 'n' conversions to always be congruent with the default conversion. It provides an intuitive semantic: 'n' is the same as "no specifier", but with locale-dependent presentation.  The integer and charT presentation types currently specify 'n' conversions that are congruent with the default conversion.  The bool and floating-point presentation types have 'n' conversions that are not congruent with the default conversion.  For C++20:  - Remove the 'n' conversion for bool.  auto s format("{:n}", 1);  // Committee Draft: s contains "1"  // Proposed: ill-formed format string  Making the 'n' conversion for floating-point match the default conversion, i.e. dependent on whether a precision is specified.  auto s format("{:n} {:2n}", 12.345678, 12.345678);  // Committee Draft: s contains "12,3456 12,34"  // Proposed: s contains "12,345678 12,34"  These changes are the minimum necessary to allow enhanced support for locale-dependent formatting in the standard library to be added in a backwards-compatible way in a future edition of C++. | | | | In [tab:format.type.bool]: Remove n.  In [tab:format.type.float]: Replace the 'Meaning' of the n specifier with:  If precision is specified, equivalent to to\_chars(first, last, value, chars\_format::general, precision), where precision is the specified formatting precision; equivalent to to\_chars(first, last, value) otherwise. The context's locale is used to insert the appropriate digit group and decimal radix separator characters. | | | | |  | | | | |
| US 227 | | |  | | 20.20.02.2 [format.  string.std] | | | | 5  Table 59 | | | | te | | | | We believe that the lack of a way to suppress the negative sign on numbers which are rounded up to zero by the specified precision is a defect which will affect most users of format string. | | | | Add: “'z' Indicates that a sign should not be used for negative numbers that display as zero (after rounding to the formatting precision).” To table 59. The details of the change will be proposed in P1496R1 in the pre-Belfast mailing. A “D” version of this paper was discussed in Kona this year. | | | | |  | | | | |
| US 228 | | |  | | 20.20.02.2 [format.string.std] | | | | Paragraph 7, Paragraph 9 | | | | te | | | | Units of width and precision are not specified which causes an ambiguity for strings in variable-length encodings. | | | | Width and precision for strings should be computed based on fixed operating system dependent encodings. If the operating system is capable of displaying Unicode text in a terminal both ordinary and wide encodings are Unicode encodings such as UTF-8 and UTF-16, respectively. [ Note: this is the case for Windows-based and many POSIX-based operating systems. -- end note ] Otherwise encodings are implementation-defined. For the given encoding, display width of a string is the number of column positions needed to display the string in a terminal [ Note: This is similar to the semantics of the POSIX wcswidth function with a fixed encoding. —] | | | | |  | | | | |
| GB 229 | | |  | | 20.20.03 | | | |  | | | | Te | | | | Formatting functions don't allow throwing on incorrect arguments  std::format is only allowed to throw if fmt is not a format string, but the intention is it also throws for errors during formatting, e.g. there are fewer arguments than required by the format string. | | | | Allow exceptions even when the format string is valid. Possibly state the \_Effects:\_ more precisely. | | | | |  | | | | |
| DE 230 | | |  | | 21.03.2.1 | | | |  | | | | te | | | | Because string::reserve() can no longer shrink the capacity, it should be added at the end of §4.2 as one function where a non-const member function can not invalidate referencess, pointers, and iterators, if it does not grow the capacity | | | | add string::reserve() at the end of §4.2 | | | | |  | | | | |
| DE 231 | | |  | | 21.03.3.5 ([string.erasure]) 22.3.8.5 ([deque.erasure]) 22.3.9.7 ([forward.list.erasure]) 22.3.10.6 ([list.erasure]) 22.3.11.6 ([vector.erasure]) 22.4.4.5 ([map.erasure]) 22.4.5.4 ([multimap.erasure]) 22.4.6.3 ([set.erasure]) 22.4.7.3 ([multiset.erasure]) 22.5.4.5 ([unord.map.erasure]) 22.5.5.4 ([unord.multimap.erasure]) 22.5.6.3 ([unord.set.erasure]) 22.5.7.3 ([unord.multiset.erasure]) | | | |  | | | | te | | | | The free erase/\_if functions were moved from LFTSv2 to the IS, but P1115  fell through the cracks. It would be awkward to ship a version of free  erase()/erase\_if() with a known API issue (returning void instead of the  number of elements removed), and then fix it up in a source-incompatible  way in C++23. | | | | Adopt P1115. | | | | |  | | | | |
| US 232 | | |  | | 21.04.2 [string.view] | | | |  | | | | Te | | | | Make string\_view constructible from contiguous ranges of character type. This is an important integration with the new ranges facility, and should not be deferred to a later standard.  We believe having basic\_string\_view be properly constructible from a range should be viewed not as a new feature but as fixing a “missing constructor” defect resulting from the integration of ranges and therefore feel this is in scope for NB comments. | | | | Apply [p1391](https://wg21.link/p1391) | | | | |  | | | | |
| US 233 | | |  | | 21.07.3  [views.span] | | | |  | | | | te | | | | span’s constructors should be harmonized with the new Ranges concepts of contiguous\_range and contiguous\_iterator, needs to be done now. | | | | Adopt [P1394](https://wg21.link/p1394) | | | | |  | | | | |
| GB 234 | | |  | | 22 | | | |  | | | | Te | | | | Adopt P1115R0 for C++20  P1209R0 added erase and erase\_if functions for the containers. P0646R1 changed the remove members of list and forward\_list to return the number of removed elements. We failed to coordinate these changes, meaning the non-member erase functions discard the useful information now returned by forward\_list::remove.  P1115R0 proposed to fix this, but isn't in the CD. | | | | Adopt P1115R0 as an obvious defect in the new erase and erase\_if functions.  Note this affects multiple locations in clause 22 | | | | |  | | | | |
| US 235 | | |  | | 22.02.7  [unord.reg] | | | | 11  17  Table 70 | | | | Te | | | | C++20 design fix: the use of Hash::transparent\_key\_equal to enable heterogeneous lookup for unordered associative containers deviates from prior art, does not address the incompatibility concerns raised in the original LEWG review, and adds more subtle and confusing corner cases and will likely surprise and confuse the user.  For details on the problem, see <https://isocpp.org/files/papers/P1690R1.html#design-minimize-confusion>  [P1690R0](https://wg21.link/p1690r0) proposed a fix that was reviewed by LEWG in Cologne, which was forwarded to LWG for C++20 (poll results below) with a suggestion to prioritize as it would be a break to do later; unfortunately LWG ran out of time before getting to it.  SF F N A SA  5 6 7 0 0 | | | | See: [P1690r1](https://wg21.link/p1690r1)  For details on the problem, see <https://isocpp.org/files/papers/P1690R1.html#design-minimize-confusion> | | | | |  | | | | |
| US 236 | | |  | | 22.02.7  [unord.req] et. al. | | | | n/a | | | | te | | | | The working paper has an implementation of heterogeneous lookup that differs substantially from existing practice. Once we ship the design currently in the CD, we will have a difficult time retrofitting the design in [P1690](https://wg21.link/p1690).  LEWG reviewed and approved [P1690](https://wg21.link/p1690) for C++20 in Cologne, but limited LWG review time prevented this from being moved. | | | | Merge [P1690](https://wg21.link/p1690) into the working paper | | | | |  | | | | |
| PL 237 | | |  | | 22.02.7 [unord.req] | | | |  | | | | te | | | | Heterogenous lookup for unordered containers requires hasher to provide the transparent\_key\_equal nested type that denotes the predicate. This design is inconsistent with the method used for the ordered containers and existing non-standard implementations, that checks for nested is\_transparent type. Furthermore, it prevents the implementation of generic hashers to be combined with dedicated type predicate. Finally, it overrides std::equal\_to equality predicate, even in a situation when it is explicitly provided by the user. | | | | Adopt P1690R0. | | | | |  | | | | | |
| US 238 | | |  | | 22.02.7 [unord.req] [N4810] | | | | 11, 17  Table 70. | | | | te | | | | C++20 design fix: the use of Hash::transparent\_key\_equal to enable heterogeneous lookup for unordered associative containers deviates from prior art, does not address the incompatibility concerns raised in the original LEWG review, and adds more subtle and confusing corner cases and will likely surprise and confuse the user.   For details on the problem, see <https://isocpp.org/files/papers/P1690R1.html#design-minimize-confusion>   [P1690R0](https://wg21.link/p1690r0) proposed a fix that was reviewed by LEWG in Cologne, which was forwarded to LWG for C++20 (poll results below) with a suggestion to prioritize as it would be a break to do later; unfortunately LWG ran out of time before getting to it.  SF F N A SA  5 6 7 0 0 | | | | For proposed wording, see [p1690R1](https://wg21.link/p1690r1) | | | | |  | | | | |
| US 239 | | |  | | 22.03.7.1 | | | | 2 | | | | Te | | | | There is no specification of whether std::array has strong structural equality. | | | | Specify that it has no non-static data members other than the obvious array (and see also comment on [class.compare.default]/4.2.1). | | | | |  | | | | |
| FR 240 | | |  | | 22.07 | | | |  | | | | **te** | | | | | span::index\_type’s name is inconsistent with the convention used by other containers and views, notably string\_view | | | | Rename span::index\_type to span::size\_type | | | | |  | | | | | |
| US 241 | | |  | | 22.07 | | | |  | | | | te | | | | Rename std::dynamic\_extent to std::dyn, as repeatedly using such a long name in the upcoming mdspan proposal ([P0009](https://wg21.link/P0009), slated for Library Fundamentals V3) is unnecessarily unwieldly. | | | | Replace dynamic\_extent with dyn throughout the subsections of 22.7. | | | | |  | | | | |
| US 242 | | |  | | 22.07  [ views ] | | | |  | | | | Ed | | | | This early view type should be editorially consolidated into the new section for views in general, rather than lying in the containers clause. | | | | Move into 24 [Ranges] | | | | |  | | | | |
| FR 243 | | |  | | 22.07.2 | | | |  | | | | **te** | | | | Both std::extent and the proposed std::static\_extent are type traits, while std::dynamic\_extent is not, which is surprising and inconsistent | | | | Rename std::dynamic\_extent to std::dynamic\_extent\_tag | | | | |  | | | | | |
| FR 244 | | |  | | 22.07.2 | | | |  | | | | **te** | | | | std::as\_bytes and std::as\_writable\_bytes encourage undefined behavior | | | | Consider removing these functions. | | | | |  | | | | | |
| US 245 | | |  | | 22.07.3 | | | |  | | | | te | | | | [P1227R2](https://wg21.link/P1227R2) changed the size and indexing operations in span from the signed type ptrdiff\_t to the unsigned type size\_t. The typedef should be changed from index\_type to size\_type to be consistent and interoperable with the rest of the standard library. | | | | Replace index\_type with size\_type as per [P1872R0](https://wg21.link/P1872R0). | | | | |  | | | | |
| US 246 | | |  | | 22.07.3   [views.span] | | | |  | | | | Te | | | | Span should be constructible from a contiguous forwarding range or iterators thereof, and not just "Container" types with a data() member function. As this has the possibility of affecting overload resolution and SFINAE, it may not be possible to repair std::span in a later standard. | | | | Apply [P1394](https://wg21.link/p1394) | | | | |  | | | | |
| PL 247 | | |  | | 22.07.3 [views.span] | | | |  | | | | te | | | | span<T> provides a const-qualified begin() method and cbegin() method that produces a different result if T is not const-qualifed:  1) begin() produces mutable iterator over T (as if T\*)  2) cbegin() preduces const iterator over T (as if T const\*)  As consequence for the object s of type span<T>, the call to the std::cbegin(s)/std::ranges::cbegin(s) produces different result than s.cbegin(). | | | | Change span<T> members cbegin()/cend()/crbegin()/crend()/const\_iterator to be equivalent to begin()/end()/rbegin()/rend()/iterator respectivelly. | | | | |  | | | | | |
| PL 248 | | |  | | 22.07.3 [views.span] | | | |  | | | | te | | | | std::span uses the name `index\_type` instead of `size\_type` for the return type of its `size` function. There is a historical reason for this; std::span used to have a signed return type of `size`. This typedef is also used as a type for "index" or "count" parameters, but since they are all unsigned at this point, it seems like an unwarranted inconsistency with the rest of the standard library. | | | | Either: 1. Rename std::span::index\_type to size\_type. 2. Add an additional alias, size\_type, aliasing index\_type, to std::span. | | | | |  | | | | | |
| US 249 | | |  | | 22.07.3.1 | | | |  | | | | te | | | | Remove const\_pointer and const\_reference from span, as they are unused. | | | | **~~using const\_pointer = const element\_type\*;~~**  using reference = element\_type&;  **~~using const\_reference = const element\_type&;~~** | | | | |  | | | | |
| PL 250 | | |  | | 22.07.3.2 [span.cons] | | | |  | | | | te | | | | The resolution of the LWG3101 prevents accidental undefined behavior caused by size mismatch between the range and constructed span, e.g.:  void processFixed(span<int, 5>); void processDynamic(span<int>); std::vector<int> v;  processFixed(v); // ILL-FORMED after 3103, UB if v.size() != 5 before  processDynamic(v); // OK  However, the resolution does not prevent similar problems in the situation when the (ptr, len) or (ptr, ptr) constructor is used:  processFixed({v.data(), v.size()}); // WELL-FORMED, UB if v.size() != 5  processFixed({v.data(), v.data() + v.size()}); // WELL-FORMED, UB if v.size() != 5  Morover, currently, the code remains ill-formed even if explicit cast is performed by the user:  processFixed(span<int, 5>(v)); // ILL-FORMED  To resolve the issue, the construction of fixed-size span from dynamic-sized range should be explicit:  processFixed(v); // ILL-FORMED  processFixed({v.data(), v.size()}); // ILL-FORMED  processFixed({v.data(), v.data() + v.size()}); // ILL-FORMED  processFixed(span<int, 5>(v)); // WELL-FORMED  processFixed(span<int, 5>{v.data(), v.size()}); // WELL-FORMED  processFixed(span<int, 5>{v.data(), v.data() + v.size()}); // WELL-FORMED  To summarize:  Source | Destination | Constructor  Fixed | Fixed | Implicit, ill-formed if size-mismatch  Fixed | Dynamic | Implicit  Dynamic | Dynamic | Implicit  Dynamic | Fixed | Explicit | | | | Add 'explicit(extent != dynamic\_extent)' specifier to the following constructors in [span.cons]:  constexpr span(pointer ptr, index\_type count);  constexpr span(pointer first, pointer last);  In the specification of constructors:  template<class Container> constexpr span(Container& cont);  template<class Container> constexpr span(const Container& cont);  \* Add 'explicit(extent != dynamic\_extent)' specifier.  \* Remove 'extent == dynamic\_extent is true' ([span.cons]p 14.1) from Constrains element.  \* Add 'If extent is not equal to dynamic\_extent, then size(cont) is equal to extent.' to Expects element.  In the specification of constructor:  template<class OtherElementType, size\_t OtherExtent>  constexpr span(const span<OtherElementType, OtherExtent>& s) noexcept;  \* Add 'explicit(extent != dynamic\_extent && OtherExtent == dynamic\_extent)' specifier.  \* Replace the 'Extent == dynamic\_extent || Extent == OtherExtent is true' constrain with 'Extent == dynamic\_extent || OtherExtent == dynamic\_extent || Extent == OtherExtent is true'.  \* Add 'If extent is not equal to dynamic\_extent, then s.size() is equal to extent.' to Expects element. | | | | |  | | | | | |
| PL 251 | | |  | | 22.07.3.2 [span.cons] | | | |  | | | | te | | | | C++20 introduced both std::contiguous\_range and std::continous\_iterator concepts, that are the generalization of the pointer to continuous sequence of objects, and type erased view for such ranges in form of std::span.  However, these two features are not integrated together, as consequence std::span cannot be directly constructed from ranges that models std::contiguous\_range and std::sized\_range, nor from the pair of std::continous\_iterator:  std::vector v{...};  std::span s = v; // OK  std::span s = v | std::take\_view(10); // ILL-FORMED  std::span s(std::to\_address(v.begin()), 2); //OK  std::span s(std::to\_address(v.begin()), std::to\_address(v.begin() + 2)); //OK  std::span s(v.begin(), 2); // IMPLEMENTATION-DEFINED  std::span s(v.begin(), .begin() + 2); // IMPLEMENTATION-DEFINED | | | | Adopt P1394R3. | | | | |  | | | | | |
| CA 252 | | |  | | 22.07.3.7 [span.objectrep] | | | |  | | | | te | | | | as\_writable\_bytes standardizes UB. In particular, pointer interconvertibility between an object and its object representation (in array form) is not established. We should not hide reinterpret\_cast inside another std function.  Also, as\_writable\_bytes and as\_bytes should not be free functions unless we plan on applying them to other std types (e.g., vector). Free functions should be designed as function overload sets or as functions acting on a concept (i.e., all containers or all views, etc.—not necessarily a C++ Concept). | | | | Preferred: Remove as\_writable\_bytes and move as\_bytes to be a member function of span.  Alternative: Rename as\_writable\_bytes to something including the word “reinterpret”, such as reinterpret\_as\_bytes, and make it a member function (along with as\_bytes as a member function). | | | | |  | | | | |
| US 253 | | |  | | 23  24  25 | | | |  | | | | te | | | | The adoption of [P1207](https://wg21.link/p1207) (movability of single-pass iterators) has left the working paper in an inconsistent state. Many places both in [range.adaptors] and [algorithms] assume copyability of input and output iterators, an assumption [P1207](https://wg21.link/p1207) invalidated by permitting input (but not forward) and output iterators to be move-only.  For instance, here are three issues in filter\_view alone:  In [range.filter.iterator]/p5, the current\_ member of a filter\_view::iterator is copied:   constexpr iterator\_t<V> base() const;   Effects: Equivalent to: return current\_;  In [range.filter.iterator]/p7 has the same problem:   constexpr iterator\_t<V> operator->() const   requires has-arrow <iterator\_t<V>>;   Effects: Equivalent to: return current\_;  In [range.filter.iterator]/p8, we are copying out of the current\_ member in a call to find\_if:   constexpr iterator& operator++();   Effects: Equivalent to:     current\_ = ranges::find\_if(++current\_, ranges::end(parent\_->base\_), ref(\*parent\_->pred\_));     return \*this;  As an example from the [algorithms] clause, here is [alg.rotate]/p11, which is shown erroneously copying a potentially move-only output iterator:   template<forward\_range R, weakly\_incrementable O>     requires indirectly\_copyable<iterator\_t<R>, O>       constexpr ranges::rotate\_copy\_result<safe\_iterator\_t<R>, O>         ranges::rotate\_copy(R&& r, iterator\_t<R> middle, O result);   Effects: Equivalent to:     return ranges::rotate\_copy(ranges::begin(r), middle, ranges::end(r), result);  [P1207](https://wg21.link/p1207) introduced an unknown but likely large number of bugs into the working paper. Tracking them all down would take time and leave us with little confidence that we had found them all.  In addition, [p1456](https://wg21.link/p1426) (Move-only views) was \_not\_ merged to the working draft, leaving us in the oddly inconsistent state where iterators could be move-only but views could not. This has caused yet more bugs. For instance, views::counted(first,n) returns a view that holds an iterator by value. If the iterator is move-only, then the resulting view is not a view because it fails to satisfy the copyability requirement of the view concept. | | | | Revert [P1207](https://wg21.link/p1207), restoring the copyability requirement to the weakly\_incrementable concept. | | | | |  | | | | |
| GB 254 | | |  | | 23 | | | |  | | | | Te | | | | Most of the ranges iterator operations should be marked [[nodiscard]]  These are equality-preserving operations that return values always intended to be used. The library should reflect this. | | | | Add [[nodiscard]] to the following operations.  ranges::iter\_move  ranges::distance  ranges::next  ranges::prev | | | | |  | | | | |
| GB 255 | | |  | | 23 | | | |  | | | | Te | | | | output\_iterator and output\_range shouldn't be concepts  Very little uses these concepts, and it's not clear if they're actually necessary at all, since they're explicitly omitted in N3351 (see §3.7). The author's understanding is that they were added to mitigate potential confusion among users familiar with STL output iterators.  We should be judicious about the concepts that we introduce. If an output\_iterator concept proves itself to be useful, then we can probably add it in C++23. | | | | Proposed Change:  Strike [iterator.concept.output] and associated references.  Strike output\_range from [range.refinements] or transform it into an exposition-only output-iterator.  Re-specify the following algorithms so that they require weakly\_incrementable && writable instead of output\_iterator, or to require *output-iterator* (an exposition-only concept):  replace\_copy  replace\_copy\_if  fill  fill\_n | | | | |  | | | | |
| GB 256 | | |  | | 23 | | | |  | | | | Te | | | | iterator concepts belong in namespace ranges  This will (hopefully) help solidify that the iterator concepts introduced by C++20 aren't a one-to-one mapping between STL iterators and ranges iterators (similarly to how ranges algorithms aren't identical to std algorithms).  It'll also let us use the name iterator instead of input\_or\_output\_iterator. | | | | Move all concepts in [iterators] into namespace ranges.  Rename input\_or\_output\_iterator to iterator. | | | | |  | | | | |
| US 257 | | |  | | 23.02  and  24.2 | | | |  | | | | te | | | | Two of the ranges opt-in variable templates are negative and checked against, the other is positive and checked for. Double negatives are needlessly difficult to understand. Make all the opt-in variable templates enable\_meow instead of disable\_meow. | | | | Adopt [P1871](https://wg21.link/p1871). | | | | |  | | | | |
| US 258 | | |  | | 23.03.1 [iterator.  requirements.general] | | | | 10 | | | | Ed | | | | It is unhelpful for the library to overload the definition of reachable with the core language definition of reachable for modules. Based on usage throughout this clause, suggest including the following 'from' in the defined words of power. This is consistent with every intended use of the current term, and no subsequent usage requires (nor uses) italics on either word. | | | | Change font to italics on the word 'from': is called reachable from an iterator i to is called reachable from an iterator i | | | | |  | | | | |
| US 259 | | |  | | 23.03.2.3 | | | | 03.3 | | | | te | | | | Types satisfying input\_iterator but not equality\_comparable look like C++17 output iterators. This issue is discussed in detail in [LWG#3283](https://cplusplus.github.io/LWG/issue3283). | | | | Adopt the proposed resolution at <https://cplusplus.github.io/LWG/issue3283> | | | | |  | | | | |
| US 260 | | |  | | 23.03.2.3/p4 23.5.4.2/p1 24.7.4.3/p3 24.7.7.3/p3 24.7.8.3 24.7.8.5/p1 | | | |  | | | | te | | | | It is currently impossible to non-intrusively opt-out of conformance to the C++17 iterator concepts without also opting out of conformance to the C++20 iterator concepts. This is a corner case that was missed when the Ranges TS was merged into namespace std. The issue is discussed in depth in [LWG#3289](https://cplusplus.github.io/LWG/issue3289). | | | | Adopt the proposed resolution in <https://cplusplus.github.io/LWG/issue3289> | | | | |  | | | | |
| US 261 | | |  | | 23.03.4.13 | | | | 02.6 | | | | te | | | | The expression ++(a + D(n - 1)) is erroneously applying pre-increment to an rvalue iterator. This is not required to be valid for random access iterators. | | | | Replace with either  [](I c){ return ++c; }(a + D(n - 1)) or with next(a + D(n - 1)).  See <https://cplusplus.github.io/LWG/issue3277> | | | | |  | | | | |
| US 262 | | |  | | 23.03.4.13 | | | | 02.7 | | | | te | | | | The semantic constraints of the random\_access\_iterator concept is accidentally promoting the difference type using unary negate. | | | | Change (b += -n) to (b += D(-n)).  See <https://cplusplus.github.io/LWG/issue3284> | | | | |  | | | | |
| US 263 | | |  | | 23.03.4.2 | | | |  | | | | te | | | | In the current spec, shared\_ptr<int> is readable, but shared\_ptr<int>& is not. That is because readable\_traits is not stripping top-level references before testing for nested typedefs. | | | | Change every occurance of iter\_value\_t<In> in the definition of the readable concept with iter\_value\_t<remove\_reference\_t<In>>.  See <https://cplusplus.github.io/LWG/issue3279> | | | | |  | | | | |
| US 264 | | |  | | 23.03.4.2 | | | |  | | | | te | | | | The readable concept is both under- and over-constrained. It is under-constrained in that it permits its associated types (iter\_value\_t, iter\_reference\_t, etc) to differ depending on whether the type is const-qualified or not. It \_might\_ make sense for iter\_reference\_t to be sensitive to const-ness if, for example, it is our intention for a type like optional to satisfy readable. Generally we use readable to constrain types that are logical indirections; e.g., pointers (smart and dumb) and iterators. For those, top-level cv-qualification should not matter.  readable is over-constrained because it only requires operator\* to be valid on a (non-const) lvalue.  See discussion at <https://github.com/ericniebler/stl2/issues/514> | | | | Change the definition of the readable concept to correct these problems. | | | | |  | | | | |
| FR 265 | | |  | | 23.03.4.6 | | | |  | | | | **te** | | | | input\_or\_output\_iterator does not denote input\_iterator<It> || output\_iterator<It>, which sets a bad precedent for concept naming and may not match the user intent.  It is also at odds with the naming used for decades including most literature and Stepanov’s work. | | | | Rename it to general\_iterator | | | | |  | | | | | |
| US 266 | | |  | | 23.03.5.3 | | | | Table 85 | | | | Te | | | | What does it mean for an output iterator to be incrementable after any number of increments? | | | | Add a note explaining the satisfaction of the property from [iterator.concept.winc]/13, provide an alternate definition, or remove the Ensures | | | | |  | | | | |
| US 267 | | |  | | 23.03.6.2  indirectcallable.  indirectinvocable,  others | | | |  | | | | te | | | | The ranges compare algorithms are over-constrained. LEWG approved [P1716](https://wg21.link/p1716) with the correct fix, but LWG ran out of time to review it. Without [P1716](https://wg21.link/p1716), safe and correct programs will be erroneously rejected. | | | | Adopt the proposed resolution in [P1716](https://wg21.link/p1716), which has already passed LEWG design review. | | | | |  | | | | |
| US 268 | | |  | | 23.03.7.4 | | | |  | | | | te | | | | The indirectly\_swappable concept is over-constrained: it requires only that iter\_swap is callable with lvalue iterators. It should be possible to call iter\_swap with rvalues as well.  See discussion at <https://github.com/ericniebler/stl2/issues/578>. | | | | Change the concept to require iter\_swap to be callable with both lvalue and rvalue iterators. | | | | |  | | | | |
| DE 269 | | |  | | 23.07 Range access | | | | Paragraph: 18 | | | | te | | | | There are ranges that model std::ranges::sized\_range, but do not provide a .size() member function. (Also) for this reason std::ranges::size() was introduced with slightly different semantics than std::size(). Now we are introducing std::ssize() with the semantics of std::size() + signed-ness. This means we get three size functions that each have different deficiencies and none that works for all sized ranges and is signed. | | | | Preferred: Make std::ssize() resolve to std::ranges::size() + signed-ness.  Alternative: Also add std::ranges::ssize(). | | | | |  | | | | |
| GB 270 | | |  | | 24 | | | |  | | | | Te | | | | P1207 provided an opportunity for us to weaken input iterators so that they don't need to be copyable. While the author thinks that this is a step in the correct direction, P1207 has left us in a partial state where iterators don't need to be copyable, but views do. Given that views have underlying iterators, we need to address this problem before C++20 ships. | | | | Either apply P1456 and evaluate all standard range adaptors to determine if they're affected (and then apply changes to bring them into accordance with P1207 and P1456), or completely roll back P1207. | | | | |  | | | | |
| GB 271 | | |  | | 24 | | | |  | | | | Te | | | | Many operations in namespace ranges should be marked nodiscard  These are equality-preserving operations that return values always intended to be used. The library should reflect this intention in the strongest way possible. | | | | Add [[nodiscard]] to the following operations.  ranges::begin  ranges::end  ranges::cbegin  ranges::cend  ranges::rbegin  ranges::rend  ranges::crbegin  ranges::crend  ranges::size  ranges::empty  ranges::data  ranges::cdata  Add [[nodiscard]] to the following view\_interface member functions.  empty  data  size  front  back  operator[]  Add [[nodiscard]] to the following subrange member functions.  begin  end  empty  Size | | | | |  | | | | |
| US 272 | | |  | | 24  Applies to §24 Ranges [ranges] | | | | §24.4.2 Ranges [range.range];  §24.6.1.2 Class template empty\_view [range.empty.view];  §24.6.3 Class template iota\_view[range.iota.view];  §24.7.6.4 views::take [range.take.adaptor];  §24.7.8.3 views::drop [range.drop.adaptor]. | | | | te | | | | Due to the ranges API being more or less fixed after shipment, API-breaking fixes have to be scheduled now rather than shipped in C++23. The papers [p1664](https://wg21.link/p1664) and [p1739](https://wg21.link/p1739)  represent important and ultimately source-breaking changes to the ranges API. Unless shipped, these API optimizations will result in source code breaking at a later date if attempted to be fixed later, and also makes it impossible to reliably simplify the return value of ranges for a wide variety of current and future adaptors and algorithms. These changes are imperative for developing better APIs and it would be unfortunate to not be able to do them post-C++20 due to source breaking changes. All of the changes except for [p1664](https://wg21.link/p1664)'s two new exposition-only concepts have been approved by LEWG during review of [p1739](https://wg21.link/p1739). However, [p1664](https://wg21.link/p1664) was discussed as part of [p1739](https://wg21.link/p1739)'s approval and this comment was to be expected. | | | | The fix and its motivations have been formalized in two papers, [[p1664](https://wg21.link/p1664) - Reconstructible Ranges] and [[p1739](https://wg21.link/p1739) - Type preservation for forwarding Ranges for "subrange-y" views]. Further fixes are applied through Corentin Jabot's and Casey Carter's [[p1391](https://wg21.link/p1391)] and [[p1394](https://wg21.link/p1394)]. | | | | |  | | | | |
| US 273 | | |  | | 24.02 | | | |  | | | | Te | | | | all\_view is not a view like the others. For the other view types, foo\_view{args...} is a valid way to construct an instance of type foo\_view. However, all\_view is just an alias to the type of view::all(arg), which could be one of several different types. all\_view feels like the wrong name. | | | | Suggest renaming all\_view to all\_t and moving it into the views:: namespace. | | | | |  | | | | |
| GB 274 | | |  | | 24.02 | | | |  | | | | Te | | | | Add range\_size\_t  LEWG asked that range\_size\_t be removed from P1035, as they were doing a good job of being neutral w.r.t whether or not size-types were signed or unsigned at the time.  Now that we've got a policy on what size-types are, and that P1522 and P1523 have been adopted, it makes sense for there to be a range\_size\_t. | | | | Add to [ranges.syn]:  template<range R>  using range\_difference\_t = iter\_difference\_t<iterator\_t<R>>;  + template<sized\_range R>  + using range\_size\_t = decltype(ranges::size(declval<R&>())); | | | | |  | | | | |
| GB 275 | | |  | | 24.03 | | | |  | | | | Te | | | | ranges::begin and ranges::end should not accept arrays of unknown bound  The current definitions of ranges::begin and ranges::end mean that an array of unknown bound is treated as an empty range. The expressions E+0 and E+extent\_v<T> are both well-formed for an array of unknown bound (with extent\_v<T> equal to zero). | | | | Make ranges::begin(E) and ranges::end(E) ill-formed when E is an array of unknown bound. | | | | |  | | | | |
| US 276 | | |  | | 24.03.1  24.3.2  24.5.3  24.7.3.1  24.6.3.2  several | | | | 01.3  1.3 | | | | te | | | | Several of the range views define non-template friend function begin/end overloads taking rvalues to satisfy the exposition-only forwarding-range concept. These have a couple of problems. First, the ones for subrange take subrange&&. That means that a const rvalue subrange fails to satisfy forwarding-range, which causes cbegin(subrange{...}) to be ill-formed.  The bigger problem is that since these functions are non-templates, whenever they get added to the overload set, the compiler will try conversions to these types (subrange, ref\_view). The attempted conversions could lead to errors in theory.  Finally, class iota\_view has iterator that can safely outlive the view that created them, so it too should be given begin/end friend functions that accept rvalues following the same pattern. | | | | In [range.access.begin]/p1.3, change the poison-pill overloads from:  template<class T> void begin(T&&) = delete;  template<class T> void begin(initializer\_list<T>&&) = delete;  ...to:  template<class T> void begin(T&&) = delete;  template<class T> void begin(initializer\_list<T>) = delete;  To the synopsis in [range.subrange]/p1, add:  template<class A, class B>  concept same-ish = // exposition only    same\_as<A const, B const>;  In the class synopsis of subrange (same section), change the `begin`/`end` friend functions from:  friend constexpr I begin(subrange&& r) { return r.begin(); }  friend constexpr S end(subrange&& r) { return r.end(); }  ...to:  friend constexpr I begin(same-ish<subrange> auto && r) { return r.begin(); }  friend constexpr S end(same-ish<subrange> auto && r) { return r.end(); }  In the synopsis of `ref\_view` in [range.ref.view]/p1, change the `begin`/`end` friend functions from this:  friend constexpr iterator\_t<R> begin(ref\_view r)  { return r.begin(); }  friend constexpr sentinel\_t<R> end(ref\_view r)  { return r.end(); }  ...to this (editors note: the use of same\_as here instead of same-ish is intentional; likewise for the use of pass-by-value):  friend constexpr iterator\_t<R> begin(same\_as<ref\_view> auto r)  { return r.begin(); }  friend constexpr sentinel\_t<R> end(same\_as<ref\_view> auto r)  { return r.end(); }  To the class synopsis of iota\_view in [range.iota.view], add the following `begin`/`end` friend functions:  friend constexpr W begin(same-ish<iota\_view> auto && r) { return r.begin(); }  friend constexpr auto end(same-ish<iota\_view> auto && r) { return r.end(); }  See <https://github.com/ericniebler/stl2/issues/592>. | | | | |  | | | | |
| GB 277 | | |  | | 24.04 | | | |  | | | | Te | | | | Adopt P1456 or change istream\_view's requirements  P1456 weakens view so that it does not require copyable. Without this, istream\_view is unable to process non-copyable types as input.  LEWG approved P1456 in Kona, but it seems that it didn't make it in time for LWG in Cologne. I don't think this is something that can be fixed in C++23. | | | | Apply the proposed wording in P1456 to the International Standard. | | | | |  | | | | |
| GB 278 | | |  | | 24.04 | | | |  | | | | Te | | | | Rename viewable\_range to viewable  The name viewable\_range doesn't communicate its intended purpose very clearly. Consider renaming to viewable, which very clearly reflects its description, and thus intended purpose. You can't convert a non-range to a view anyway. | | | | Rename viewable\_range to viewable. Affects:  [ranges.syn]  [range.adaptor.object]  [range.filter.view]  [range.transform.view]  [range.take.view]  [range.join.view]  [range.common.view]  [range.reverse.view] | | | | |  | | | | |
| US 279 | | |  | | 24.04.2  [range.  range] | | | |  | | | | te | | | | The *forwarding-range* concept opt-in is too subtle and just adds complexity to overload resolution. The other range concepts use variable templates to opt-in, this one should do. | | | | Adopt [P1870](https://wg21.link/p1870). | | | | |  | | | | |
| GB 280 | | |  | | 24.04.2 | | | |  | | | | Ed | | | | forwarding-range is too-easily confused with forward\_range | | | | Please find an alternative name for forwarding-range, as it is extremely similar to forward\_range, which is (a) different in definition, and (b) a user-facing concept. | | | | |  | | | | |
| FR 281 | | |  | | 24.04.4 | | | |  | | | | **te** | | | | The View concept requires copy-ability. There are reasons why this is overly restrictive - for example a predicate might not be copyable or a view might need to hold a coroutine\_handle which should not be copied.  As concepts are hardly modifiable after the publication of the standard, it is important to relax this constraint while we still can | | | | Adopt P1456 which was approved by LEWG | | | | |  | | | | | |
| DE 282 | | |  | | 24.04.4 | | | |  | | | | te | | | | "Since the difference between range and view is largely semantic, the two are differentiated with the help of enable\_view." (§3)  enable\_view is designed as on opt-in trait to specify that a type is a view. It defaults to true for types derived from view\_base (§4.2) which is clearly a form of opt-in. But it also employs a heuristic assuming that anything with iterator == const\_iterator is also view (§4.3). This is a very poor heuristic, the same paragraph already needs to define six exceptions from this rule for standard library types (§4.2).  Experience in working with range-v3 has revealed multiple of our own library types as being affected from needing to opt-out from the "auto-opt-in", as well. This is counter-intuitive: something that was never designed to be a view shouldn't go through hoops so that it isn't treated as a view. | | | | Make enable\_view truly be opt-in by relying only on explicit specialisation or inheritance from view\_base. This means removing 24.4.4 §4.2 - §4.4 and introducing new §4.2 "Otherwise, false".  Double-check if existing standard library types like basic\_string\_view and span need to opt-in to being a view now. | | | | |  | | | | |
| US 283 | | |  | | 24.05.1 | | | |  | | | | te | | | | The exposition-only has-arrow concept is ill-formed. It has a constrained template parameter, which is not valid C++20. | | | | Change the concept to:  template<class I> concept has-arrow = // exposition only   input\_iterator && (is\_pointer\_v<I> || requires(I i) { i.operator->(); }); | | | | |  | | | | |
| US 284 | | |  | | 24.05.3 | | | |  | | | | te | | | | Conversion from pair-like types to subrange is a silent semantic promotion. Just because a pair is holding two iterators does not mean those two iterators denote a valid range. Permitting that pair to be implicitly converted to a subrange is error prone. | | | | In the synopsis of subrange, strike the definition of the exposition-only \_pair-like-convertible-to\_ concept, and the following two subrange constructors:  template<not-same-as<subrange> PairLike>   requires pair-like-convertible-to<PairLike, I, S> constexpr subrange(PairLike&& r) requires (!StoreSize);  template<pair-like-convertible-to<I, S> PairLike> constexpr subrange(PairLike&& r, make-unsigned-like-t(iter\_difference\_t<I>) n)   requires (K == subrange\_kind::sized);  See <https://cplusplus.github.io/LWG/issue3281> | | | | |  | | | | |
| US 285 | | |  | | 24.05.3 | | | | 1 | | | | te | | | | The subrange converting constructors permit derived-to-base slicing errors. See detailed discussion of this issue in [LWG#3282](https://cplusplus.github.io/LWG/issue3282). | | | | Adopt the proposed resolution at <https://cplusplus.github.io/LWG/issue3282> | | | | |  | | | | |
| US 286 | | |  | | 24.06.3.2  [range.iota.  view] | | | |  | | | | te | | | | Iota\_view is currently under-constrained and does not behave as a forwarding-range. | | | | Adopt LWG 3292, and add the correct opt-in for *forwarding-range* (dependent on earlier NB comment) | | | | |  | | | | |
| US 287 | | |  | | 24.06.3.3 | | | |  | | | | te | | | | iota\_view::iterator has the wrong iterator\_category. Depending on the capabilities of the template parameter W, the category could be anything from input\_iterator\_tag to random\_access\_iterator\_tag. However, according to the \_Cpp17InputIterator\_ requirements, iota\_view::iterator cannot satisfy any of the old iterator concepts stronger than input. That is because its operator\* returns a prvalue. | | | | Adopt the proposed resolution in <https://cplusplus.github.io/LWG/issue3291>. | | | | |  | | | | |
| DE 288 | | |  | | 24.07 | | | |  | | | | te | | | | By fully specifying the types returned by view adaptors, the standard forces the return type of multiple chained view operations to be an increasingly nested template. In general this is not avoidable, but for certain combinations of input ranges and view adaptors, one can simply create a modified object of the original type (e.g. with different bounds). Not addressing this is a design flaw that needlessly complicates working with views. | | | | 1. Adopt P1739 and in this context also P1391 and P1394 which are strongly suggested by P1739. All three papers have been seen and approved by LEWG. P1739 cannot be adopted after C++20 without breaking API. 2. Discuss whether P1664 (or parts of it) should also be adopted. P1664 generalises the notion of "reconstructible ranges" (those affected by P1739). 3. Evaluate whether any other combinations of range and view adaptor should get special treatment; any such changes after C++20 are breaking.generalises the notion of "reconstructible ranges" (those affected by P1739). 3. Evaluate whether any other combinations of range and view adaptor should get special treatment; any such changes after C++20 are breaking. | | | | |  | | | | |
| DE 289 | | |  | | 24.07.1 | | | |  | | | | te | | | | "Given an additional range adaptor closure object D, the expression C | D is well-formed and produces another range adaptor closure object"  Experience in combining range-v3 with our library's views has revealed that it is very difficult to satisfy the above requirement in a generic way since it is not defined how code can identify "range adaptor closure objects" and which entity is responsible for combining the two closures into one. This leads to incompatible implementations and conflicting overloads. It may even become impossible for developers to target different standard library implementations at the same time -- depending on how these chose to implement the above rule. | | | | 1) Introduce a boolean trait called enable\_range\_adaptor\_closure that must be specialised for the type of all range adaptor closure objects.  2) Specify that the standard library implements the aforementioned combining of two-into-one in an implementation-defined manner (e.g. a free function operator| that works on any two objects whose types satisfy enable\_range\_adaptor\_closure and that returns the respective combined closure object).  [Note that this only affects operator| for combining two closure objects -- not for piping a range into a closure object. The trait would however also enable the standard library to provide a generic implementation of the latter so that users need only provide operator() for their closure objects. This in turn could make it possible to create closure objects from lambdas.] | | | | |  | | | | |
| US 290 | | |  | | 24.07.10 | | | |  | | | | te | | | | The strange is\_reference\_v<iter\_reference\_t<iterator\_t<V>>> || view<iter\_value\_t<iterator\_t<V>>> constraint was correct before [P0970](http://wg21.link/P0970) and the *forwarding-range* concept. Now it is inexact and wordy. What we are really looking for is a *forwarding-range;* that is, a range on which we can call view::all to get a view which we can store within the join\_view cheaply. | | | | Change the requirements on the join\_view class template from:  template<input\_range V>   requires view<V> && input\_range<range\_reference\_t<V>> &&     (is\_reference\_v<range\_reference\_t<V>> || view<range\_value\_t<V>>) class join\_view;  to:  template<input\_range V>   requires view<V> && input\_range<range\_reference\_t<V>> &&     forwarding-range<iter\_reference\_t<iterator\_t<V>>> class join\_view; | | | | |  | | | | |
| US 291 | | |  | | 24.07.10.2 | | | |  | | | | te | | | | The non-const join\_view::begin() returns iterator<simple-view<V>>. If simple-view<V> is true, then the iterator stores a const join\_view\* named parent\_. iterator::satisfy() will try to write to parent\_->inner\_ if ref\_is\_glvalue is false. That doesn't work because the inner\_ field is not marked mutable. | | | | In [range.join.view], change the join\_view<V>::inner\_ member to be mutable. This is safe because this exposition-only member is only used when the join\_view is single-pass and only modified by operations that invalidate other iterators.  See <https://cplusplus.github.io/LWG/issue3278> | | | | |  | | | | |
| US 292 | | |  | | 24.07.10.2 | | | |  | | | | te | | | | join\_view::iterator's constructor is incorrect. In join\_view<V>::iterator<Const>, we see the constructor:    constexpr iterator(Parent& parent, iterator\_t<V> outer)  V above is the non-const-qualified view template parameter. We will then try to initialize the outer\_ data member with outer, which has type iterator\_t<Base>, where Base is const V when Const is true, and V otherwise. This is broken; there is no required conversion if the types are different. Fixing this will probably require changes also to join\_view's begin() and end() members. | | | | In [range.join.view], change the join\_view<V>::inner\_ member to be mutable. This is safe because this exposition-only member is only used when the join\_view is single-pass and only modified by operations that invalidate other iterators.  See <https://cplusplus.github.io/LWG/issue3278> | | | | |  | | | | |
| US 293 | | |  | | 24.07.10.2 | | | |  | | | | te | | | | join\_view is missing a base() member for returning the underlying view. All the other range adaptors provide this. | | | | To the join\_view class template add the member:    constexpr V base() const { return base\_; } | | | | |  | | | | |
| US 294 | | |  | | 24.07.10.3 | | | | 14,15 | | | | te | | | | join\_view::iterator::operator-- is improperly constrained. In the Effects clause in paragraph 14, we see the statement:    inner\_ = ranges::end(\*--outer\_);  However, this only well-formed when end returns an iterator, not a sentinel. This requirement is not reflected in the constraints of the function(s). | | | | Change join\_view::iterator::operator--() and operator--(int) to the following:  constexpr iterator& operator--() requires ref\_is\_glvalue && bidirectional\_range<Base> && bidirectional\_range<range\_reference\_t<Base>> && common\_range<range\_reference\_t<Base>>;  constexpr iterator operator--(int) requires ref\_is\_glvalue && bidirectional\_range<Base> && bidirectional\_range<range\_reference\_t<Base>> && common\_range<range\_reference\_t<Base>>; | | | | |  | | | | |
| US 295 | | |  | | 24.07.10.3 | | | | 2,3 | | | | ed | | | | Paras 2.1 and 3.2 do not say what the iterator\_(category|concept) should be if neither of the two sub-bullets hold. Presumably the author intended those bullets to fall through to p2.2 and p3.3 respectively, but I don't think it works that way. | | | | Add a 2.1.3 that reads, "Otherwise, iterator\_concept denotes input\_iterator\_tag." Add a 3.2.3 that reads, "Otherwise, iterator\_category denotes input\_iterator\_tag.“ | | | | |  | | | | |
| US 296 | | |  | | 24.07.11.3 | | | |  | | | | te | | | | split\_view::outer\_iterator converting constructor is slightly wrong. In split\_view::outer\_iterator<V, Pattern>, when V is not const-iterable, we must avoid forming the type iterator\_t<const V> since it will fail to compile. | | | | For the converting constructor:  constexpr outer\_iterator(outer\_iterator<!Const> i)   requires Const && ConvertibleTo<iterator\_t<V>, iterator\_t<const V>>;  change the requirement to:    requires Const && ConvertibleTo<iterator\_t<V>, iterator\_t<Base>>; | | | | |  | | | | |
| US 297 | | |  | | 24.07.11.4 | | | |  | | | | te | | | | The value\_type of the split\_view iterator is a view; however, unlike all the other view types in the ranges clause, this one does not inherit from view\_interface. This is an oversight. This must be corrected now as doing so later would change ABI. | | | | Change the synopsis of split\_view's outer\_iterator to show struct split\_view<V, Pattern>::outer\_iterator<Const>::value\_type inheriting from view\_interface<value\_type>.  See  <https://cplusplus.github.io/LWG/issue3276> | | | | |  | | | | |
| US 298 | | |  | | 24.07.13.3 | | | | 01.1 | | | | Te | | | | The behavior of views::common cannot depend on modeling a concept. | | | | Change "models" to "satisfies". | | | | |  | | | | |
| GB 299 | | |  | | 24.07.16.2 | | | |  | | | | Te | | | | has-tuple-element helper concept needs convertible\_to  The exposition-only has-tuple-element concept (for elements\_view) is defined as  template<class T, size\_t N>  concept has-tuple-element = *exposition only*  requires(T t) {  typename tuple\_size<T>::type;  requires N < tuple\_size\_v<T>;  typename tuple\_element\_t<N, T>;  { get<N>(t) } -> const tuple\_element\_t<N, T>&;  };  However, the return type constraint for { get<N>(t) } is no longer valid under the latest concepts changes | | | | Change to:  template<class T, size\_t N>  concept has-tuple-element = *exposition only*  requires(T t) {  typename tuple\_size<T>::type;  requires N < tuple\_size\_v<T>;  typename tuple\_element\_t<N, T>;  { get<N>(t) } -> convertible\_to<const tuple\_element\_t<N, T>&>;  }; | | | | |  | | | | |
| US 300 | | |  | | 24.07.2 | | | | 1 | | | | Te | | | | The behavior of semiregular-box cannot depend on modeling a concept. | | | | Change "model{s,ed}" to "satisfie{s,d}". Add semantic constraints on the use of semiregular-box if necessary. | | | | |  | | | | |
| GB 301 | | |  | | 24.07.4.2 | | | |  | | | | Te | | | | filter\_view has no pred() member  Other views taking predicates (take\_while\_view and drop\_while\_view) have a pred() member returning (a const reference to) the contained predicate object, but filter\_view does not | | | | In [range.filter.view], add  constexpr const Pred& pred() const;  Effects: Equivalent to: return \*pred\_; | | | | |  | | | | |
| US 302 | | |  | | 24.07.4.2 24.7.5.2 24.7.6.2 24.7.10.2 24.7.11.2 24.7.14.2 | | | |  | | | | te | | | | Several of the view class templates in the [range.adaptors] section have converting constructors from compatible ranges. These were originally added in a misguided effort to support CTAD, but as described in [LWG#3280](https://cplusplus.github.io/LWG/issue3280), these constructors can cause recursion in the type constraints, leading to spurious compile errors. | | | | Adopt the proposed resolution in <https://cplusplus.github.io/LWG/issue3280> | | | | |  | | | | |
| US 303 | | |  | | 24.07.5.2 | | | |  | | | | Te | | | | The transform\_view does not constrain the return type of the transformation function. It is invalid to pass a void-returning transformation function to the transform\_view, which would cause its iterators’ operator\* member to return void. | | | | Change the constraints on transform\_view to the following:  template<input\_range V, copy\_constructible F>   requires view<V> && is\_object\_v<F> &&     regular\_invocable<F&, range\_reference\_t<V>> &&     can-reference<invoke\_result\_t<F&, range\_reference\_t<V>>> class transform\_view | | | | |  | | | | |
| US 304 | | |  | | 24.07.6.2 24.5.3.1/p6 | | | |  | | | | te | | | | On an input (but not forward) range, begin(rng) is not required to be an equality-preserving expression (24.4.2 [[range.range]](https://wg21.link/range.range)/3.3). If the range is *also* sized, then it is not valid to call size(rng) after begin(rng) (24.4.3 [[range.sized]](https://wg21.link/range.sized)/2.2). In several places in the ranges clause, this precondition is violated. A trivial re-expression of the effects clause fixes the problem. | | | | Adopt the proposed resolution in <https://cplusplus.github.io/LWG/issue3286> | | | | |  | | | | |
| FR 305 | | |  | | 25 | | | |  | | | |  | | | | The range version of some algorithms are missing | | | | Adopt P1243 | | | | |  | | | | | |
| US 306 | | |  | | 25  [algorithms] | | | |  | | | | te | | | | The ranges comparison algorithms are overconstrained – they require symmetric comparison functions even though the algorithm doesn’t need them. The constraints should be lowered.  See also <https://github.com/ericniebler/stl2/issues/610> | | | | Adopt [P1716](https://wg21.link/p1716) | | | | |  | | | | |
| US 307 | | |  | | 25  [algorithms] | | | |  | | | | te | | | | Some algorithms do not have ranges:: counterparts. | | | | Adopt [P1243](https://wg21.link/p1243). | | | | |  | | | | |
| GB 308 | | |  | | 25 | | | |  | | | | Te | | | | All half ranges should be fully rangified  It seems odd that we're not offering full ranges for the ranges that we write to. We can potentially eliminate a class of error by requiring all ranges have bounds, and implementations can optimise for the unreachable\_sentinel\_t case.  This is already the case for the uninitialised memory algorithms. | | | | Redesign all ranges algorithms with half-ranges so that they're fully bounded.  Example:  // Current  template<input\_iterator I, sentinel\_for<I> S, weakly\_incrementable O>  requires indirectly\_copyable<I, O>  constexpr ranges::copy\_result<I, O> ranges::copy(I first, S last, O result);  template<input\_range R, weakly\_incrementable O>  requires indirectly\_copyable<iterator\_t<R>, O>  constexpr ranges::copy\_result<safe\_iterator\_t<R>, O> ranges::copy(R&& r, O result);  // Proposed  template<input\_iterator I, sentinel\_for<I> S1,  input\_or\_output\_iterator O, sentinel\_for<O> S2>  requires indirectly\_copyable<I, O>  constexpr ranges::copy\_result<I, O> ranges::copy(I first, S1 last, O result, S2 result\_last);  template<input\_range R, range O>  requires indirectly\_copyable<iterator\_t<R>, iterator\_t<O>>  constexpr ranges::copy\_result<safe\_iterator\_t<R>, safe\_iterator\_t<O>> ranges::copy(R&& r, O&& result); | | | | |  | | | | |
| GB 309 | | |  | | 25 | | | |  | | | | Te | | | | Strike ranges::\*\_n algorithms  Most \*\_n algorithms become unnecessary in the wake of counted\_iterator, and relevant range adaptors. | | | | Remove the following ranges algorithm overloads:  copy\_n  fill\_n  generate\_n  (Note: search\_n omitted, as it appears to be different to the others.) | | | | |  | | | | |
| GB 310 | | |  | | 25 | | | |  | | | | Te | | | | Some new algorithms should be marked [[nodiscard]]  These algorithms are equality-preserving, and are ultimately read operations. They should be marked as [[nodiscard]] to reflect that their result is always intended to be used.  Consider similar change for lexicographical\_compare\_three\_way. | | | | Add [[nodiscard]] to the following ranges:: algorithms.  All algorithms in [alg.nonmodifying].  remove (not a read operation, but the result should rarely be discarded)  remove\_if (not a read operation, but the result should rarely be discarded)  is\_sorted  is\_sorted\_until  All algorithms in [alg.binarysearch].  is\_parititoned  partition\_point  includes  is\_heap  All algorithms in [alg.min.max].  lexicographical\_compare  Possibly also add nodiscard to lexicographical\_compare\_three\_way. | | | | |  | | | | |
| CZ 311 | | |  | | 25 | | | | 24.5 | | | | te | | | | Due to the ranges API being more or less fixed after shipment, API-breaking fixes have to be scheduled now rather than shipped in C++23. Currently, many of the new algorithms in std::ranges algorithms copy their boolean-returning predecessors by returning a single boolean value. And while this makes perfect sense as an independent unit, individuals composing these algorithms with iterators that are bidirectional or worse sometimes need to perform additional actions around and because of the return of one of these boolean-returning algorithms. For example, if an iterator is advanced to its corresponding "last" value by std::equal, that advancement is lost upon returning just a boolean from the algorithm. Any work that wanted to continue from the "last" value supplied into the algorithm must re-increment the iterator, resulting in duplicated work. | | | | The algorithms (the new ones in std::ranges) should be changed to have a result type which is (explicitly) convertible to boolean and also retains the Iterator value at its state. If the algorithm is "successful" (e.g., std::equal returns true), the iterator must point at the end. Otherwise, the value of the returned iterator is unspecified. Purportedly, an upcoming paper [P1877](https://wg21.link/p1877) will handle this. | | | | |  | | | | |
| PL 312 | | |  | | 25 [algorithms] | | | |  | | | | te | | | | The ranges::is\_permutation, ranges::unique, ranges::unique\_copy algorithm are currently underconstrained, as they do not require supplied functor to model equivalence relation.  Other compare algorithms (like ranges::equal, ranges::mismatch, ranges::search), are overconstrained. They require supplied functor to model relation<T, U> concept instead of predicate<T, U>, where T and U are reference types of supplied ranges. As consequence, supplied functor needs to be callable with four combinations of arguments: (T, U), (U, T), (U, U), (U, U), instead of just (T, U). This makes them less general than existing non-range overloads, and complicates code migration. | | | | Adopt P1716R2. | | | | |  | | | | | |
| US 313 | | |  | | 25.02 [algorithms.  requirements] | | | |  | | | | Ge | | | | This subclause describes general purpose wording that applies to all algorithms in the standard, without defining algorithm. It generally applies such definitions to algorithms "in this clause", but the wording for specialized algorithms in 20.10.11 [specialized.algorithms] relies on this wording too, especially to provide definitions for its ranges overloads. | | | | Revise this subclause to include a definition of algorithm, so that all the wording that applies to this subclause instead applies to all algorithms. Reasonable definitions of algorithm (for the purposes of the standard library) might be all function templates in clause 25, and clause 20.10.11, or some definition constructed around function templates having arguments of iterator or range type. The former is likely a simpler fix for C++20, the latter would avoid having to update the list of locations for algorithms in the future, | | | | |  | | | | |
| JP8 314 | | |  | | 25.03.1 | | | |  | | | | ed | | | | Notes in terminological entries should start with different element, namely,  "Note # to entry", according to Clauses 24 and 16.5.9 in the Directives Part 2. | | | | Replace "[Note:" with "[Note 1 to entry:". | | | | |  | | | | |
| GB 315 | | |  | | 25.04 | | | |  | | | | Te | | | | next\_permutation\_result has no conversion operators  The other \*\_result classes have conversion operators defined (for the case where the range-based overload returns dangling) but next\_permutation\_result does not. It is also missing the [[no\_unique\_address]] attribute for its iterator member. | | | | In [algorithm.syn], change  template<class I>  struct next\_permutation\_result {  bool found;  [[no\_unique\_address]] I in;  template <class I2>  requires convertible\_to<const I&, I2>  operator next\_permutation\_result<I2>() const & {  return {found, in};  }  template <class I2>  requires convertible\_to<I, I2>  operator next\_permutation\_result<I2>() && {  return {found, std::move(in)};  }  }; | | | | |  | | | | |
| GB 316 | | |  | | 25.04 | | | |  | | | | Te | | | | Algorithm result types should be distinct types; not aliases  Each algorithm should have its own result type (that might be derived from some common exposition-only type).  It will be probably be confusing for a diagnostic to report copy\_result as the return type when a user is using move\_backward, or for mismatch\_result to appear when a user is using swap\_ranges! | | | | Add three exposition-only types:  template<class I1, class I2>  struct in1-in2-result {  [[no\_unique\_address]] I1 in1;  [[no\_unique\_address]] I2 in2;  template<class II1, class II2>  requires convertible\_to<const I1&, II1> && convertible\_to<const I2&, II2>  operator in1-in2-result<II1, II2>() const & {  return {in1, in2};  }  template<class II1, class II2>  requires convertible\_to<I1, II1> && convertible\_to<I2, II2>  operator in1-in2-result<II1, II2>() && {  return {std::move(in1), std::move(in2)};  }  };  template<class I, class O>  struct in-out-result {  [[no\_unique\_address]] I in;  [[no\_unique\_address]] O out;  template<class I2, class O2>  requires convertible\_to<const I&, I2> && convertible\_to<const O&, O2>  operator in-out-result<I2, O2>() const & {  return {in, out};  }  template<class I2, class O2>  requires convertible\_to<I, I2> && convertible\_to<O, O2>  operator in-out-result<I2, O2>() && {  return {std::move(in), std::move(out)};  }  };  template<class I1, class I2, class O>  struct in1-in2-out-result {  [[no\_unique\_address]] I1 in1;  [[no\_unique\_address]] I2 in2;  [[no\_unique\_address]] O out;  template<class II1, class II2, class OO>  requires convertible\_to<const I1&, II1> &&  convertible\_to<const I2&, II2> && convertible\_to<const O&, OO>  operator in1-in2-out-result<II1, II2, OO>() const & {  return {in1, in2, out};  }  template<class II1, class II2, class OO>  requires convertible\_to<I1, II1> &&  convertible\_to<I2, II2> && convertible\_to<O, OO>  operator in1-in2-out-result<II1, II2, OO>() && {  return {std::move(in1), std::move(in2), std::move(out)};  }  };  Each of the algorithm\_result types should be privately derived from the relevant exposition-only type, with the members being made publicly available.  next\_permutation\_result should be renamed to permutation\_result. (Note that for\_each\_result, partition\_copy\_result, minmax\_result, and next\_permutation\_result don't have an exposition-only type, since their use-cases are mostly unique.) | | | | |  | | | | |
| FR 317 | | |  | | 25.06.14 | | | |  | | | | **te** | | | | The names of shift\_left and shift\_right will be misleading when specialized for bit proxy iterators (shift\_left will call the bitwise right shift operator >>, and shift\_right will call the bitwise left shift operator <<). As a result, the names of the algorithms shift\_left and shift\_right would benefit from being adjusted to less misleading names. | | | | Alternatives include having only one shift algorithm taking a signed integer for the shift amount, and shifting to the beginning for a negative amount, and to the end for a positive amount. Changing names is another alternative:  shift\_next/shift\_prev,  backshift/foreshift, back\_shift/fore\_shift,  shift\_forward/shift\_backward,  shift\_front/shift\_back,  shift\_begin/shift\_end | | | | |  | | | | | |
| US 318 | | |  | | 25.06.14  [alg.shift] | | | |  | | | | te | | | | Currently, shift by a negative value is simply ignored. This runs totally counter to user expectation. Make it a precondition to provide a non-negative value. This would be a behaviour change so needs to be done now. | | | | Adopt [P1243](https://wg21.link/p1243).  If not that, adopt [P1233](https://wg21.link/p1233). | | | | |  | | | | |
| JP9 319 | | |  | | 25.07.8 | | | | p22 | | | | ed | | | | "Let X be the return type. Returns Xx, y, where ..." needs braces for constructing the value. In addition, other descriptions for "Returns" doesn't specify the return type explicitly. It would be better to make consistent. | | | | {x, y}, where ... | | | | |  | | | | |
| US 320 | | |  | | 25.08  [numeric.  ops.  overview] | | | |  | | | | te | | | | We made lots of algorithms constexpr, but not the ones in <numeric>. We really should be more thorough and not just forget these. | | | | Adopt [P1645](https://wg21.link/p1645). | | | | |  | | | | |
| FR 321 | | |  | | 25.09 | | | |  | | | |  | | | | All non-allocating algorithms have been made constexpr except for the ones in numeric, which seems like an oversight | | | | Adopt P1645 | | | | |  | | | | | |
| PL 322 | | |  | | 25.09 [numeric.ops] | | | |  | | | | te | | | | ● The specification of GENERALIZED\_\*\_SUM is overly restrictive and suggests that a serial cutoff is not allowed.  ● The intermediate type for numeric algorithms is unclear.  ● The type requirements for numeric algorithms are unclear.  ● The requirements on function objects for numeric algorithms are overly restrictive. | | | | Adopt P0571. | | | | |  | | | | | |
| FR 323 | | |  | | 26.05 | | | |  | | | | **te** | | | | The current wording of low level bit manipulation functions specified by "P0553R4: Bit operations" and by "P0556R3: Integral power-of-2 operations" make these functions unusable with std::byte. Since the whole purpose of the introduction of std::byte was to break the ambiguity between pure integers vs memory storage, the incompatibility between low level bit functions and std::byte reintroduces this pre-C++17 ambiguity for users. The current design may make future evolution of the <bit> header more complicated. | | | | Introduce a machine word “concept” as well as related type traits (binary\_digits, is\_word) that unsigned integers, extended unsigned integers, and std::byte satisfy. This mechanism should also constitute a customization point for advanced users who want to provide their own words types. Have the low-level bit operations take machine words as inputs and not only unsigned integers. See paper P1856R1. | | | | |  | | | | | |
| US 324 | | |  | | 26.05  Applies to §26.5 Numerics Library, Bit manipulation [bit] | | | | §26.5.4 Integral powers of 2 [bit.pow.two];  §26.5.5 Rotating [bit.rotate];  §26.5.6 Counting [bit.count] | | | | te | | | | By strict interpretation of the wording, none of these new bit-oriented interfaces work with std::byte. Given the discussion in the minutes it seems like this was something intentionally left out, to be patched later. A cast to unsigned integral type for std::byte makes a type which already suffers from lack of math operations and similar even more verbose when working with operations it is absolutely supposed to apply to. This is not a useful restriction. | | | | For sections §26.5.4 Integral powers of 2 [bit.pow.two]; §26.5.5 Rotating [bit.rotate]; §26.5.6 Counting [bit.count], change the "Remarks" text to be as follows -- "Remarks: This function shall not participate in overload resolution unless T is an unsigned integer type ([basic.fundamental]) or byte ([cstddef.syn]). If T is byte, then the expression is equivalent to std:: ???(static\_cast<underlying\_type\_t<byte>>(value))." Substitute the name of each function from the section for ???. | | | | |  | | | | |
| US 325 | | |  | | 26.05 [bit] | | | |  | | | | Ed | | | | The contents of the header <bit> relate to inspecting and manipulating memory patterns directly, rather than numeric operations. It better belongs under clause 20, general utilities. | | | | Move 26.5 [bit] to a new subsection under class 20 [utilities] | | | | |  | | | | |
| PL 326 | | |  | | 26.05 [bit] | | | |  | | | | te | | | | The name of std::log2p1 clashes with an IEEE-754 algorithm of the same name, which has been included in a C TS meant for inclusion in a future C standard (and probably also in C++). There is no ambiguity between the overloads, however having a name overloaded for two completely different mathematical formulas is not a good thing. For reference, the log2p1 algorithm from the CD is `log2p1(x) = (x == 0) ? 0 : 1 + floor(log2(x))`, while the log2p1 algorithm from IEEE-754 `log2p1(x) = log2(1 + x)`.  Additionally, the function the CD calls log2p1 is commonly - at the very least in Python, Ruby, Java, and Dart - called `bit\_length`. To the author of this comment, it seems preferable to follow the name that is used in other languages, to make it easier for programmers to communicate and reason about code across the different languages. | | | | Rename `std::log2p1` to `std::bit\_length`. | | | | |  | | | | | |
| US 327 | | |  | | 26.05.4  [bit.pow.two] | | | |  | | | | ed | | | | `log2p1` collides with an IEEE-754 operation. | | | | Rename `log2p1`. We suggest a different semantically meaningful name such as `bit\_width` or `base2digits`. | | | | |  | | | | |
| US 328 | | |  | | 26.05.4  [bit.pow.two] | | | |  | | | | ed | | | | `ceil2` and `floor2`'s names are unintuitive. Meaning that most programmers reading the code won't know what's meant. | | | | Rename `ceil2` and `floor2`. | | | | |  | | | | |
| US 329 | | |  | | 26.05.4  [bit.pow.two] | | | |  | | | | te | | | | The behavior of `ceil2` and `floor2` at 0 is unlikely to be something programmers use correctly. | | | | Change the constraints for both functions around 0. | | | | |  | | | | |
| US 330 | | |  | | 26.05.4  [bit.pow.two] | | | |  | | | | ed | | | | The specification says: "The minimal value `y` such that `ispow2(y)` is `true` and `y >= x`; if `y` is not representable as a value of type `T`, the result is an unspecified value.“ `y` is an argument to `ispow2(y)`. It is necessarily representable.  Note - the above words are not in the CD, SC22 N5410 (WG21 N4830). See 26.5.4 | | | | Rephrase. Better wording might be "if no such `y` exists", but that doesn't seem particularly useful: how does one differentiate "no such `y`" from a real answer? | | | | |  | | | | |
| GB 331 | | |  | | 26.05.4 | | | | 3,4,5,6 | | | | Te | | | | std::ceil2() & std::floor2() produce conceptual confusion  std::ceil() is a linear operation, but std::ceil2() is an exponential operation. The '2' suffix does not provide any hint as to its fundamental difference from the std::ceil() function.  std::floor2() suffers from the same defect.  Spell out what the operations are exactly by renaming to ceil\_power\_of\_two() and floor\_power\_of\_two(). | | | | Rename std::ceil2() to std::ceil\_power\_of\_two().  Rename std::floor2() to std::floor\_power\_of\_two(). | | | | |  | | | | |
| GB 332 | | |  | | 26.05.4 | | | | 7, 8 | | | | Te | | | | std::log2p1() from P0556 introduces a possible name collision.  It is defined as:  log2p1(x) = (x == 0) ? 0 : 1 + floor(log2(x))  In IEEE754-2008, and WG14 TS 18661-4a (targeted for C2X):  log2p1(x) = log2(1+x)  The intention of P0556's log2p1() function is to facilitate bit manipulation algorithms by computing the number of bits needed to represent an unsigned integer. Give it a more descriptive name, such as std::bits\_needed(). | | | | Rename std::log2p1() to std::bits\_needed(). | | | | |  | | | | |
| GB 333 | | |  | | 27.02.2.1 | | | |  | | | | Te | | | | UTC epoch is not correctly defined  UTC has an officially recorded epoch of 1/1/1972 00:00:00 and is 10 seconds behind TAI.  This can be confirmed through reference to the BIPM (the body that oversees international metrology)  ​https://www.bipm.org/cc/CCTF/Allowed/18/CCTF\_09-32\_noteUTC.pdf  Specifically page 6  "The defining epoch of 1 January 1972, 0 h 0 m 0 s UTC was set 10 s behind TAI, which was the approximate accumulated difference between TAI and UT1 since the inception of TAI in 1958, and a unique fraction of a second adjustment was applied so that UTC would differ from TAI by an integral number of seconds. The recommended maximum departure of UTC from UT1 was 0.7 s. The term “leap second” was introduced for the stepped second." | | | | utc\_clock and utc\_timepoint should correctly report relative to the official UTC epoch.  27.2.2.1 footnote 1 should read  In contrast to sys\_time, which does not take leap seconds into account, utc\_clock and its associated time\_point, utc\_time, count time, including leap seconds, since 1972-01-01 00:00:00 UTC. [Example:  clock\_cast<utc\_clock>(sys\_seconds{sys\_days{1972y/January/1}}).time\_since\_epoch() is 0s.  clock\_cast<utc\_clock>(sys\_seconds{sys\_days{2000y/January/1}}).time\_since\_epoch()  is 883'612'822, which is 10’197 \* 86’400s + 22s.  — end example] | | | | |  | | | | |
| US 334 | | |  | | 27.05.10  [time.  duration.io] | | | |  | | | | te | | | | operator<< for floating-point durations always produces output with six digits after the decimal point, and doesn’t use the stream’s locale either. | | | | Rewrite the specification to not rely on to\_string() for floating-point formatting. | | | | |  | | | | |
| GB 335 | | |  | | 27.07.1.1 | | | |  | | | |  | | | | Wording for clocks should be unified unless they are intended to behave differently  In 27.7.1.1 note 1 for system\_clock it is stated  "Objects of type system\_clock represent wall clock time from the system-wide realtime clock. Objects of  type sys\_time<Duration> measure time since (and before) 1970-01-01 00:00:00 UTC"  The express statement of "since (and before)" is important given the time epoch of these clocks. If all the clocks support time prior to their zero-time then this should be stated explicitly. If not then likewise that should be noted. No change is proposed yet, clarification required over the intended behaviour when using values prior to a given clock's epoch is needed before the appropriate change can be suggested. | | | | Unify the wording | | | | |  | | | | |
| GB 336 | | |  | | 27.07.4 | | | |  | | | | Te | | | | Use of specific clocks may create expectations that are not delivered (GPS)  The "gps" clock has nothing to do with the GNSS service known as GPS except for sharing a common anchor point (epoch) there is no calibration, no feed, no expectation that the "clock" correlates to GPS data streams. It seems a very niche use and given some of the other issues around its interpretation I would suggest it is removed. | | | | delete gps\_clock | | | | |  | | | | |
| GB 337 | | |  | | 27.07.4 | | | |  | | | | Te | | | | gps\_clock is a unilateral reference to a US service and has no place alone in the ISO standard  The GPS GNSS service is owned, maintained and controlled by the US government and while Satellite timing and position usage has become all but ubiquitous in many applications a modern GNSS receiver is capable of receivng updates from multiple GNSS constellations to ensure coverage and security. If GPS is represented then other national and interenationally maintained services should be included, Glonass, Beidou, Galileo to name but 3. Each of these have different operating parameters, most notably the epoch and the application or not of leap seconds.  ​https://gssc.esa.int/navipedia/index.php/Time\_References\_in\_GNSS  While Galileo and Beidou do not respect leap seconds, GLONASS does and GLONASS transmits at a constant offset of 3 hours relative to UTC (being Russian standard time).  Other features of gps, such as the week rollover (an epoch defining event and which occurs every 19 years, the most recent being this April 2019) are not represented in the gps\_clock. If one purpose of providing a gps\_clock is to allow the comparison of gps data to other clocks then epoch rollover probably ought to be recognised, though hopefully it will become a thing of the past as increasingly gps satellites are upgrading to a large week counter) | | | | Remove gps\_clock | | | | |  | | | | |
| JP10 338 | | |  | | 27.08.3.3 | | | | p10 | | | | ed | | | | This is different from the declaration in 27.2. | | | | constexpr chrono::day operator""d(unsigned long long d) noexcept; | | | | |  | | | | |
| JP12 339 | | |  | | 27.08.4.2 | | | | p14 | | | | ed | | | | Class name is not required. | | | | constexpr explicit ~~month::~~operator unsigned() const noexcept; | | | | |  | | | | |
| JP13 340 | | |  | | 27.08.4.2 | | | | p15 | | | | ed | | | | Class name is not required. | | | | constexpr bool ~~month::~~ok() const noexcept; | | | | |  | | | | |
| JP11 341 | | |  | | 27.08.4.2 | | | | p2 | | | | ed | | | | Class name is not required. | | | | constexpr month& ~~month::~~operator++() noexcept; | | | | |  | | | | |
| GB 342 | | |  | | 27.08.4.3 | | | |  | | | | Te | | | | std::chrono::month is the only duration-like type without a UDL, which makes constructing objects such as this look a little off.  chrono::year\_month\_day{1815y, chrono::month{12}, 10d}} | | | | Add operator""month, which behaves similarly to operator""y.  The earlier expression could then read as:  chrono::year\_month\_day{1815y, 12month, 10d}} | | | | |  | | | | |
| JP14 343 | | |  | | 27.08.5.3 | | | | p10 | | | | ed | | | | This is different from the declaration in 27.2. | | | | constexpr chrono::year operator""y(unsigned long long y) noexcept; | | | | |  | | | | |
| DE 344 | | |  | | 27.11.01 | | | | paragraph 1 | | | | ge | | | | This paragraph says "27.11 describes an interface for accessing the IANA Time Zone database described in RFC 6557, ..."  However, RFC 6557 does not describe the database itself; it only describes the maintenance procedures for that database, as its title implies (quoted in clause 2). | | | | Add a reference to a specification of the database itself, or excise all references to the IANA time zone database. | | | | |  | | | | |
| DE 345 | | |  | | 27.11.08 | | | |  | | | | te | | | | The class name "leap" to designate a UTC leap second event is too generic and not sufficiently descriptive. | | | | Rename the class to "utc\_leap" or "utc\_leap\_second", consistent with the naming of "utc\_clock" and "utc\_time" for other UTC-related classes. | | | | |  | | | | |
| DE 346 | | |  | | 27.11.09 | | | |  | | | | te | | | | The class name "link" to designate an alias for a named time zone is too generic and not sufficiently descriptive. | | | | Rename the class to "zone\_link", consistent with the fact that all other classes related to time zones contain "zone" in their name. | | | | |  | | | | |
| FR 347 | | |  | | 27.12.10 | | | |  | | | |  | | | | local\_time\_format should use optional<string> and optional<seconds> instead of pointers | | | | change local\_time\_format signature to local\_time\_format(local\_time<Duration> time, optional<string> = {},  optional<seconds> offset\_sec = {}); | | | | |  | | | | | |
| JP15 348 | | |  | | 29.10.01 | | | |  | | | | ed | | | | The default template arguments are missing. | | | | template<class charT, class traits = char\_traits<charT>, class Allocator = allocator<charT>> class basic\_syncbuf; template<class charT, class traits = char\_traits<charT>, class Allocator = allocator<charT>> class basic\_osyncstream; | | | | |  | | | | |
| JP16 349 | | |  | | 29.10.02.1 | | | |  | | | | ed | | | | The default template arguments are missing. | | | | template<class charT, class traits = char\_traits<charT>, class Allocator = allocator<charT>> class basic\_syncbuf : public basic\_streambuf<charT, traits> { | | | | |  | | | | |
| JP17 350 | | |  | | 29.10.03.1 | | | |  | | | | ed | | | | The default template arguments are missing. | | | | template<class charT, class traits = char\_tratis<charT>, class Allocator = allocator<charT>> class basic\_osyncstream : public basic\_ostream<charT, traits> { | | | | |  | | | | |
| US 351 | | |  | | 31  [atomics] | | | |  | | | | te | | | | Atomic initialization has been broken since C++11. | | | | Adopt [P0883R1](https://wg21.link/p0883r1). | | | | |  | | | | |
| US 352 | | |  | | 31  [atomics] | | | |  | | | | te | | | | It is not possible to include C’s <stdatomic.h> in C++ today, which makes it difficult to use atomics in code that needs to be compiled as both C and C++. C++ should support inclusion of <stdatomic.h>. | | | | Adopt [P0943](https://wg21.link/P0943). | | | | |  | | | | |
| CA 353 | | |  | | 31 [atomics] | | | |  | | | | te | | | | Atomic initialization doesn’t work as expected. | | | | Adopt P0883R1. | | | | |  | | | | |
| US 354 | | |  | | 31.03  [atomics.  alias] | | | |  | | | | Te | | | | In [P1135r3](https://wg21.link/P1135r3), the atomic\_int\_fast\_wait\_t and atomic\_uint\_fast\_wait\_t type aliases were removed. The paper’s changelog explains why:  Removed atomic\_int\_fast\_wait\_t and atomic\_uint\_fast\_wait\_t, because LEWG at San Diego 2018 felt that the use case was uncommon and the types had high potential for misuse.  We think this decision warrants reconsideration. On some platforms, certain implementation strategies for wait/notify are only available for certain sized integer types (for example, Linux’s futex is for int only) | | | | Re-add [P1135r2](https://wg21.link/P1135r2)’s atomic\_int\_fast\_wait\_t and atomic\_uint\_fast\_wait\_t. | | | | |  | | | | |
| US 355 | | |  | | 31.07  [atomics.ref.generic]  31.7.1  [atomics.ref.operations]  31.7.2  [atomics.ref.  int]  31.7.3  [atomics.ref.  float]  31.7.4  [atomics.ref.pointer] | | | | 25-28  1  1    1 | | | | te | | | | atomic\_ref<T>::notify\_one and atomic\_ref<T>::notify\_all should be const member functions (in the generic class and all the specializations), since it is the atomic object that is (conceptually) modified, not the atomic\_ref<T> object. | | | | Make atomic\_ref<T>::notify\_one and atomic\_ref<T>::notify\_all const. | | | | |  | | | | |
| US 356 | | |  | | 31.07.1  [atomics.ref.operations] | | | | 10 | | | | te | | | | atomic\_ref::is\_lock\_free should require that the result only depend on the type of the object, not the specific object. The current specification is inconsistent with atomic<T>. (See [atomics.lockfree] 31.5p3.) This test is primarily useful to determine whether a particular algorithm can or should be used. If the result can vary based on object identity, that is not possible. There is no way to ask whether the property holds for all relevant objects until all of the objects are actually available for testing. Note that is\_always\_lock\_free does not fully serve this purpose, since is\_lock\_free() may vary at run time depending on hardware characteristics, while still being consistent per type. This was the subject of recent reflector discussion. | | | | Apply the PR for LWG3249, and replicate the equivalent wording here.  Possibly consider making the is\_lock\_free member function static. That appears to make the member function behavior less surprising, at the cost of an inconsistency with the C-constrained free function. | | | | |  | | | | |
| US 357 | | |  | | 31.07.2 [atomics.ref.  int]  31.8.2  [atomics.  types.int] | | | | 1  1 | | | | ed | | | | The note at the end of [atomics.ref.int] (31.7.2) paragraph 1:  [ *Note*: For the specialization atomic\_ref<bool>, see 31.7. — *end note* ]  refers to [atomics.ref.generic](31.7). There is no mention of atomic\_ref<bool> in that subclause, so the reason for the cross reference is not obvious.  The note at the end of [atomics.types.int] (31.8.2) paragraph 1:  [ *Note*: For the specialization atomic<bool>, see 31.8. -- *end note* ]  has a similar issue. [atomics.types.generic] (31.8) does mention atomic<bool>, but not in a way that makes the reason for the cross reference obvious. | | | | Change the note at the end of [atomics.ref.int] (31.7.2) paragraph 1 as follows:  [ *Note*: For the specialization atomic\_ref<bool>, see 31.7. — *end note* ]  [ *Note*: The specialization atomic\_ref<bool> is based on the primary template ([atomic.ref.generic]), and is not included among the integer-specific specializations. — *end note* ]  Change the note at the end of [atomics.types.int] (31.8.2) paragraph 1 as follows:  [ *Note*: For the specialization atomic<bool>, see 31.8. -- *end note* ]  [ *Note*: The specialization atomic<bool> is based on the primary template ([atomics.types.generic]), and is not included among the integer-specific specializations. -- *end note* ] | | | | |  | | | | |
| US 358 | | |  | | 31.07.3  [atomics.ref.  float] | | | | 1 | | | | te | | | | In the atomic\_ref<*floating-point*> synopsis in [atomic.ref.float] (31.7.3) paragraph 1:  *floating-point* operator=(*floating-point*) noexcept;  should be a const member function, like all other atomic\_ref<T> assignment operators. | | | | Change the atomic\_ref<*floating-point*> synopsis in [atomic.ref.float] (31.7.3) paragraph 1 as follows:  *floating-point* operator=(*floating-point*) const noexcept; | | | | |  | | | | |
| US 359 | | |  | | 31.07.5  [atomics.ref.memop] | | | | 1-4 | | | | ed | | | | In the specification of member operations common to atomic\_ref<*integral*> and atomic\_ref<T\*> specializations in [atomics.ref.memop] (31.7.5), all of the member functions are specified to return T\*, which is only correct for the atomic\_ref<T\*>. The corresponding member operations for atomic<*integral*> and atomic<T\*> in [atomics.types.memop] (31.8.5) return *T*.  Additionally, there is an extra int parameter in the specification of the second operator--. That declaration is supposed to be the predecrement operator, not the postdecrement operator. | | | | Change the specification of member operations common to atomic\_ref<*integral*> and atomic\_ref<T\*> specializations in [atomics.ref.memop] (31.7.5) as follows:  T\* operator++(int) const noexcept;  *Effects:* Equivalent to: return fetch\_add(1);  T\* operator--(int) const noexcept;  *Effects:* Equivalent to: return fetch\_sub(1);  T\* operator++() const noexcept;  *Effects:* Equivalent to: return fetch\_add(1) + 1;  T\* operator--(int) const noexcept;  *Effects:* Equivalent to: return fetch\_sub(1) - 1;  Alternatively, we could consider rewording these member operations entirely for both atomic<T> and atomic\_ref<T>. For example, we could just add them to both the integer and pointer specializations, which would be clearer, but would duplicate the wording. | | | | |  | | | | |
| GB 360 | | |  | | 32 | | | |  | | | | Te | | | | Too many new headers  The CD includes new headers <stop\_token>, <semaphore>, <latch>, and <barrier>. It seems unhelpful to split the synchronization utilities across so many different headers. | | | | Consolidate them into a single <sync> header or similar. | | | | |  | | | | |
| GB 361 | | |  | | 32.04.3 | | | |  | | | | Te | | | | Rename jthread  The name jthread, while concise, is cryptic. By expanding the j, the 'smart-thread' type will be in sync with its smart-pointer cousins. | | | | Proposed Change:  Consider expanding jthread into a name that more directly represents its intention. Examples include, but are not limited to:  joining\_thread  join\_thread | | | | |  | | | | |
| JP18 362 | | |  | | 32.04.3.5 | | | |  | | | | ed | | | | This is different from the declaration in 32.4.3. | | | | [[nodiscard]] unsigned hardware\_concurrency() noexcept; | | | | |  | | | | |
| PL 363 | | |  | | 32.06.4 [thread.condition.condvarany] | | | |  | | | | te | | | | The conditition\_variable\_any::wait\_until that accepts lock and stop\_token, is inconsistent with the [thread.req.timing] p4, that reserves '\_until' suffix for functions that accepts time\_point. Furthermore, all interruptible waits functions, are accepting stop\_token as the last argument, following the predicate, thus making them harder to format in situations when lambda is passed as a predicate. | | | | Change the interruptible waits interface as proposed in P1869R0. | | | | |  | | | | | |
| US 364 | | |  | | 32.07.2 | | | | Paragraph 13 | | | | te | | | | The phrasing of the spurious failure case of semaphore try\_acquire can confuse readers, who may parse it as being about blocking guarantees or a statement about QoI, rather than capturing various memory model subtleties as intended. Better would be to word this case similarly to mutex try\_lock. The proposed change does so. | | | | Replace with "Effects: Attempts to atomically check if the counter is positive and decrement it by one if so, without blocking. If the counter is not decremented, there is no effect and try\_acquire immediately returns. An implementation by fail to decrement the counter even if it is positive. [ Note: This spurious failure is normally uncommon, but allows interesting implementations based on a simple compare and exchange ([atomic]). -- end note] An implementation should ensure that try\_acquire() does not consistently return false in the absence of contending semaphore operations.“ | | | | |  | | | | |
| US 365 | | |  | | 32.08  [thread.  coord] | | | |  | | | | te | | | | latch and barrier currently take a ptrdiff\_t as their expected count parameter and thus must support any expected count (larger than or equal to 0) that will fit in a ptrdiff\_t. This limits implementation freedom; some platforms can provide a much more efficient implementation of latch and barrier if they can restrict the maximum possible expected count. | | | | Adopt [P1865](https://wg21.link/p1865), which adds a static constexpr ptrdiff\_t max() noexcept; member to both classes that returns the expressible range of the object, like the one on counting\_semaphore. | | | | |  | | | | |
| CA 366 | | |  | | 32.08.1 [thread.latch] | | | |  | | | | ed | | | | Subclauses 32.8.2 [latch.syn] and 32.8.3 [thread.latch.class] should be subclauses of subclause 32.8.1 [thread.latch] instead of at the same level (e.g., to be consistent with subclause 32.8.4 [thread.barriers]). | | | | Make subclauses 32.8.2 and 32.8.3 subclauses of 32.8.1. | | | | |  | | | | |
| US 367 | | |  | | 6-15 | | | |  | | | | ge | | | | Requirements that a header be included before a language feature is functional should also allow for importing that header (unit), While we will file additional comments for the cases we identify, this is a catch-all comment to adopt the principle and similarly fix any places we miss. | | | | When a header be included before a language feature is functional should also allow for importing that header (unit). | | | | |  | | | | |
| US 368 | | |  | | 6-15 | | | |  | | | | ge | | | | Undefined behavior lexing and tokenizing the program text has no place in a modern standard. Unnecessary undefined behaviour in our standard raises a wide variety of concerns, not least with the security community, and all concerns related to turning source code into tokens for subsequent analysis should be either diagnosable errors, or (conditionally) supported behavior. This comment is a principle statement for more detailed comments that follow, and as a catch-all for any further occurrences that were missed. | | | | Remove Undefined Behavior lexing and tokenizing the program text. | | | | |  | | | | |
| BG3 369 | | | P 118 | | 7.06.2.3 | | | | 6 | | | | ge | | | | (Related to BG2) The code example uses the void-returning variant of await\_suspend(). | | | | Change the return type of my\_future::await\_suspend() and awaiter::await\_suspend() to coroutine\_handle. | | | | |  | | | | |
| US 370 | | |  | | All clauses affected by P0912R5  Beh ???  9.04.4  et al | | | | n/a | | | | ge | | | | WG21 has received four independent usage reports on efforts to adopt Coroutines in production code: [P0054](https://wg21.link/p0054), [P0973](https://wg21.link/p0973), [P1471](https://wg21.link/p1471), and [P1745](https://wg21.link/p1475). All of these early adopters identified major problems that could not be fixed in a backwards-compatible way; the problems identified in [P0054](https://wg21.link/p0054) were addressed via incompatible changes, and the problems in the other papers remain unaddressed in the CD. On the basis of this experience, we believe it would be premature to standardize Coroutines. | | | | Revert the application of [P0912R5](https://wg21.link/p0912r5) (Merge Coroutines TS into C++20 working draft). | | | | |  | | | | |
| RU 371 | |  | | | | | All the library clauses | | | |  | | | | ge | | | | Apply all the wordings from all the “Mandating the Standard Library” papers. | | | | | Apply wordings from P1505, P1622, P1686, P1718-P1723. | | | | |  | | | | | |
| US 372 | | |  | | Annex C  [c.compat] | | | |  | | | | te | | | | C and C++ atomics haven't worked together properly since first being standardized, even if the intent was for them to interoperate. | | | | Adopt [P0943R4](https://wg21.link/p0943r4). | | | | |  | | | | |
| JP19 373 | | |  | | C.5.01 | | | | p2.6 | | | | ed | | | | It's good to have a reference as in p2.1 to p2.5 | | | | Add a reference to 7.5.7. | | | | |  | | | | |
| NL 374 | | | 3 | | Cross references from ISO C ++ 2017" | | | |  | | | | ed | | | | Typo ‘fmtflatgs’ should be ‘fmtflags’. | | | | Change to ‘fmtflags’ | | | | |  | | | | |
| NL 375 | | | 2 | | D.19 [depr.fs.path.factory] | | | | Sub 4 | | | | te | | | | Example in deprecated section implies that std::string is the type to use for utf8 strings.  [Example: A string is to be read from a database that is encoded in UTF-8, and used to create a  directory using the native encoding for filenames:  namespace fs = std::filesystem;  std::string utf8\_string = read\_utf8\_data();  fs::create\_directory(fs::u8path(utf8\_string)); | | | | Add clarification that std::string is the wrong type for utf8 strings | | | | |  | | | | |
| JP20 376 | | |  | | D.7 | | | | p1 | | | | ed | | | | 9.4 is less appropriate as a reference for "these implicit definitions could become deleted" | | | | Replace 9.4 with 9.4.3. | | | | |  | | | | |
| CA 377 | | |  | | General | | | |  | | | | te | | | | C / C++ interop for atomics is buggy. | | | | Adopt P0943. | | | | |  | | | | |
| CA 378 | | |  | | General | | | |  | | | | ge | | | | How constraints work with non-templated functions is still under heavy construction during this late stage in the process. While we have provided various comments that build in a direction where supporting such constructs (including ordering between multiple constrained functions based on their constraints) would become possible, we acknowledge that WG 21 might not find a solution with consensus in time for the DIS. We ask WG 21 to evaluate the risk of shipping the feature in such a state and consider removing the ability to declare such functions. | | | |  | | | | |  | | | | |

Programming languages -- C++